

# THE MAXIMAL DIMENSIONS OF PATH AND GRAPH ALGEBRAS

Łukasz Kaczmarczyk joint work with  
P.M. Hajac, M. Lowiel, E.A. Pacheco

University of Warsaw

29.07.2026

# Notation

## Definition: Directed Graph

A directed graph is a quadruple  $E := (E^0, E^1, s, t)$  consisting of

- ① the set of vertices  $E^0$  and the set of edges  $E^1$ ,
- ② the source and target maps  $s, t: E^1 \rightarrow E^0$  assigning to each edge its source and target vertex, respectively.

## Definition: Sources and Sinks

A vertex  $v$  in a directed graph  $E$  is called

- a source if  $t^{-1}(\{v\}) = \emptyset$ ,
- a sink, if  $s^{-1}(\{v\}) = \emptyset$ .

# Notation

## Definition: Path

A finite path in a directed graph  $E$  is a single vertex or a finite tuple  $p_n := (e_1, \dots, e_n)$  of edges satisfying

$$t(e_1) = s(e_2), \quad t(e_2) = s(e_3), \quad \dots, \quad t(e_{n-1}) = s(e_n).$$

- The beginning  $s(p_n)$  of  $p_n$  is  $s(e_1)$ ,
- The end  $t(p_n)$  of  $p_n$  is  $t(e_n)$ ,
- The length of a path is 0 when it's a vertex or the size of the tuple otherwise,
- Path  $p_n$  is called a cycle if  $s(p_n) = t(p_n)$ ,
- A directed graph without cycles is called acyclic,

# Notation

## Definition: Connectedness

The directed graph  $(E^0, E^1, s, t)$  is connected if it is connected as an undirected graph.

## Definition: Forests, Trees and Trunks

Let  $(E^0, E^1, s, t)$  be an acyclic graph with  $N \geq 1$  edges, then

- if it has no isolated vertexes we call it a *N-forest graph*,
- if it's a connected *N-forest* we call it a *N-tree graph*,
- if its set of vertexes can indexed by  $\{1, \dots, n\}$  and satisfy:

$$\forall e \in E^1 \exists 1 \leq i \leq n-1 \text{ } s(e) = v_i \text{ and } t(e) = v_{i+1}$$

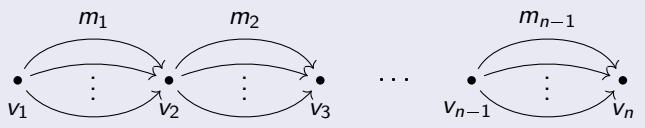
we call it a *N-trunk graph*.

We denote the set of all forest, tree and trunk graphs with  $N \geq 1$  edges by  $\mathcal{F}^N$ ,  $\mathcal{T}^N$ , and  $\mathcal{J}^N$ , respectively.

# Notation

## Presentation of trunks

Given a trunk  $E$  with vertices  $E^0 = \{v_1, \dots, v_n\}$  we will depict it in the following way:



The trunk graph as in the above picture may be described by its associated sequence  $(m_1, \dots, m_{n-1})$ .

## Induced graphs: Trunks

If  $E$  is a trunk graph with  $E^0 = \{v_1, \dots, v_n\}$  then for  $1 \leq i \leq j \leq n$  let  $E_{i,j}$  denote a full subgraph of  $E$  induced by vertices  $\{v_i, \dots, v_j\}$ .

# Calculating the number of paths

A. Chirvasitu, P.M. Hajac 2019

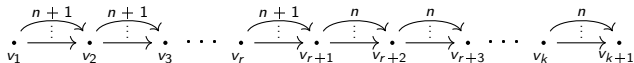
Let  $E$  be an acyclic directed graph with  $N \geq 1$  edges, and let

$$1 \leq k \leq N =: nk + r, \quad 0 \leq r \leq k - 1.$$

Then there are at most

$$P_k^N := (n + 1)^r n^{k-r}$$

different paths of length  $k$ , and the bound is optimal.



## Definition: Path algebra

### Definition: Path algebra of a graph

Let  $E = (E^0, E^1, s, t)$  be a directed graph,  $\mathbb{K}$  a field.

Let's consider the free algebra  $\mathbb{K}\langle E^0 \sqcup E^1 \rangle$  together with relations:

- $(V) \quad \forall v, u \in E^0 \quad vu = \delta_{v,u}v,$
- $(E) \quad \forall e \in E^1 \quad s(e)e = e = et(e),$

then the algebra  $\mathbb{K}\langle E^0 \sqcup E^1 | V, E \rangle$  is called the *path algebra* of  $E$  over  $\mathbb{K}$  and will be denoted by  $\mathbb{K}E$ .

# Notation

## Number of paths

Let  $E$  be a forest graph, then

- $P(v, w)$  - number of paths from  $v$  to  $w$ ,
- $P_s(v)$  - amount of paths beginning in  $v$ ,
- $P_t(v)$  - number of paths ending in  $v$ ,
- $P(E)$  - number of paths in  $E$ .

Moreover if  $E$  has unique source/sink then

- $P_s(E)$  - number of paths beginning at the source of  $E$ ,
- $P_t(E)$  - number of paths ending at the sink of  $E$ .

## Fact:

For a directed graph  $E$  the following equality holds:

$$\dim \mathbb{K}E = P(E).$$



## Calculating the number of paths

P.M. Hajac, O.M. Stachowiak

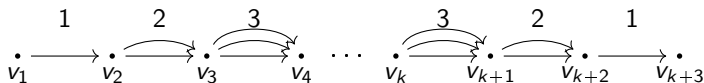
For  $E \in \mathcal{E}^N$  the following inequality holds:

$$P(E) \leq N \left( e^{\frac{1}{e}} \right)^N.$$

# Three Sisters: Path Algebras

## Theorem: Dimension of the Path Algebra

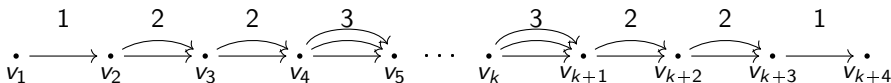
If  $N \geq 16$ ,  $N = 3k$  and  $F$  maximizes the dimension of  $\mathbb{K}E$  among  $\mathcal{E}^N$ , then  $P(F) = \frac{1}{4} (121 \cdot 3^{k-2} - 2k + 7)$



# Three Sisters: Path Algebras

## Theorem: Dimension of the Path Algebra

If  $N \geq 16$ ,  $N = 3k + 1$  and  $F$  maximizes the dimension of  $\mathbb{K}E$  among  $\mathcal{E}^N$ , then  $P(F) = \frac{1}{4} (529 \cdot 3^{k-3} - 2k + 25)$



# Three Sisters: Path Algebras

## Theorem: Dimension of the Path Algebra

If  $N \geq 16$ ,  $N = 3k + 2$  and  $F$  maximizes the dimension of  $\mathbb{K}E$  among  $\mathcal{E}^N$ , then  $P(F) = \frac{1}{4} (253 \cdot 3^{k-2} - 2k + 15)$



## Definition: Levitt Path Algebra

### Definition: Doubled Graph

For a directed graph  $E = (E^0, E^1, s, t)$  we define the doubled graph  $\tilde{E} = (E^0, E^1 \sqcup (E^1)^*, \tilde{s}, \tilde{t})$  via formulas:

- $(E^1)^* = \{e^* : e \in E^1\},$
- $\forall e \in (E^1)^* \quad \tilde{s}(e^*) = t(e) \text{ and } \forall e \in E^1 \quad \tilde{s}(e) = s(e),$
- $\forall e \in (E^1)^* \quad \tilde{t}(e^*) = s(e) \text{ and } \forall e \in E^1 \quad \tilde{t}(e) = t(e).$

# Leavitt Path Algebra

## Definition: Leavitt path algebra of a graph

Let  $E = (E^0, E^1, s, t)$  be a directed graph,  $\mathbb{K}$  a field.

Let's consider the path algebra of the doubled graph  $\tilde{E}$ :

$\mathbb{K}\langle E^0 \sqcup E^1 \sqcup (E^1)^* \mid V, E \rangle$  together with relations:

- (CK1)  $\forall e^* \in (E^1)^* \forall f \in E^1 \quad e^* f = \delta_{e,f} t(e),$
- (CK2)  $\forall v \in E^0 \quad 0 < |s^{-1}(\{v\})| < +\infty \Rightarrow v = \sum_{e \in s^{-1}(\{v\})} ee^*.$

The algebra  $\mathbb{K}\langle E^0 \sqcup E^1 \sqcup (E^1)^* \mid V, E, CK1, CK2 \rangle$  is called the *Leavitt path algebra* of  $E$  over  $\mathbb{K}$  and is denoted by  $L_{\mathbb{K}}(E)$ .

# Leavitt Path Algebras properties

## Theorem

Let  $E$  be an  $N$ -forest graph and  $\{u_i\}_1^n$  be the set of all its sinks. Then the following equality holds

$$L_{\mathbb{K}}(E) = \bigoplus_{i=1}^n M_{P_t(u_i)}(\mathbb{K}).$$

## Lemma

Let  $\mathcal{E}_s^N$  denote the set of all  $N$ -forests with a single sink. Then the following equality holds:

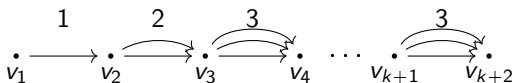
$$\sup_{E \in \mathcal{E}^N} \dim L_{\mathbb{K}}(E) = \sup_{E \in \mathcal{E}_s^N} \dim L_{\mathbb{K}}(E) = \sup_{E \in \mathcal{E}_s^N} (P_t(E))^2.$$

Note that  $\mathcal{T}^N \subset \mathcal{E}_s^N$ .

# Three Sisters: Leavitt Path Algebras

## Theorem: Dimension of the Leavitt Path Algebra

If  $N \geq 8$ ,  $N = 3k$  and  $F$  maximizes the dimension of  $L_{\mathbb{K}}(E)$  among  $\mathcal{E}^N$ , then  $P_t(F) = \frac{1}{2} (11 \cdot 3^{k-1} - 1)$ ,

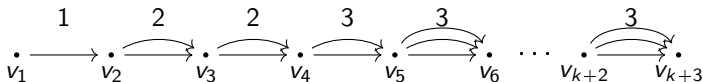




# Three Sisters: Leavitt Path Algebras

## Theorem: Dimension of the Leavitt Path Algebra

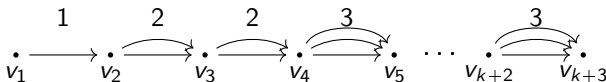
If  $N \geq 8$ ,  $N = 3k + 1$  and  $F$  maximizes the dimension of  $L_{\mathbb{K}}(E)$  among  $\mathcal{E}^N$ , then  $P_t(F) = \frac{1}{2} (47 \cdot 3^{k-2} - 1)$ ,



# Three Sisters: Leavitt Path Algebras

## Theorem: Dimension of the Leavitt Path Algebra

If  $N \geq 8$ ,  $N = 3k + 2$  and  $F$  maximizes the dimension of  $L_{\mathbb{K}}(E)$  among  $\mathcal{E}^N$ , then  $P_t(F) = \frac{1}{2} (23 \cdot 3^{k-1} - 1)$ ,



# Path Algebra

## Path Algebra Case - Methods

# From Trees to Trunks

## Lemma:

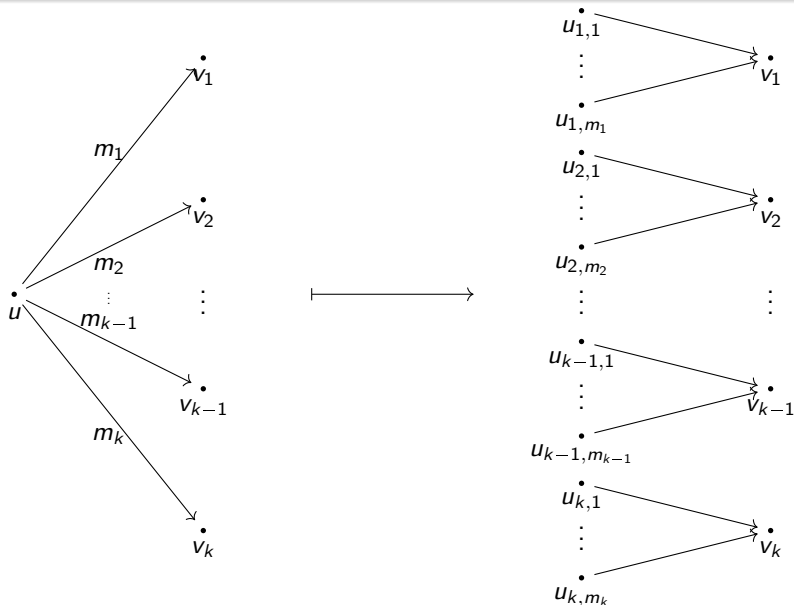
The following equality holds:

$$\sup_{E \in \mathcal{E}^N} P(E) = \sup_{E \in \mathcal{T}^N} P(E).$$

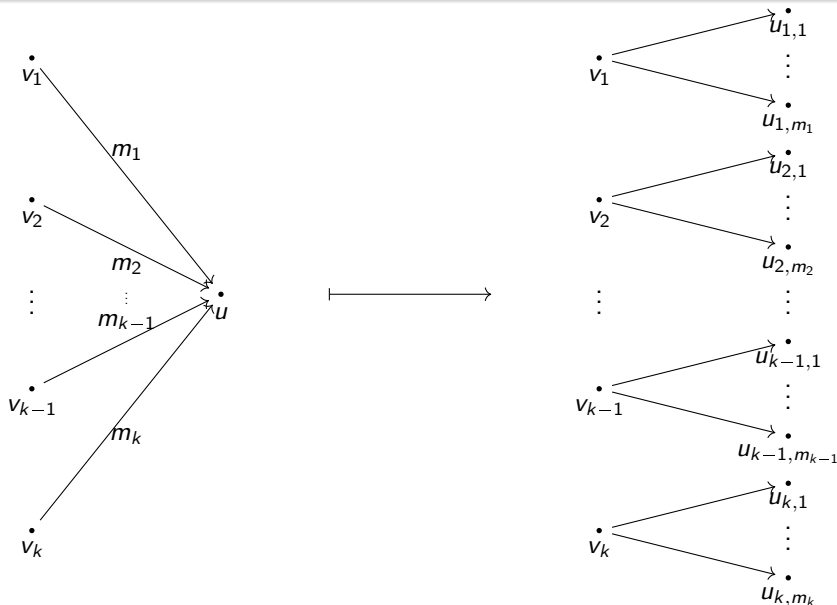
## Moves

- Splitting sinks and sources,
- Resolving alternative paths,
- Combining stray paths,
- Fusion of trunks.

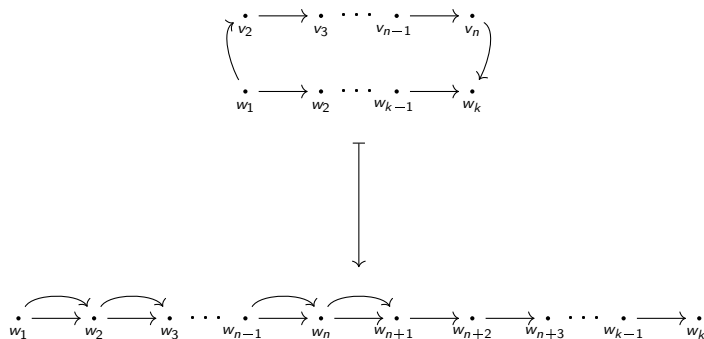
# Moves - Splitting Sources



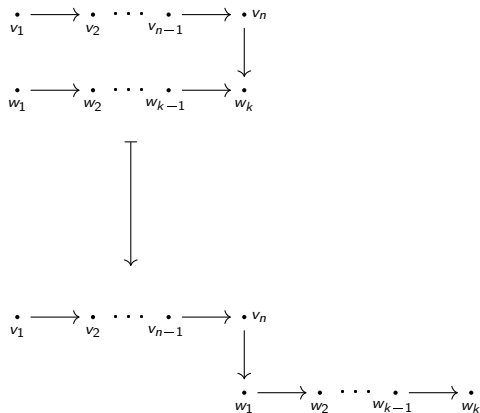
# Moves - Splitting Sinks



# Moves - Resolving alternative paths

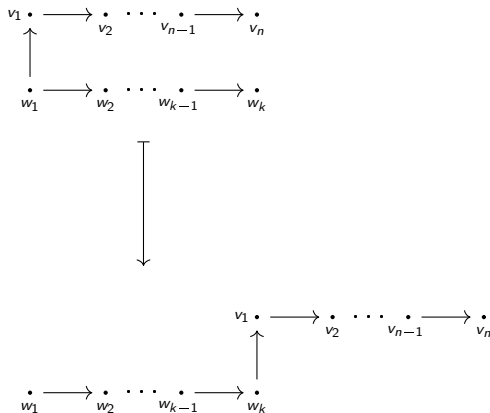


# Moves - Combing stray paths from sources

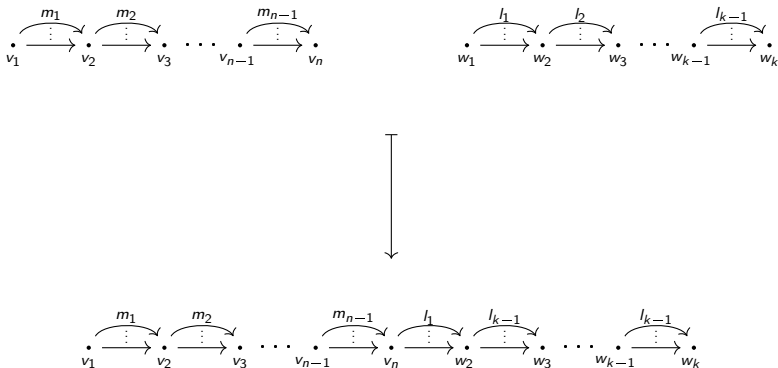




# Moves - Combining stray paths to sinks



# Moves - Fusion of trunks



# Python

| $N$ edges | Maximal number of paths | Optimal trunks                               | Length of longest path |
|-----------|-------------------------|----------------------------------------------|------------------------|
| 3         | 3                       | (1,1,1)                                      | 3                      |
| 4         | 15                      | (1,1,1,1)                                    | 4                      |
| 5         | 21                      | (1,1,1,1,1), (1,1,2,1), (1,2,1,1)            | 4 or 5                 |
| 6         | 31                      | (1,2,2,1)                                    | 4                      |
| 7         | 44                      | (1,1,2,2,1), (1,2,2,1,1)                     | 5                      |
| 8         | 66                      | (1,2,2,2,1)                                  | 5                      |
| 9         | 91                      | (1,1,2,2,2,1), (1,2,2,2,1,1), (1,2,3,2,1)    | 5 or 6                 |
| 10        | 137                     | (1,2,2,2,2,1)                                | 6                      |
| 11        | 192                     | (1,2,2,3,2,1), (1,2,3,2,2,1)                 | 6                      |
| 12        | 280                     | (1,2,2,2,2,2,1)                              | 7                      |
| 13        | 401                     | (1,2,2,3,2,2,1)                              | 7                      |
| 14        | 571                     | (1,2,2,3,3,2,1), (1,2,3,3,2,2,1)             | 7                      |
| 15        | 820                     | (1,2,2,2,3,2,2,1), (1,2,2,3,2,2,2,1)         | 8                      |
| 16        | 1194                    | (1,2,2,3,3,2,2,1)                            | 8                      |
| 17        | 1709                    | (1,2,2,3,3,3,2,1), (1,2,3,3,3,2,2,1)         | 8                      |
| 18        | 2449                    | (1,2,3,3,3,3,2,1)                            | 8                      |
| 19        | 3574                    | (1,2,2,3,3,3,2,2,1)                          | 9                      |
| 20        | 5124                    | (1,2,2,3,3,3,3,2,1), (1,2,3,3,3,3,2,2,1)     | 9                      |
| 21        | 7349                    | (1,2,3,3,3,3,3,2,1)                          | 9                      |
| 22        | 10715                   | (1,2,2,3,3,3,3,2,2,1)                        | 10                     |
| 23        | 15370                   | (1,2,2,3,3,3,3,3,2,1), (1,2,3,3,3,3,3,2,2,1) | 10                     |
| 24        | 22050                   | (1,2,3,3,3,3,3,3,2,1)                        | 10                     |
| 25        | 32139                   | (1,2,2,3,3,3,3,3,2,2,1)                      | 11                     |

# From Trunks to bounded Trunks

## Definition: Width, Weight, Length

Let  $E$  be a trunk graph with  $N \geq 1$  edges and  $n$  vertices.

- The *width* of  $E$  is  $w(E) := \max_{v \in E^0} |s_E^{-1}(v)| = \max_{1 \leq k \leq n-1} m_k$ ,
- The *weight* of  $E$  is  $N$ ,
- The *length* of  $E$  is  $n - 1$ .

We denote the set of all trunk graphs with weight  $N$  and width  $w$  by  $\mathcal{T}_w^N$ .

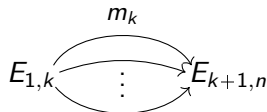
## Theorem:

The following equality holds:

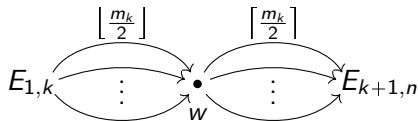
$$\sup_{E \in \mathcal{T}^N} P(E) = \sup_{E \in \mathcal{T}_3^N} P(E).$$

# Proof

Let us consider graph  $E \in \mathcal{T}^N \setminus \mathcal{T}_3^N$  and let  $v_k$  be a vertex satisfying  $m_k > 3$ .



and compare it with a graph  $G$  given by decomposition:



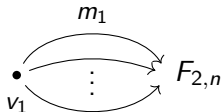
$$\begin{aligned}
 P(G) - P(E) &= 1 + P_t(v_k) \cdot \left\lfloor \frac{m_k}{2} \right\rfloor + \left\lceil \frac{m_k}{2} \right\rceil \cdot P_s(v_{k+1}) \\
 &\quad + P_t(v_k) \cdot \left\lfloor \frac{m_k}{2} \right\rfloor \cdot \left\lceil \frac{m_k}{2} \right\rceil \cdot P_s(v_{k+1}) - P_t(v_k) \cdot m_k \cdot P_s(v_{k+1}).
 \end{aligned}$$

# Properties

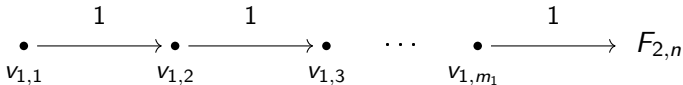
Lemma:

Let  $F \in \mathcal{T}_3^N$  be the optimal graph then:

$$m_1 = m_{n-1} = 1.$$



and compare it with graph  $G$ :

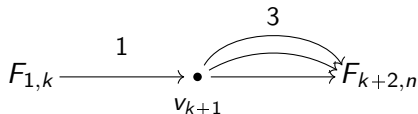


# Properties

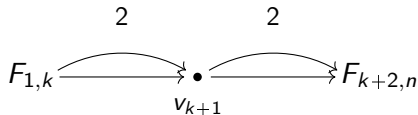
Lemma:

Let  $F \in \mathcal{T}_3^N$  be the optimal graph then:

$$\forall 1 \leq k \leq n-2 \quad |m_k - m_{k+1}| \leq 1.$$



and compare

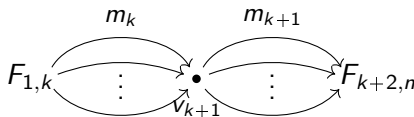


# Properties

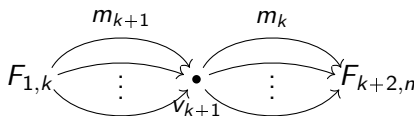
## Lemma:

Let  $F \in \mathcal{T}_3^N$  be the optimal graph and  $k^* = \min\{k: P_t(v_k) \geq P_s(v_{k+2})\}$ , then:

- $\forall 1 \leq k < k^* \quad m_k \leq m_{k+1},$
- $\forall k^* \leq k \leq n-2 \quad m_{k+1} \leq m_k.$



and compare it with a graph  $G$  given by decomposition:





# Properties

## Lemma: Summary

Let  $F \in \mathcal{T}_3^N$  be the optimal graph and  $k^* = \min\{k: P_t(v_k) \geq P_s(v_{k+2})\}$ , then:

- $m_1 = m_{n-1} = 1$ ,
- $\forall_{1 \leq k \leq n-2} |m_k - m_{k+1}| \leq 1$ ,
- $\forall_{1 \leq k < k^*} m_k \leq m_{k+1}$ ,
- $\forall_{k^* \leq k \leq n-2} m_{k+1} \leq m_k$ .

# Sequence

## Definition: Associated Sequence

All lemmas that we have proven thus far yield a very particular form of the associated sequence of edges  $I := (m_k)_{k=1}^{n-1}$  in an optimal graph. Namely, we know that  $I$  is a sequence of the form:

$$I = (1, \dots, 1, 2, \dots, 2, 3, \dots, 3, 2, \dots, 2, 1, \dots, 1).$$

Say that there are  $n_1$  copies of 1's at the beginning, followed by  $n_2$  2's,  $n_3$  3's,  $n_4$  2's and finally  $n_5$  1's at the end.

We will use the following shorthand for the sequence  $I$ :

$$I =: (n_1, n_2, n_3, n_4, n_5).$$

## Notation:

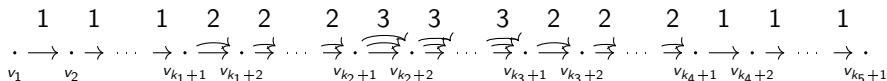
Note that when  $n_3 = 0$  then  $n_2, n_4$  are not well defined, similarly when  $n_2 = n_4 = 0$  then  $n_1, n_5$  are not well defined.

# Trunks

## Notation: "Critical Index"

For  $j \in \{1, 2, 3, 4, 5\}$  let  $k_j$  denote the  $j$ -th index such that  $m_j \neq m_{j+1}$ .

For critical indices the following equality holds:  $k_j = \sum_{i=1}^j n_i$ .



# Formula for paths

## Formula:

If  $E$  has associated sequence  $(n_1, n_2, n_3, n_4, n_5)$  then

$$P_t(E) = \left( \left( (n_1 + 2) 2^{n_2} - \frac{1}{2} \right) 3^{n_3} + \frac{1}{2} \right) 2^{n_4} + n_5 - 1 \quad (1)$$

$$\begin{aligned} P(E) = & [(n_1 + 2) 2^{n_2} (n_5 + 2) 2^{n_4}] 3^{n_3} - \frac{1}{2} [(n_1 + 2) 2^{n_2} + (n_5 + 2) 2^{n_4}] 3^{n_3} \\ & + \frac{1}{4} 3^{n_3} + \frac{1}{2} [(n_1 + 2) 2^{n_2} + (n_5 + 2) 2^{n_4}] \\ & + \frac{1}{2} [n_1 (n_1 - 1) + n_5 (n_5 - 1)] - [n_2 + n_4] - \frac{1}{2} n_3 - \frac{13}{4} \end{aligned} \quad (2)$$

# Sketch

## Proof

Let  $E$  have associated sequence  $(n_j)_1^5$  and  $\{k_j\}_1^5$  be the critical indices, then

$$\begin{aligned}
 P_t(E) &= \sum_{k=1}^n P(v_k, v_n) \\
 &= \sum_{k=k_5+1}^n P(v_k, v_n) + \sum_{k=k_4+1}^{k_5} P(v_k, v_n) + \sum_{k=k_3+1}^{k_4} P(v_k, v_n) \\
 &\quad + \sum_{k=k_2+1}^{k_3} P(v_k, v_n) + \sum_{k=k_1+1}^{k_2} P(v_k, v_n) + \sum_{k=1}^{k_1} P(v_k, v_n) \\
 &= 1 + \sum_{k=1}^{n_5} 1^k + \sum_{k=1}^{n_4} 2^k \\
 &\quad + \sum_{k=1}^{n_3} 3^k \cdot 2^{n_4} + \sum_{k=1}^{n_2} 2^k \cdot 3^{n_3} \cdot 2^{n_4} + \sum_{k=1}^{n_1} 1^k \cdot 2^{n_2} \cdot 3^{n_3} \cdot 2^{n_4}
 \end{aligned}$$

# Sketch

## Proof

$$\begin{aligned}
 P(F) &= \sum_{k=1}^n \sum_{l=1}^k P(v_l, v_k) = \sum_{k=1}^n P_t(v_k) = \sum_{k=1}^n P_t(E_{1,k}) \\
 &= P_t(E_{1,1}) + \sum_{k=2}^{k_1+1} P_t(E_{1,k}) + \sum_{k=k_1+2}^{k_2+1} P_t(E_{1,k}) \\
 &\quad + \sum_{k=k_2+2}^{k_3+1} P_t(E_{1,k}) \\
 &\quad + \sum_{k=k_3+2}^{k_4+1} P_t(E_{1,k}) \\
 &\quad + \sum_{k=k_4+2}^{k_5+1} P_t(E_{1,k})
 \end{aligned}$$

# Sketch

## Proof

Applying formula (1) we get:

$$\begin{aligned}
 P(E) = & 1 + \sum_{k=1}^{n_1} [k + 1] + \sum_{k=1}^{n_2} \left[ (n_1 + 2) 2^k - 1 \right] \\
 & + \sum_{k=1}^{n_3} \left[ \left( (n_1 + 2) 2^{n_2} - \frac{1}{2} \right) 3^k - \frac{1}{2} \right] \\
 & + \sum_{k=1}^{n_4} \left[ \left( \left( (n_1 + 2) 2^{n_2} - \frac{1}{2} \right) 3^{n_3} + \frac{1}{2} \right) 2^k - 1 \right] \\
 & + \sum_{k=1}^{n_5} \left[ \left( \left( (n_1 + 2) 2^{n_2} - \frac{1}{2} \right) 3^{n_3} + \frac{1}{2} \right) 2^{n_4} + k - 1 \right]
 \end{aligned}$$

# Sequences and their properties

## Lemma:

If  $F$  is the optimal graph then its sequence  $(n_1, n_2, n_3, n_4, n_5)$  satisfies:

- $n_2 \geq 1$  whenever  $N \geq 5$ ,
- $n_1 = n_5 = 1$  whenever  $N \geq 8$ ,
- $n_3 \geq 1$  whenever  $N \geq 13$ ,
- $n_3 \geq 2$  whenever  $N \geq 16$ .

Moreover

- $n_1, n_5 \leq 3$  whenever  $n_2 \geq 1$ ,
- $|n_2 - n_4| \leq 1$  whenever  $n_3 \geq 1$ ,
- $\max\{n_2, n_4\} \leq 2$  whenever  $n_3 \geq 2$ .



# Moves

## List

- ①  $(n_1, n_2, n_3, n_4, n_5) \mapsto (n_1 - 2, n_2 + 1, n_3, n_4, n_5)$
- ②  $(n_1, n_2, n_3, n_4, n_5) \mapsto (n_1 - 1, n_2 - 1, n_3 + 1, n_4, n_5)$
- ③  $(n_1, n_2, n_3, n_4, n_5) \mapsto (n_1 - 1, n_2, n_3 + 1, n_4 - 1, n_5)$
- ④  $(n_1, n_2, n_3, n_4, n_5) \mapsto (n_1 - 1, n_2 + 1, n_3 - 1, n_4 + 1, n_5)$
- ⑤  $(n_1, n_2, n_3, n_4, n_5) \mapsto (n_1, n_2 - 2, n_3 + 2, n_4 - 1, n_5)$
- ⑥  $(n_1, n_2, n_3, n_4, n_5) \mapsto (n_1, n_2 - 1, n_3, n_4 + 1, n_5)$

# Moves

## Move 1

$$(n_1, n_2, n_3, n_4, n_5) \mapsto (n_1 - 2, n_2 + 1, n_3, n_4, n_5)$$

$$E_{1, n_1-1} \xrightarrow{1} \underset{v_{n_1}}{\bullet} \xrightarrow{1} E_{k_1+1, n}$$

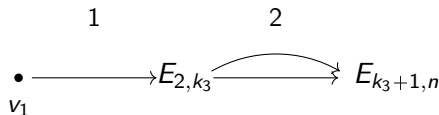
And let's compare it with graph  $G$

$$E_{1, n_1-1} \xrightarrow{2} E_{k_1+1, n}$$

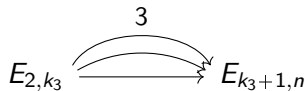
# Moves

## Move 2

$$(n_1, n_2, n_3, n_4, n_5) \mapsto (n_1 - 1, n_2 - 1, n_3 + 1, n_4, n_5)$$



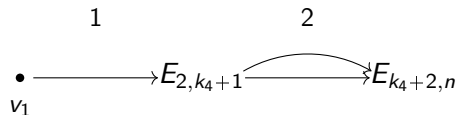
Compare it with



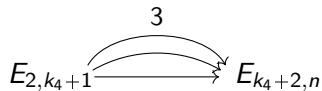
# Moves

## Move 3

$$(n_1, n_2, n_3, n_4, n_5) \mapsto (n_1 - 1, n_2, n_3 + 1, n_4 - 1, n_5)$$



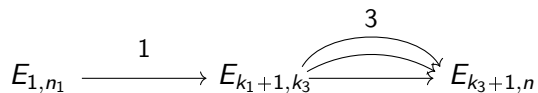
Compare it with



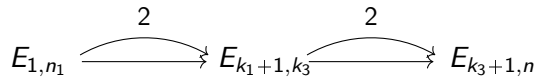
# Moves

## Move 4

$$(n_1, n_2, n_3, n_4, n_5) \mapsto (n_1 - 1, n_2 + 1, n_3 - 1, n_4 + 1, n_5)$$



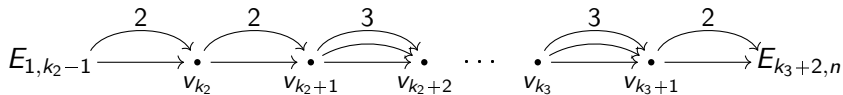
Compare it with



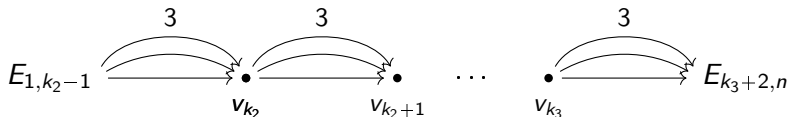
# Moves

## Move 5

$$(n_1, n_2, n_3, n_4, n_5) \mapsto (n_1, n_2 - 2, n_3 + 2, n_4 - 1, n_5)$$



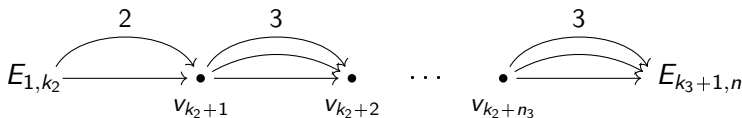
Let us consider the graph  $G$



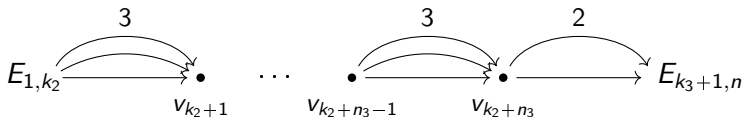
# Moves

## Move 6

$$(n_1, n_2, n_3, n_4, n_5) \mapsto (n_1, n_2 - 1, n_3, n_4 + 1, n_5)$$



and compare it with graph  $G$



## Three Sisters

### Theorem: Form of Optimal Graph

If  $N \geq 16$  and  $F$  maximizes the dimension of  $\mathbb{K}E$  among  $\mathcal{E}^N$ , then

- $F = (1, 1, \frac{N-6}{3}, 1, 1) = (1, 1, \lfloor \frac{N}{3} \rfloor - 2, 1, 1),$
- $F = (1, 2, \frac{N-10}{3}, 2, 1) = (1, 2, \lfloor \frac{N}{3} \rfloor - 3, 2, 1),$
- $F = (1, 2, \frac{N-8}{3}, 1, 1) = (1, 2, \lfloor \frac{N}{3} \rfloor - 2, 1, 1),$

depending on the remainder of division of  $N$  by 3.

### Theorem: Dimension of the Path Algebra

If  $N \geq 16$  and  $F$  maximizes the dimension of  $\mathbb{K}E$  among  $\mathcal{E}^N$ , then

- $P(F) = \frac{1}{4} \left( 121 \cdot 3^{\lfloor \frac{N}{3} \rfloor - 2} - 2 \lfloor \frac{N}{3} \rfloor + 7 \right),$
- $P(F) = \frac{1}{4} \left( 529 \cdot 3^{\lfloor \frac{N}{3} \rfloor - 3} - 2 \lfloor \frac{N}{3} \rfloor + 25 \right),$
- $P(F) = \frac{1}{4} \left( 253 \cdot 3^{\lfloor \frac{N}{3} \rfloor - 2} - 2 \lfloor \frac{N}{3} \rfloor + 15 \right),$

depending on the remainder of division of  $N$  by 3.



# Leavitt Path Algebra

## Leavitt Path Algebra Methods

# From Trees to Trunks

## Theorem

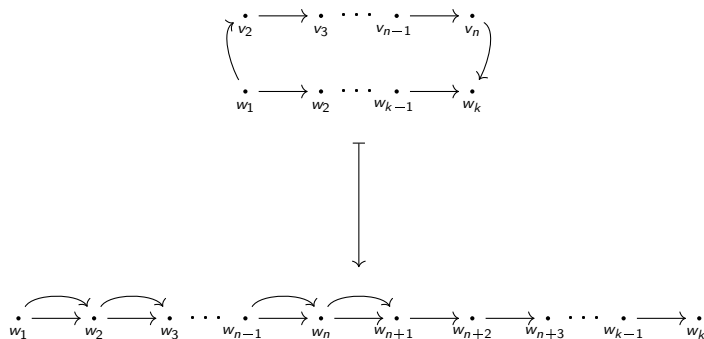
The following equality holds:

$$\sup_{E \in \mathcal{O}_S^N} P_t(E) = \sup_{E \in \mathcal{T}^N} P_t(E).$$

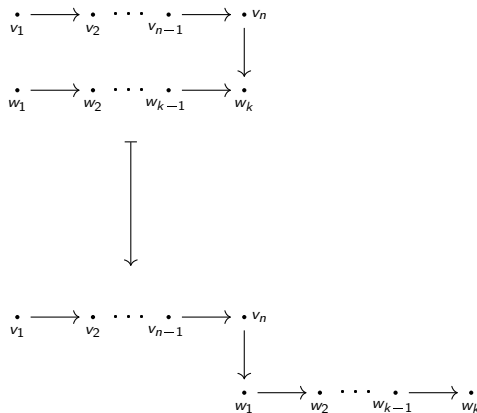
## Moves

- Resolving alternative paths,
- Combing stray paths.

# Moves - Resolving alternative paths



# Moves - Combing stray paths from sources



# Python

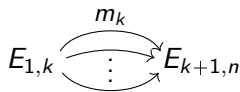
| $N$ edges | Maximal number of paths | Optimal trunks        | Length of longest path |
|-----------|-------------------------|-----------------------|------------------------|
| 3         | 5                       | (1,2)                 | 2                      |
| 4         | 7                       | (1,1,2),(1,3),(2,2)   | 2 or 3                 |
| 5         | 11                      | (1,2,2)               | 3                      |
| 6         | 16                      | (1,2,3)               | 3                      |
| 7         | 23                      | (1,2,2,2)             | 4                      |
| 8         | 34                      | (1,2,2,3)             | 4                      |
| 9         | 49                      | (1,2,3,3)             | 4                      |
| 10        | 70                      | (1,2,2,2,3)           | 5                      |
| 11        | 103                     | (1,2,2,3,3)           | 5                      |
| 12        | 148                     | (1,2,3,3,3)           | 5                      |
| 13        | 211                     | (1,2,2,2,3,3)         | 6                      |
| 14        | 310                     | (1,2,2,3,3,3)         | 6                      |
| 15        | 445                     | (1,2,3,3,3,3)         | 6                      |
| 16        | 634                     | (1,2,2,2,3,3,3)       | 7                      |
| 17        | 931                     | (1,2,2,3,3,3,3)       | 7                      |
| 18        | 1336                    | (1,2,3,3,3,3,3)       | 7                      |
| 19        | 1903                    | (1,2,2,2,3,3,3,3)     | 8                      |
| 20        | 2794                    | (1,2,2,3,3,3,3,3)     | 8                      |
| 21        | 4009                    | (1,2,3,3,3,3,3,3)     | 8                      |
| 22        | 5710                    | (1,2,2,2,3,3,3,3,3)   | 9                      |
| 23        | 8383                    | (1,2,2,3,3,3,3,3,3)   | 9                      |
| 24        | 12028                   | (1,2,3,3,3,3,3,3,3)   | 9                      |
| 25        | 17131                   | (1,2,2,2,3,3,3,3,3,3) | 10                     |

# Ones and differences

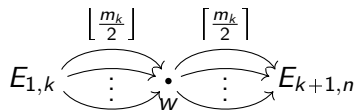
## Lemma

The following equality holds:

$$\sup_{E \in \mathcal{T}^N} P_t(E) = \sup_{E \in \mathcal{T}_3^N} P_t(E).$$



and compare it with a graph  $G$  given by decomposition:



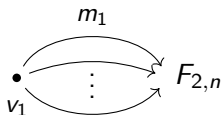
$$\begin{aligned} P_t(G) - P_t(E) &= \left\lfloor \frac{m_k}{2} \right\rfloor \cdot P_s(v_{k+1}) + P_t(v_k) \cdot \left\lfloor \frac{m_k}{2} \right\rfloor \cdot \left\lfloor \frac{m_k}{2} \right\rfloor \cdot P_s(v_{k+1}) \\ &\quad - P_t(v_k) \cdot m_k \cdot P_s(v_{k+1}). \end{aligned}$$

# Properties

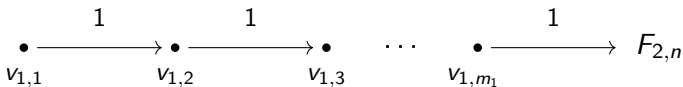
## Lemma:

Let  $F \in \mathcal{T}_3^N$  be the optimal graph then:

$$m_1 = 1.$$



and compare it with graph  $G$ :

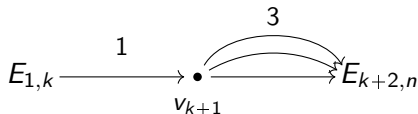


# Properties

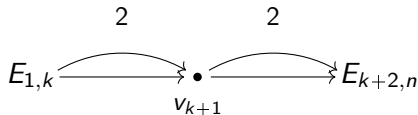
Lemma:

Let  $F \in \mathcal{T}_3^N$  be the optimal graph then:

$$\forall 1 \leq k \leq n-2 \quad |m_k - m_{k+1}| \leq 1.$$



and compare



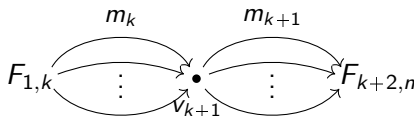


# Properties

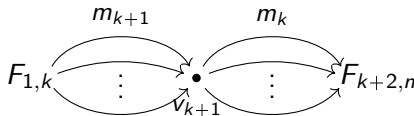
Lemma:

Let  $F \in \mathcal{T}_3^N$  be the optimal graph then:

$$\forall 1 \leq k \leq n-2 \quad m_k \leq m_{k+1}.$$



and compare it with a graph  $G$  given by decomposition:



# Properties

## Lemma: Summary

Let  $F \in \mathcal{T}_3^N$  be the optimal graph then:

- $m_1 = 1$ ,
- $\forall_{1 \leq k \leq n-2} |m_k - m_{k+1}| \leq 1$ ,
- $\forall_{1 \leq k \leq n-2} m_k \leq m_{k+1}$ .

# Paths to the sink

## Formula

If  $E$  has associated sequence  $(n_1, n_2, n_3)$  then

$$P_t(E) = \left( (n_1 + 2) 2^{n_2} - \frac{1}{2} \right) 3^{n_3} - \frac{1}{2}$$

## Proof

We apply formula (1)

# Sequences and their properties

## Lemma:

If  $F$  is the optimal graph then its sequence  $(n_1, n_2, n_3)$  satisfies:

- $n_2 \geq 1$  whenever  $N \geq 3$ ,
- $n_1 = 1$  whenever  $N \geq 5$ ,
- $n_3 \geq 1$  whenever  $N \geq 8$ .

Moreover

- $n_2 \leq 3$  whenever  $n_3 \geq 1$ .

# Moves

## Moves

Let us consider a trunk with associated sequence  $(n_1, n_2, n_3)$ .

- ①  $(n_1, n_2, n_3) \mapsto (n_1 - 2, n_2 + 1, n_3)$
- ②  $(n_1, n_2, n_3) \mapsto (n_1, n_2 - 3, n_3 + 2)$
- ③  $(n_1, n_2, n_3) \mapsto (n_1 - 1, n_2 - 1, n_3 + 1)$

# Moves

## Move 1

$$(n_1, n_2, n_3) \mapsto (n_1 - 2, n_2 + 1, n_3)$$

$$E_{1, n_1 - 1} \xrightarrow{1} \underset{v_{n_1}}{\bullet} \xrightarrow{1} E_{k_1 + 1, n}$$

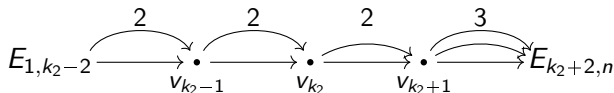
And let's compare it with graph  $G$

$$E_{1, n_1 - 1} \xrightarrow{2} E_{k_1 + 1, n}$$

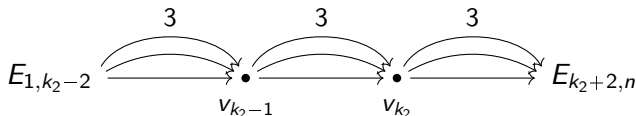
# Moves

## Move 2

$$(n_1, n_2, n_3) \mapsto (n_1, n_2 - 3, n_3 + 2)$$



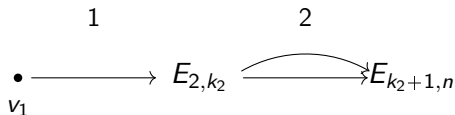
Let us consider the graph  $G$



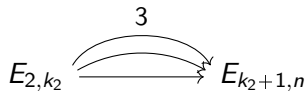
# Moves

## Move 3

$$(n_1, n_2, n_3) \mapsto (n_1 - 1, n_2 - 1, n_3 + 1)$$



Compare it with





# Optimal graphs and Three sisters

## Theorem:

If  $N \geq 8$  and  $F$  maximizes the dimension of  $L_{\mathbb{K}}(E)$  among  $\mathcal{E}_s^N$  then

- $F = (1, 1, \frac{N-3}{3}) = (1, 1, \lfloor \frac{N}{3} \rfloor - 1)$ ,  $P_t(F) = \frac{1}{2} \left( 11 \cdot 3^{\lfloor \frac{N}{3} \rfloor - 1} - 1 \right)$
- $F = (1, 3, \frac{N-7}{3}) = (1, 3, \lfloor \frac{N}{3} \rfloor - 2)$ ,  $P_t(F) = \frac{1}{2} \left( 47 \cdot 3^{\lfloor \frac{N}{3} \rfloor - 2} - 1 \right)$
- $F = (1, 2, \frac{N-5}{3}) = (1, 2, \lfloor \frac{N}{3} \rfloor - 1)$ ,  $P_t(F) = \frac{1}{2} \left( 23 \cdot 3^{\lfloor \frac{N}{3} \rfloor - 1} - 1 \right)$

depending on the value of  $N \bmod 3$ .

## Open Problem

A new problem appears!

# Gap Problem

## Necessary and Sufficient numbers

For any natural number  $n \in \mathbb{N}$  let us consider numbers:

$$m(n) = \max \left\{ N : \sup_{E \in \mathcal{E}_s^N} P_t(E) < n \right\} + 1$$

$$M(n) = \min \{ N : M_n(\mathbb{K}) = L_{\mathbb{K}}(E), |E^1| = N \}$$

called necessary and sufficient numbers for construction respectively.

## Gap problem

What is the behavior of  $M(n)$  compared to  $m(n)$ ?

In particular:

$$\sup_{n \in \mathbb{N}} (M(n) - m(n)) < +\infty.$$

# The End

Thank you!!!