

Quantitative Stability in Optimal Transport: From Classical Potentials to Stable Transport Maps

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Entry points into optimal transport

There are already excellent introductions to the standard theory.

- ▶ C. Villani, *Topics in Optimal Transportation; Optimal Transport: Old and New*.
- ▶ F. Santambrogio, *Optimal Transport for Applied Mathematicians*.
- ▶ L. Ambrosio, N. Gigli, G. Savaré, *Gradient Flows*.
- ▶ A. Figalli, F. Glaudo, *An Invitation to Optimal Transport, Wasserstein Distances, and Gradient Flows*.
- ▶ L. Ambrosio, E. Brué, D. Semola, *Lectures on Optimal Transport*.
- ▶ G. Peyré, M. Cuturi, *Computational Optimal Transport*.

These lectures do not try to replace these references.

Course philosophy

The standard route often begins with the (Monge-)Kantorovich primal problem:

transport plans $\longrightarrow$ duality $\longrightarrow$ Kantorovich potentials.

We will try to reverse it.

What is the potential behind optimal transport?

Lectures 01–02 roadmap

In the first two lectures we build the dual viewpoint step by step.

1. **Dual** potentials as prices with a no-arbitrage condition.
2. Hard c -**transforms**.
3. Softmin **entropic regularization** and zero-temperature limit
4. **Wasserstein distance** of order p .
5. Cyclical monotonicity and **Rockafellar theorem**.

Economic starting point: replacement costs

Suppose that a unit of mass at $x \in X$ can be replaced by a unit of mass at $y \in Y$.

$v(y)$ = price of using target point y ,

$c(x, y)$ = cost of replacing x by y .

Then the total replacement cost of x through y is

$$v(y) + c(x, y).$$

Admissible valuation No arbitrage

A source value $u(x)$ is admissible only if it is no larger than every possible replacement cost.

The admissibility condition is therefore

$$u(x) \leq v(y) + c(x, y) \quad \forall x \in X, y \in Y.$$

Equivalently,

$$u(x) - v(y) \leq c(x, y).$$

If this failed for some pair (x, y) , then

$$u(x) > v(y) + c(x, y),$$

so x would be valued more than the cost of replacing it by y .

Arbitrage-free valuations

$$u(x) \leq v(y) + c(x, y) \quad \forall x, y.$$

$u(x)$ = value assigned to source point x ,

$v(y)$ = price assigned to target point y ,

$v(y) + c(x, y)$ = cost of replacing x by y .

Kantorovich potentials are arbitrage-free valuations.

Dual gain

Let $\rho \in \mathcal{P}(X)$, $\mu \in \mathcal{P}(Y)$,

ρ = distribution of items to be sold / valued,

μ = distribution of items to be bought / used as replacements.

For admissible prices

$$u(x) \leq v(y) + c(x, y),$$

the dual gain is

$$\int_X u \, d\rho - \int_Y v \, d\mu.$$

The dual problem

The Kantorovich (dual) cost is

$$W_c(\rho, \mu) := \sup_{u(x) \leq v(y) + c(x, y)} \left\{ \int_X u \, d\rho - \int_Y v \, d\mu \right\}.$$

Maximize value discrepancy (gain) under the no-arbitrage constraint.

Sign convention

Standard notation:

$$\varphi(x) + \psi(y) \leq c(x, y).$$

Our notation:

$$u(x) - v(y) \leq c(x, y), \quad \text{equivalently} \quad u(x) \leq v(y) + c(x, y).$$

$$u = \varphi, \quad v = -\psi.$$

This convention is adapted to the c -transform operator

$$Qv(x) = \inf_y \{v(y) + c(x, y)\},$$

the cheapest replacement value induced by v .

The primal problem Monge - Kantorovich

Couplings:

$$\Pi(\underline{\rho}, \underline{\mu}) = \{\pi \in \mathcal{P}(X \times Y) : (p_X)_\# \pi = \underline{\rho}, (p_Y)_\# \pi = \underline{\mu}\}.$$

Primal problem:

$$W_c(\underline{\rho}, \underline{\mu}) \stackrel{?}{=} \inf_{\pi \in \Pi(\underline{\rho}, \underline{\mu})} \int_{X \times Y} c(x, y) d\pi(x, y).$$

primal = feasible matchings,
dual = feasible prices.

Weak duality

$$\begin{aligned}
 W_c(s, \mu) &= \sup_{u \in \mathcal{U} \leq c} \int_X u d\gamma - \int_Y v d\mu \\
 \forall \pi \in \Pi(s, \mu) \quad & \left| \begin{aligned}
 \int_X u d\gamma &= \int_{X \times Y} u(x) d\pi(x, y) & \int_Y v d\mu &= \int_{X \times Y} v(y) d\pi(x, y) \\
 &= \sup_{u \in \mathcal{U} \leq c} \int_{X \times Y} \underbrace{(u(x) - v(y))}_{\leq c(x, y)} d\pi(x, y) \\
 &\leq \sup_{u \in \mathcal{U} \leq c} \int_{X \times Y} c(x, y) d\pi(x, y) \leftarrow \inf_{\pi}
 \end{aligned} \right.
 \end{aligned}$$

Theorem: Kantorovich duality

Let X, Y be Polish spaces, let

$$\rho \in \mathcal{P}(X), \quad \mu \in \mathcal{P}(Y),$$

and let

$$c : X \times Y \rightarrow (-\infty, +\infty]$$

be lower semicontinuous and bounded from below. Then,

$$\rightarrow \inf_{\pi \in \Pi(\rho, \mu)} \int_{X \times Y} c(x, y) d\pi(x, y) = \sup_{u(x) \leq v(y) + c(x, y)} \left\{ \int_X u d\rho - \int_Y v d\mu \right\},$$

where the supremum is taken over integrable admissible potentials.

Moreover, if the primal value is finite, an optimal plan exists.

References: Villani, *Topics in Optimal Transportation*, Thm. 1.3;

Ambrosio–Gigli–Savaré, *Gradient Flows*, Ch. 6;

Santambrogio, *Optimal Transport for Applied Mathematicians*, Ch. 1.

Formal derivation: introduce multipliers

Primal:

$$\inf_{\pi \geq 0} \left\{ \int c \, d\pi : (\text{pr}_X)_\# \pi = \rho, (\text{pr}_Y)_\# \pi = \mu \right\}.$$

$$\pi(A \times Y) = \rho(A) \quad \forall A$$

↑

$$\int_{X \times Y} u(x,y) \, d\pi(x,y) = \int_X u(x) \, d\rho(x)$$

∀ u

Lagrange multipliers for the two marginal constraints:

$$u : X \rightarrow \mathbb{R}, \quad v : Y \rightarrow \mathbb{R}.$$

$$\mathcal{L}(\pi, u, v) = \int c \, d\pi + \underbrace{\int u \, d\rho - \int u(x) \, d\pi}_{=0} - \underbrace{\int v \, d\mu + \int v(y) \, d\pi}_{=0}.$$

Formal derivation

$$\mathcal{L}(\pi; (u, v)) = \int c d\pi - \int u(x) d\pi + \int u(x) g(x) + \int v(y) d\pi - \int v(y) d\mu(y)$$

$$\sup_{(u, v)} \mathcal{L}(\pi; (u, v)) = \int c d\pi + \sup_u \left(\underbrace{- \int u(x) d\pi + \int u d\mathcal{G}}_{+\infty \text{ unless } (P_x)_\# \pi = \mathcal{G}} \right) + \sup_v \left(\underbrace{- \int v(y) d\pi + \int v d\mu}_{+\infty \text{ unless } (P_y)_\# \pi = \mu} \right)$$

$$\inf_{\pi} \sup_{(u, v)} \mathcal{L}(\pi; (u, v)) = \inf_{\pi \in \Pi(\mathcal{G}, \mu)} \int c d\pi$$

$$\boxed{\sup_{(u, v)} \inf_{\pi} \mathcal{L}(\pi; (u, v)) = \inf_{\substack{\pi \geq 0 \\ \pi \in \Pi(\mathcal{G}, \mu)}} \int (c - u(x) + v(y)) d\pi + \left| \int u d\mathcal{G} - \int v d\mu \right|}$$

$\downarrow$
 $x \times y \rightarrow \text{unless } \boxed{u(x) \leq v(y) + c} \rightarrow 0$

Formal duality

The min-max principle (formally) gives

$$\begin{aligned}
 \inf_{\pi \in \Pi(\rho, \mu)} \int c \, d\pi &= \inf_{\pi \geq 0} \sup_{u, v} \mathcal{L}(\pi, u, v) \\
 &= \sup_{u, v} \inf_{\pi \geq 0} \mathcal{L}(\pi, u, v) \\
 &= \sup_{u(x) \leq v(y) + c(x, y)} \left\{ \int u \, d\rho - \int v \, d\mu \right\}.
 \end{aligned}$$

The (hard) c -transform

Given v , admissibility says

$$u(x) \leq v(y) + c(x, y) \quad \forall y \in Y.$$

Therefore

$$u(x) \leq \inf_{y \in Y} \{v(y) + c(x, y)\}.$$

Define the c -transform of v

$$Q_c v(x) := \inf_{y \in Y} \{v(y) + c(x, y)\}.$$

write $Q = Q_c$ if c is clear. Standard notation ($v = -\psi$)

$$Q_c v \longleftrightarrow \psi^c.$$

Semi-dual formulation

The two-potential problem can be reduced one potential replacing u with Qv .

Kantorovich functional:

$$\mathcal{D}_\mu(v) := \int_X Qv \, d\rho - \int_Y v \, d\mu.$$

We make the problem less symmetric but reduce to a single potential v .

Eliminating either potential

Given v , the best admissible u is

$$u = Qv, \quad Qv(x) := \inf_{y \in Y} \{v(y) + c(x, y)\}.$$

Given u , admissibility

$$u(x) \leq v(y) + c(x, y)$$

is equivalent to

$$v(y) \geq u(x) - c(x, y) \quad \forall x.$$

So the smallest admissible v is

$$\hat{Q}u(y) := \sup_{x \in X} \{u(x) - c(x, y)\}.$$

We can eliminate either variable: $u = Qv$ or $v = \hat{Q}u$.

The two transforms encode the same constraint

$$Qv(x) := \inf_{y \in Y} \{v(y) + c(x, y)\}, \quad \hat{Q}u(y) := \sup_{x \in X} \{u(x) - c(x, y)\}.$$

The following are equivalent:

$$u(x) \leq v(y) + c(x, y) \quad \forall x, y,$$

$$u \leq Qv,$$

$$\hat{Q}u \leq v.$$

General properties of Q and $\hat{Q}$

$$v \leq w \implies Qv \leq Qw, \quad u \leq z \implies \hat{Q}u \leq \hat{Q}z.$$

$$Q(v + a) = Qv + a, \quad \hat{Q}(u + a) = \hat{Q}u + a. \quad a \in \mathbb{R}$$

$$\|Qv - Qw\|_{\infty} \leq \|v - w\|_{\infty}, \quad \|\hat{Q}u - \hat{Q}z\|_{\infty} \leq \|u - z\|_{\infty}.$$

$$u \leq Q\hat{Q}u, \quad \hat{Q}Qv \leq v.$$

$$Q\hat{Q}Q = Q, \quad \hat{Q}Q\hat{Q} = \hat{Q}.$$

Soft c -transform / Entropic c -transform

The hard transform is

$$Q_c v(x) = \inf_{y \in Y} \{v(y) + c(x, y)\}.$$

Its soft version is for $\varepsilon > 0$ (temperature):

$$Q_c^\varepsilon v(x) := -\varepsilon \log \int_Y \exp \left(-\frac{v(y) + c(x, y)}{\varepsilon} \right) d\mu(y).$$

Interpretation: soft replacement value of x . Notice that μ explicitly appears!

Gibbs variational principle

For a function $F : Y \rightarrow \mathbb{R}$,

$$-\varepsilon \log \int_Y e^{-F/\varepsilon} d\mu = \inf_{\gamma \in \mathcal{P}(Y)} \left\{ \int_Y F d\gamma + \varepsilon \text{Ent}(\gamma \mid \mu) \right\}.$$

Apply this with

$$F_x(y) = v(y) + c(x, y).$$

Then

$$Q_c^\varepsilon v(x) = \inf_{\gamma \in \mathcal{P}(Y)} \left\{ \int_Y (v(y) + c(x, y)) d\gamma(y) + \varepsilon \text{Ent}(\gamma \mid \mu) \right\}.$$

As $\varepsilon \downarrow 0$,

$$\boxed{Q_c^\varepsilon v(x) \longrightarrow Q_c v(x).}$$

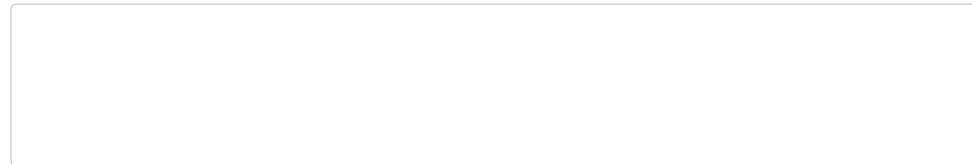
Gibbs kernel

The minimizer in the Gibbs variational principle is ~~Gibbs distribution~~ $\frac{e^{-f/\epsilon}}{Z}$

$$d\gamma_x^\epsilon(y) = \frac{\exp(-(v(y) + c(x, y))/\epsilon)}{\int_Y \exp(-(v(z) + c(x, z))/\epsilon) d\mu(z)} d\mu(y).$$

Equivalently,

$$d\gamma_x^\epsilon(y) = e^{\overbrace{(Q_\epsilon^\epsilon v(x) - v(y) - c(x, y))/\epsilon}} d\mu(y).$$



The soft adjoint transform

Hard adjoint transform:

$$\hat{Q}_c u(y) = \sup_{x \in X} \{u(x) - c(x, y)\}.$$

Its soft version is

$$\hat{Q}_c^\varepsilon u(y) := \varepsilon \log \int_X \exp \left(\frac{u(x) - c(x, y)}{\varepsilon} \right) d\rho(x).$$

For the soft transforms the exact identity is lost:

$$Q_c^\varepsilon \hat{Q}_c^\varepsilon Q_c^\varepsilon \neq Q_c^\varepsilon \quad \text{in general.}$$

Instead of one exact projection, one iterates the two soft transforms:
this is the Sinkhorn mechanism.

Maslov dequantization

The soft minimum is based on Laplace principle

$$-\varepsilon \log \int e^{-F/\varepsilon} \longrightarrow \inf F \quad \text{as } \varepsilon \downarrow 0.$$

This is often called the *Maslov dequantization*.

The usual operations are transformed as

$$+ \rightsquigarrow \min, \quad \times \rightsquigarrow +$$

Classical probability becomes idempotent, or tropical, probability.

(From the soft-max one can obtain a max, + idempotent probability)

From heat to Hopf–Lax

Let P_t^ε be the heat semigroup at diffusivity ε :

$$P_t^\varepsilon f(x) = \int_{\mathbb{R}^d} f(y) \frac{1}{(2\pi\varepsilon t)^{d/2}} e^{-|x-y|^2/(2\varepsilon t)} dy.$$

Apply it to

$$f = e^{-v/\varepsilon}.$$

Then

$$-\varepsilon \log P_t^\varepsilon(e^{-v/\varepsilon})(x) = -\varepsilon \log \int \exp\left(-\frac{v(y) + |x-y|^2/(2t)}{\varepsilon}\right) \frac{dy}{(2\pi\varepsilon t)^{d/2}}.$$

This is a soft version of the Hopf–Lax operator.

By the Laplace-(Varadhan) principle,

$$-\varepsilon \log P_t^\varepsilon(e^{-v/\varepsilon})(x) \longrightarrow \inf_y \left\{ v(y) + \frac{|x - y|^2}{2t} \right\} = Q_t v(x)$$

with $c(x, y) = |x - y|^2/2t$.

Hopf–Lax is the tropical shadow of the heat semigroup.

Why this is useful? we can guess many properties, e.g.

$$P_{t+s} = P_t P_s \Rightarrow Q_{t+s} = Q_t Q_s.$$

Wasserstein distance of order p

Let $(X = Y, d)$ be a (bounded) metric space and $p \geq 1$. For the cost

$$c(x, y) = d(x, y)^p,$$

the dual formulation gives

$$W_p^p(\rho, \mu) := \sup_{u(x) \leq v(y) + d(x, y)^p} \left\{ \int_X u d\rho - \int_X v d\mu \right\}$$

By Kantorovich duality, this is equal to the usual primal value

$$\inf_{\pi \in \Pi(\rho, \mu)} \int_{X \times X} d(x, y)^p d\pi(x, y).$$

$W_p(\rho, \mu)$

In this lecture, we use the dual formula as the definition and the primal formula as its representation.

A dual proof of the triangle inequality : $W_p(s, \mu) \leq W_p(s, \eta) + W_p(\eta, \mu)$

Fix $p \geq 1$. For $t > 0$, define the rescaled cost

$$c_t(x, y) := \frac{d(x, y)^p}{t^{p-1}}. \quad \leftarrow \boxed{p=2}$$

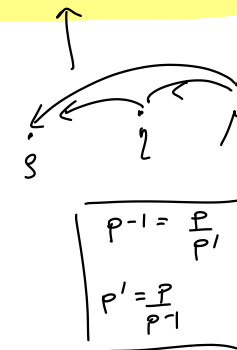
Correspondingly, define

$$Q_t v(x) := \inf_y \left\{ v(y) + \frac{d(x, y)^p}{t^{p-1}} \right\}.$$

Then

$$\boxed{W_{c_t}(\rho, \mu) = t^{1-p} W_p^p(\rho, \mu).}$$

The parameter t plays the role of time in the Hopf–Lax formula.



The key semigroup (inequality) $Q_{(s+t)} v \leq Q_t(Q_s v)$

$$Q_{s+t} v(x) = \inf_y \left\{ v(y) + \frac{d(x,y)^p}{(s+t)^{p-1}} \right\}$$

$$\frac{d(x,y)^p}{(t+s)^{p-1}} \leq \frac{d(x,z)^p}{t^{p-1}} + \frac{d(z,y)^p}{s^{p-1}} \leq \left(\frac{d(x,z)^p}{t^{p/p'}} + \frac{d(z,y)^p}{s^{p/p'}} \right)^{1/p'} \frac{(t+s)^{1/p'}}{(t+s)^{1/p'}}$$

$$\left| \frac{d(x,y)^p}{(t+s)^{p-1}} \leq \frac{d(x,z)^p}{t^{p-1}} + \frac{d(z,y)^p}{s^{p-1}} \right| \quad \forall s, t > 0$$

Lemma 2 $\frac{A^p}{(t+s)^{p-1}} \leq \frac{B^p}{t^{p-1}} + \frac{C^p}{s^{p-1}} \quad \forall s, t > 0 \Rightarrow A \leq B+C$

$$\begin{aligned}
Q_{s+t} v(x) &= \inf_y \left\{ v(y) + \frac{d(x,y)^p}{(s+t)^{p-1}} \right\} \\
&\leq \inf_z \inf_y \left\{ v(y) + \frac{d(x,z)^p}{s^{p-1}} + \frac{d(z,y)^p}{t^{p-1}} \right\} \\
&= \inf_z \left\{ \underbrace{\inf_y \left\{ v(y) + \frac{d(z,y)^p}{t^{p-1}} \right\}}_{Q_t v(z)} + \frac{d(x,z)^p}{s^{p-1}} \right\} \\
&= \inf_z \left\{ Q_t v(z) + \frac{d(x,z)^p}{s^{p-1}} \right\} = Q_s(Q_t v)(x)
\end{aligned}$$

Dual triangle inequality

$$W_{\frac{d}{t}P}(s, \mu) = \frac{1}{(s+t)^{p-1}} \underbrace{W_P(s, \mu)}_{A^p}$$

$$\begin{aligned} &= \sup_v \left\{ \int \underbrace{Q_{(s+t)} v}_{\substack{\uparrow \\ \text{or } s, t \geq 0}} d\mu - \int v d\mu \right\} \leq \sup_v \left\{ \int Q_t(Q_s v) d\mu - \int v d\mu \right\} \\ &= \sup_v \left[\int \underbrace{Q_t(Q_s v)}_{\substack{\uparrow \\ \text{or } s, t \geq 0}} d\mu - \int Q_s v d\mu \right] + \left[\int Q_s v d\mu - \int v d\mu \right] \\ &\leq W_{\frac{d}{t}P}(s, \eta) + \sup_v \left[\int Q_s v d\mu - \int v d\mu \right] \\ &= W_{\frac{d}{t}P}(s, \eta) + W_{\frac{d}{s}P}(\eta, \mu) = \frac{1}{t^{p-1}} \underbrace{W_P^p(s, \eta)}_{B^p} + \frac{1}{s^{p-1}} \underbrace{W_P^p(\eta, \mu)}_{C^p} \end{aligned}$$

Taking the supremum

Taking the supremum over v , we obtain

$$(s+t)^{1-p} W_p^p(\rho, \mu) \leq t^{1-p} W_p^p(\rho, \eta) + s^{1-p} W_p^p(\eta, \mu).$$

Optimize in s, t to conclude.

Remark $p=1$ $Q_\mu v(x) = \inf_y \{v(y) + d(x, y)\}$
 $Q_\mu v = v$ iff v is Lip $Lip(v) = 1$ $|v(x) - v(y)| \leq d(x, y)$
 $W_1(\rho, \mu) = \sup_{Lip(v) \leq 1} \left\{ \int v d\rho - \int v d\mu \right\}$

Exercises: quadratic cost and Gaussian examples

Let $c(x, y) = \frac{1}{2}|x - y|^2$ (or equivalently $c(x, y) = -x \cdot y$)

1. **Translations.** Let $\mu = (\tau_m)_\# \rho$, where

$$\tau_m(x) = x + m.$$

Show directly from the dual definition that τ_m is optimal.

2. **Dilations.** Let $T_a(x) = ax$, $a > 0$. Show that T_a is optimal from

$$\rho = \mathcal{N}(0, \sigma_0^2) \quad \text{to} \quad \mu = \mathcal{N}(0, \sigma_1^2), \quad a = \frac{\sigma_1}{\sigma_0}.$$

Is T_a optimal for arbitrary ρ ?

3. **Gaussian to Gaussian.** If

$$\rho = \mathcal{N}(m_0, \Sigma_0), \quad \mu = \mathcal{N}(m_1, \Sigma_1),$$

show that the optimal potentials are quadratic and the optimal map is affine.

Quadratic cost turns Gaussian OT into completing squares.

Hints for the quadratic exercises

Translation. Try affine potentials

$$u(x) = -m \cdot x, \quad v(y) = -m \cdot y + \frac{1}{2}|m|^2.$$

Check

$$\frac{1}{2}|x - y|^2 - (u(x) - v(y)) = \frac{1}{2}|y - x - m|^2.$$

Dilation in one dimension. Try

$$u(x) = \frac{1-a}{2}x^2, \quad v(y) = \frac{1-a}{2a}y^2.$$

Check

$$\frac{1}{2}(x - y)^2 - (u(x) - v(y)) = \frac{1}{2a}(y - ax)^2.$$

General Gaussians. Try affine $T(x) = m_1 + A(x - m_0)$ and quadratic potentials.

Contact set and c -subdifferentials

Recall

$$Qv(x) = \inf_y \{v(y) + c(x, y)\}, \quad \hat{Q}u(y) = \sup_x \{u(x) - c(x, y)\}.$$

Assume

$$u = Qv, \quad v = \hat{Q}u,$$

Define the subdifferentials

$$\partial_c u(x) := \{y \in Y : u(x) = v(y) + c(x, y)\},$$

$$\hat{\partial}_c v(y) := \{x \in X : u(x) = v(y) + c(x, y)\}.$$

$$\boxed{u(x) \leq v(y) + c(x, y)}$$

↑

Then the contact set is defined as

$$\Gamma(u, v) = \{(x, y) : y \in \partial_c u(x)\} = \{(x, y) : x \in \hat{\partial}_c v(y)\}.$$

Complementary slackness

Let $\pi \in \Pi(\rho, \mu)$ and let (u, v) be admissible:

$$c(x, y) - u(x) + v(y) \geq 0.$$

If π is primal optimal and (u, v) is dual optimal, then

$$\int_{X \times Y} (c(x, y) - u(x) + v(y)) d\pi(x, y) = 0.$$

Since the integrand is nonnegative,

$$c(x, y) - u(x) + v(y) = 0 \quad \pi\text{-a.e.}$$

Equivalently,

$$u(x) = v(y) + c(x, y) \quad \pi\text{-a.e.}$$

Optimal transport plans are concentrated on the contact set.

Definition: c -cyclical monotonicity

A set $\Gamma \subset X \times Y$ is c -cyclically monotone if for every finite family

$$(x_i, y_i) \in \Gamma, \quad i = 1, \dots, N,$$

with $y_{N+1} = y_1$,

$$\sum_{i=1}^N c(x_i, y_i) \leq \sum_{i=1}^N c(x_i, y_{i+1}).$$



Contact implies c -cyclical monotonicity

For $(x_i, y_i) \in \Gamma$, contact gives

$$u(x_i) - v(y_i) = c(x_i, y_i). \quad \leftarrow$$

Admissibility gives

$$u(x_i) - v(y_{i+1}) \leq c(x_i, y_{i+1}).$$

$$\begin{aligned} \sum_{i=1}^N c(x_i, y_i) &= \sum_{i=1}^N u(x_i) - \sum_{i=1}^N v(y_i) = \sum_{i=1}^N u(x_i) - \sum_{i=1}^N v(y_{i+1}) \\ &= \sum_{i=1}^N u(x_i) - v(y_{i+1}) \leq \sum_{i=1}^N c(x_i, y_{i+1}) \end{aligned}$$

Example: monotonicity

Let

$$c(x, y) = \frac{1}{2}|x - y|^2, \quad x, y \in \mathbb{R}^d.$$

For two contact pairs,

$$(x_1, y_1), (x_2, y_2) \in \Gamma,$$

cyclical monotonicity gives

$$|x_1 - y_1|^2 + |x_2 - y_2|^2 \leq |x_1 - y_2|^2 + |x_2 - y_1|^2.$$

$$\begin{aligned} -x_1 \cdot y_1 - x_1 \cdot y_2 &\leq -x_1 \cdot y_2 - x_2 \cdot y_1 \\ x_1(y_2 - y_1) - x_2(y_2 - y_1) &\leq 0 \\ (x_1 - x_2) \cdot (y_2 - y_1) &\geq 0 \end{aligned} \quad \left| \begin{array}{l} \leadsto y = y(x) \\ \text{is increasing} \end{array} \right.$$

Theorem: optimal plans are c -cyclically monotone

Let X, Y be Polish spaces and let

$$c : X \times Y \rightarrow (-\infty, +\infty]$$

be lower semicontinuous and bounded from below. Let

$$\pi \in \Pi(\rho, \mu)$$

be an optimal transport plan.

Then π is concentrated on a c -cyclically monotone set: there exists $\Gamma \subset X \times Y$ such that

$$\pi(\Gamma) = 1,$$

and for every $N \geq 1$, every $(x_i, y_i)_{i=1}^N \subset \Gamma$, and every (cyclic) permutation $\sigma \in S_N$,

$$\sum_{i=1}^N c(x_i, y_i) \leq \sum_{i=1}^N c(x_i, y_{\sigma(i)}).$$

References: Villani, *Topics in Optimal Transportation*, Thm. 2.3; Ambrosio–Gigli–Savaré, *Gradient Flows*, Ch. 6; Santambrogio, *Optimal Transport for Applied Mathematicians*, Ch. 1.

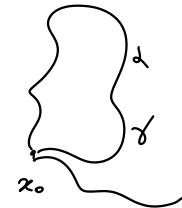
Closed one-forms versus cyclic work

Smooth Poincare lemma:

$$d\alpha = 0 \implies \alpha = d\phi.$$

Work along closed curves:

$$\int_{\gamma} \alpha = 0.$$



Nonsmooth convex analogue (Rockafellar's theorem):

$$\Gamma \text{ cyclically monotone} \implies \boxed{\Gamma \subset \partial\phi.} \quad \not\subset \text{convex}$$

$w.r.t. \quad -x \cdot y = c(x, y)$

Cyclical monotonicity is nonsmooth closedness.

Classical Poincare construction

If $d\alpha = 0$ on a simply connected domain, define

$$\phi(x) = \phi(x_0) + \int_{x_0}^x \alpha.$$



Closedness gives path independence.

closedness $\implies$ path-independent work $\implies$ potential.

Rockafellar construction

Let $\Gamma \subset \mathbb{R}^d \times \mathbb{R}^d$ be cyclically monotone (w.r.t $c(x, p) = -x \cdot p$).

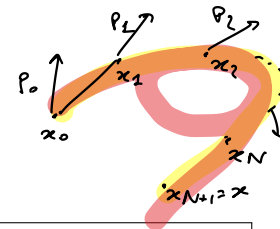
Fix $(x_0, p_0) \in \Gamma$. Define

$$\phi(x) = \sup_{\text{chains}} \sum_i p_i \cdot (x_{i+1} - x_i).$$

"work along the chain"

Here

$$(x_i, p_i) \in \Gamma, \quad x_{N+1} = x.$$



Line integral $\int \alpha$ becomes discrete work along chains.

Cyclic monotonicity is the closedness condition

For a cyclically monotone set, $\sum_{i=0}^N p_i \cdot (x_{i+1} - x_i) \leq 0$ whenever $x_{N+1} = x_0$.

$\Rightarrow$ closed discrete loops have no positive work.

This prevents the chain formula

$$\phi(x) = \sup_{\text{chains ending at } x} \sum_i p_i \cdot (x_{i+1} - x_i)$$

from being made larger by inserting artificial cycles.

Cyclical monotonicity is the no-positive-work condition.
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The potential has the right subgradients

Define

$$\phi(x) := \sup_{\text{chains ending at } x} \sum_i p_i \cdot (x_{i+1} - x_i).$$

Take a point $(\bar{x}, \bar{p}) \in \Gamma$. Any chain ending at $\bar{x}$ can be extended to end at x by adding the step

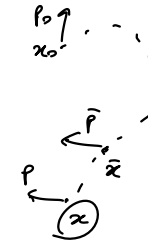
$$(\bar{x}, \bar{p}) : \quad \bar{p} \cdot \underbrace{(x - \bar{x})}.$$

Therefore

$$\phi(x) \geq \phi(\bar{x}) + \bar{p} \cdot (x - \bar{x}).$$

This means

$$\boxed{\bar{p} \in \partial\phi(\bar{x}).}$$



Theorem: c -Rockafellar

Let X, Y be Polish spaces and let

$$c : X \times Y \rightarrow [0, +\infty]$$

be lower semicontinuous.

If

$$\Gamma \subset X \times Y$$

is c -cyclically monotone and $c < +\infty$ on Γ , then there exist admissible potentials u, v such that

$$u(x) \leq v(y) + c(x, y) \quad \forall x, y,$$

and

$$u(x) = v(y) + c(x, y) \quad \forall (x, y) \in \Gamma.$$

In other words,

$$\Gamma \subset \{(x, y) : u(x) = v(y) + c(x, y)\}.$$

Conversely, every subset of the contact set of admissible potentials is c -cyclically monotone.

Reference: Villani, *Topics in Optimal Transportation*, Thm. 2.3.

The c -Rockafellar construction

Given $\Gamma \subset X \times Y$, c -cyclically monotone. Fix a base point $(x_0, y_0) \in \Gamma$.

With our sign convention, the chain formula is an infimum.

Define u by a infimum over chains:

$$u(x) := \inf_{\text{chains}} \left\{ \sum_{i=0}^N \overbrace{[c(x_{i+1}, y_i) - c(x_i, y_i)]}^{\text{work from } x_i \text{ to } x_{i+1} \rightsquigarrow -p_i(x_{i+1}-x_i)} \right\},$$

where $(x_i, y_i) \in \Gamma$, $x_{N+1} = x$.

Then define $\underline{v} = \hat{Q}u = \sup_x \{u(x) - c(x, y)\}$.

Why the c -construction works

By construction

$$u(x) \leq v(y) + c(x, y) \quad \forall x, y. \quad \text{ok}$$

Now take $(\bar{x}, \bar{y}) \in \Gamma$. The chain construction gives

$$\hat{Q}_u(\bar{y}) = \sup_x \{u(x) - c(x, \bar{y})\} \leq u(\bar{x}) - c(\bar{x}, \bar{y}) \quad (\forall x)$$

Taking the supremum over x ,

$$\underline{v}(\bar{y}) = \sup_x \{u(x) - c(x, \bar{y})\} \leq u(\bar{x}) - c(\bar{x}, \bar{y}).$$

Together with admissibility,

$$u(\bar{x}) \leq v(\bar{y}) + c(\bar{x}, \bar{y}),$$

we obtain

$$u(\bar{x}) = v(\bar{y}) + c(\bar{x}, \bar{y}).$$



Duality from Rockafellar

From optimality:

$\text{supp } \pi$ is c -cyclically monotone.

By c -Rockafellar, there exist u, v such that

$$u(x) \leq v(y) + c(x, y),$$

$=$

and equality holds on $\text{supp } \pi$.

Therefore

$$\int c \, d\pi = \int (u(x) - v(y)) \, d\pi = \int u \, d\rho - \int v \, d\mu.$$

Main route:

optimality $\Rightarrow$ c -cyclical monotonicity $\Rightarrow$ potential $\Rightarrow$ dual equality.

Lectures 01–02 recap

1. **Dual potentials as prices.** No-arbitrage condition:

$$u(x) \leq v(y) + c(x, y)$$

2. **Hard c -transforms.**

$$Q_c v(x) = \inf_y \{v(y) + c(x, y)\}, \quad \hat{Q}_c u(y) = \sup_x \{u(x) - c(x, y)\}.$$

3. **Softmin and entropy.**

$$Q_c^\varepsilon v = -\varepsilon \log \int e^{-(v+c)/\varepsilon} d\mu.$$

4. **Zero-temperature limit.**

$$Q_c^\varepsilon \longrightarrow Q_c, \quad \text{Gibbs kernels} \longrightarrow \text{contact sets}.$$

5. **Rockafellar theorem.**

$$c\text{-cyclical monotonicity} \iff \text{existence of contact potentials}.$$