Conference on Discrete Geometry and Optimization Toronto, September 2011

Exploiting Polyhedral Symmetries in Social Choice Theory

Achill Schürmann (University of Rostock)

arXiv:1109.1545

(based on work "supported" by two Bachelor projects at TU Delft)

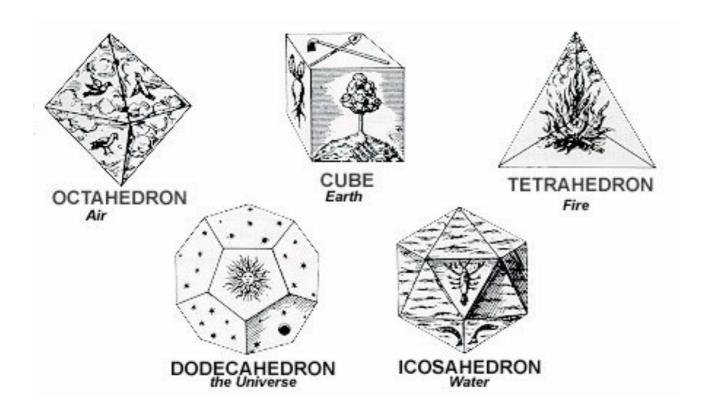
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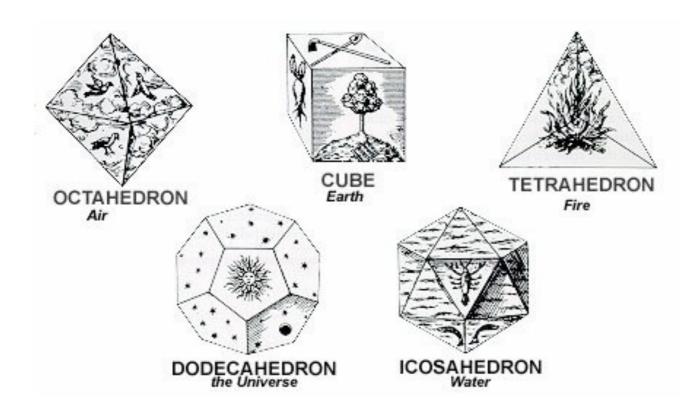
Exploiting Polyhedral Symmetries in Social Choice Theory and elsewhere

Achill Schürmann (University of Rostock)

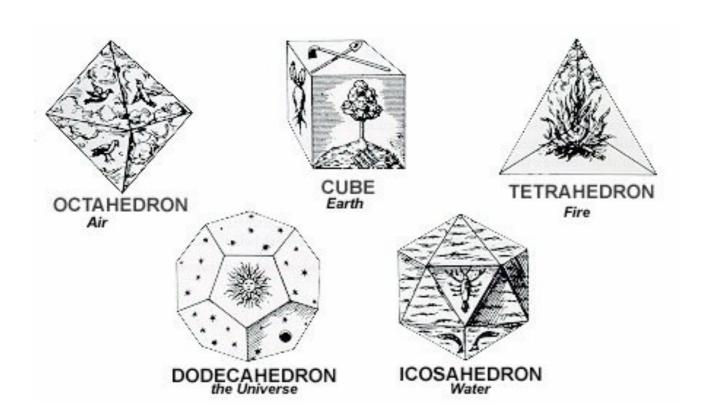
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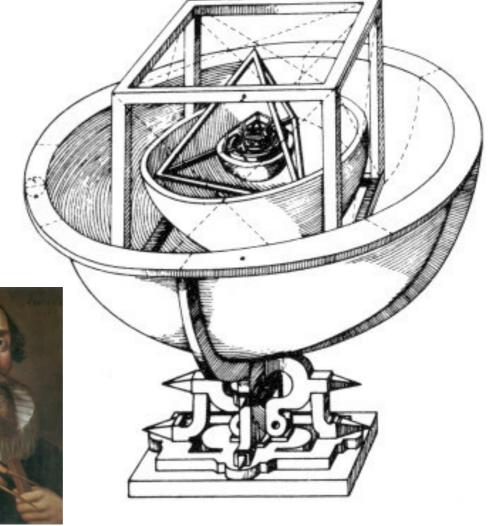


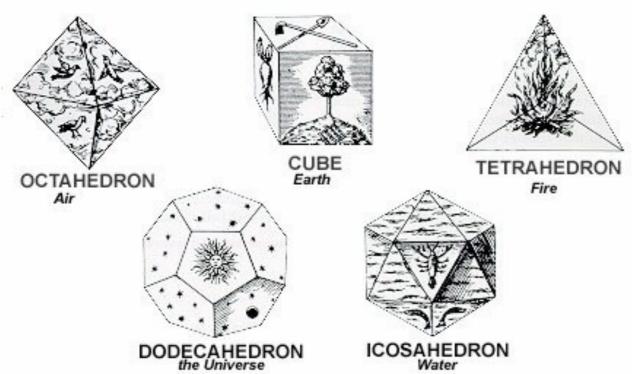




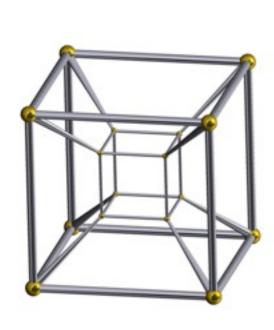


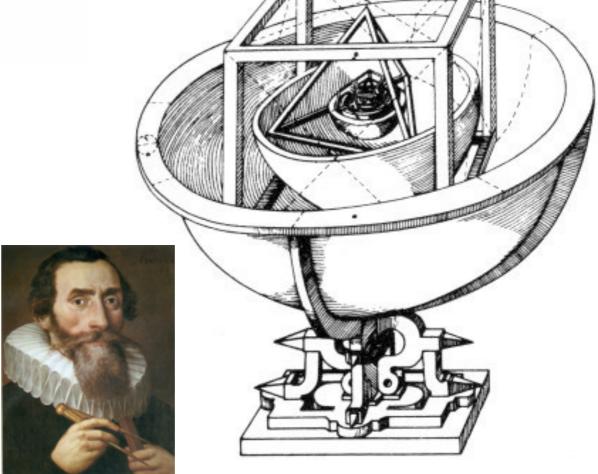


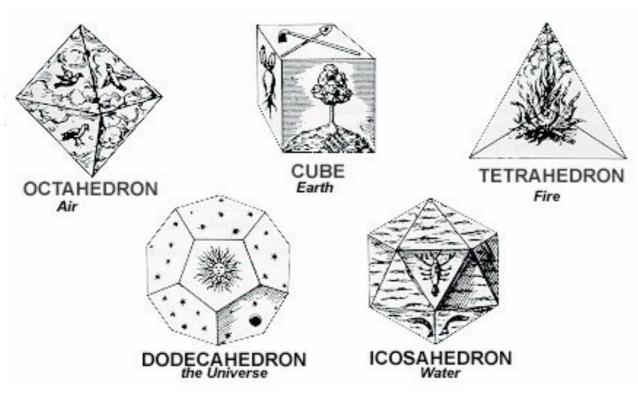


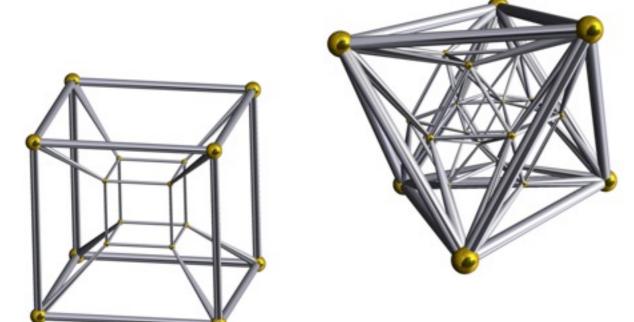




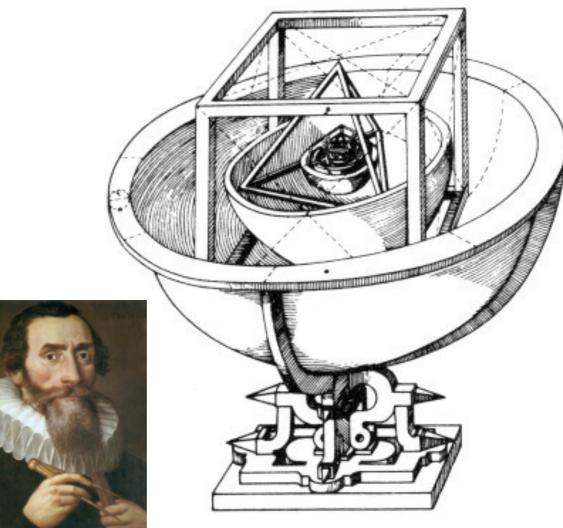












individual choices

collective choice

```
a b c b a  
> > > > >  
b c a c c ...

b c a c c ...

> > > > > >  
c a b a b
```

individual choices

collective choice

```
a b c b a
> > > >
b c a c c c
> > > >
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```







individual choices

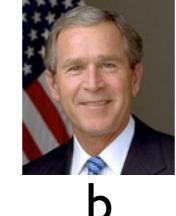
 \longrightarrow

collective choice

a b c b a
> > > >
b c a c c
> > > >
c a b a b

a > b >







individual choices

 \longrightarrow

collective choice

a b c b a
> > > >
b c a c c
> > > >
c a b a b

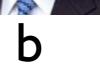


a













Kenneth Arrow (Nobel prize 1972)

THM: There is no fair voting system.



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THM: There is no voting system, which is (for at least three choices)

not a dictatorship



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- respecting binary preferences made by all



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(preference between a and b depends only on individual preferences between a and b)



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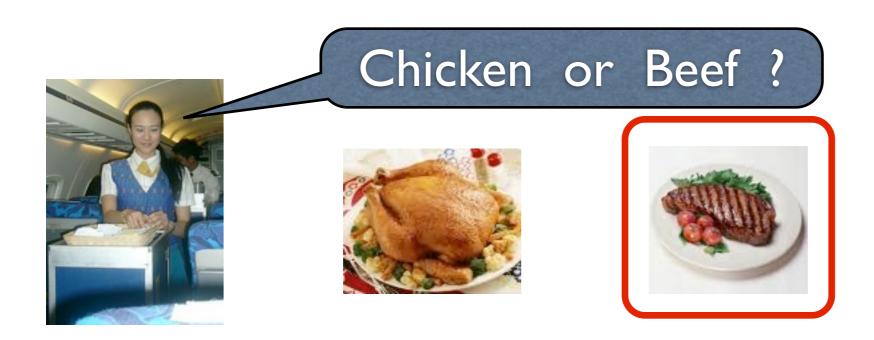
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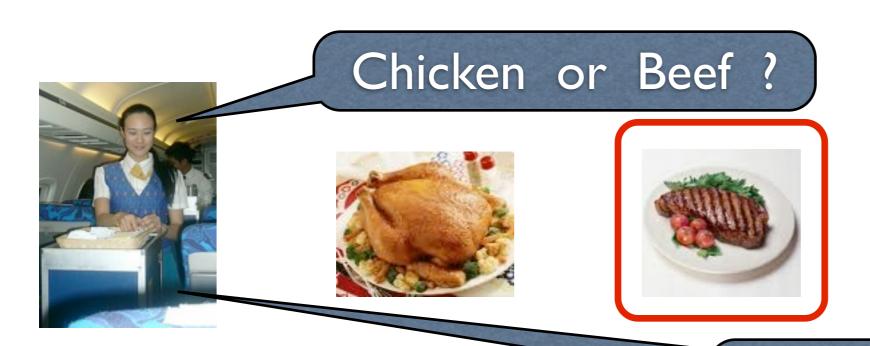
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We also have fish...



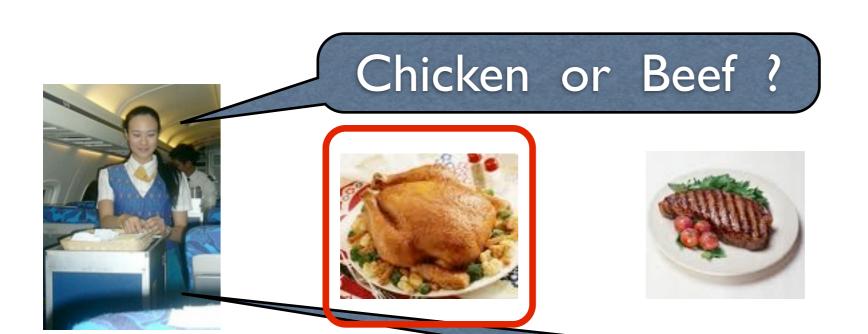
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We also have fish...

collective choice can be intransitive!



Marquis de Condorcet (1743-1793)

collective choice can be intransitive!

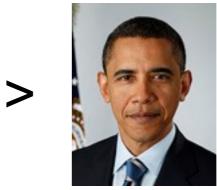


Marquis de Condorcet (1743-1793)









•••

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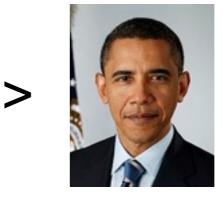


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THUS: There may be no "pairwise winner"! (Condorcet winner)

 Impartial Anonymous Culture (IAC) assumption: every voting situation is equally likely

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- for three candidates a, b and c, let

```
n_{\rm ab} number of voters with choice \, {\rm a} > {\rm b} > {\rm c} \, n_{\rm ac} number of voters with choice \, {\rm a} > {\rm c} > {\rm b} \, n_{\rm ba} number of voters with choice \, {\rm b} > {\rm a} > {\rm c} \,
```

• • •

- Impartial Anonymous Culture (IAC) assumption: every voting situation is equally likely
- for three candidates a, b and c, let

```
n_{\rm ab} number of voters with choice a>b>c n_{\rm ac} number of voters with choice a>c>b n_{\rm ba} number of voters with choice b>a>c
```

• • •

 $(n_{ab}, n_{ac}, n_{ba}, n_{bc}, n_{ca}, n_{cb})$ describes a voting situation

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 n_{ac} number of voters with choice a > c > b

 $n_{\rm ba}$ number of voters with choice b > a > c

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 $(n_{ab}, n_{ac}, n_{ba}, n_{bc}, n_{ca}, n_{cb})$ describes a voting situation

$$N = n_{ab} + n_{ac} + n_{ba} + n_{bc} + n_{ca} + n_{cb}$$

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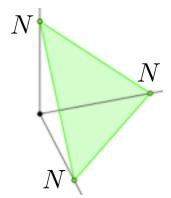
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Counting Lattice Points

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Candidate a is a Condorcet winner if

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$$n_{\mathsf{ab}} + n_{\mathsf{ac}} + n_{\mathsf{ca}} > n_{\mathsf{ba}} + n_{\mathsf{bc}} + n_{\mathsf{cb}}$$
 (a beats b)

Counting Lattice Points

Candidate a is a Condorcet winner if

$$n_{\sf ab}+n_{\sf ac}+n_{\sf ca}>n_{\sf ba}+n_{\sf bc}+n_{\sf cb}$$
 (a beats b) and $n_{\sf ab}+n_{\sf ac}+n_{\sf ba}>n_{\sf ca}+n_{\sf cb}+n_{\sf bc}$ (a beats c)

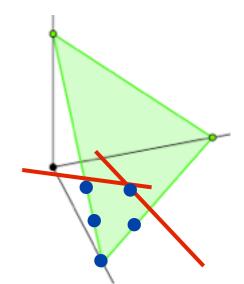
Counting Lattice Points

Candidate a is a Condorcet winner if

(1)
$$n_{ab} + n_{ac} + n_{ca} > n_{ba} + n_{bc} + n_{cb}$$
 (a beats b)

(2) and
$$n_{ab}+n_{ac}+n_{ba}>n_{ca}+n_{cb}+n_{bc}$$
 (a beats c)

That is: $(n_{\mathsf{ab}}, n_{\mathsf{ac}}, n_{\mathsf{ba}}, n_{\mathsf{bc}}, n_{\mathsf{ca}}, n_{\mathsf{cb}}) \in \mathbb{Z}_{\geq 0}^6$



is in the polyhedron

$$P_N = \left\{ \ n \in \mathbb{R}^6 \mid \ N = \sum_{\mathsf{xy}} n_{xy}, \ n_{xy} \geq 0 \ \ \mathsf{and} \ \ \underline{(1),(2)} \
ight\}$$

$$\#(P_N \cap \mathbb{Z}^d) = a_{d-1}N^{d-1} + \ldots + a_1N + a_0$$

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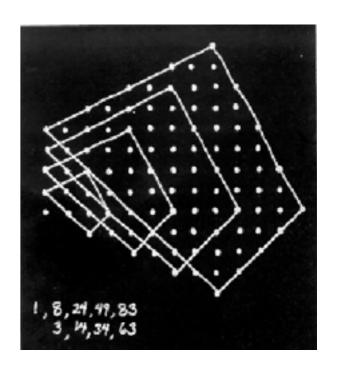
Eugène Ehrhart (1906-2000)

$$\#(P_N \cap \mathbb{Z}^d) = a_{d-1}N^{d-1} + \ldots + a_1N + a_0$$

• P_1 integral \Rightarrow polynomial



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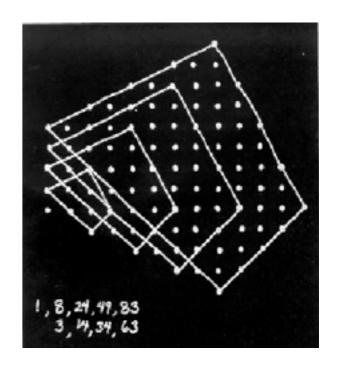
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$$P_1$$
 integral \Rightarrow polynomial

Ex:
$$P_1 = \text{conv}\{e_1, \dots, e_d\} \implies \#(P_N \cap \mathbb{Z}^d) = \binom{N+d-1}{d-1}$$



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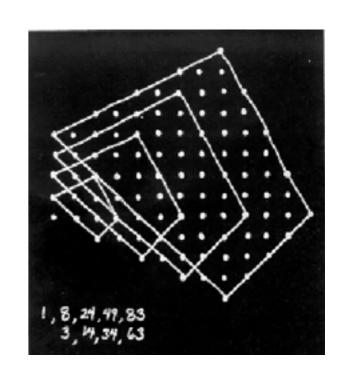


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$$\text{vol}_{d-1}(P_1)$$

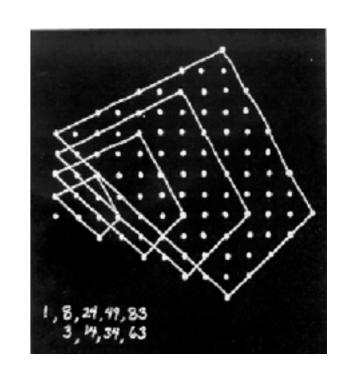


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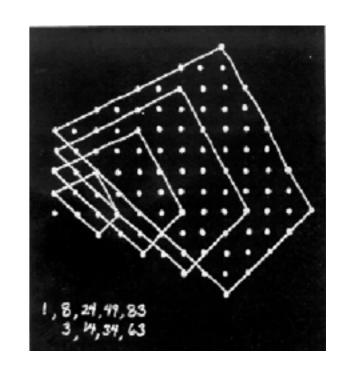


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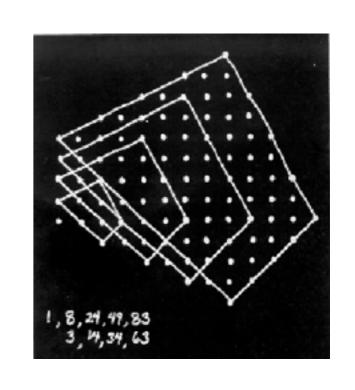


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- P_1 rational \Rightarrow quasi-polynomial
- "Reinvented" in Social Choice Theory by Chua and Huang (2000)
- Parallelity of Approach discovered in 2006 (by Lepelley et al. and Wilson / Pritchard)



Quasi-polynomial for $\#(P_N\cap \mathbb{Z}^6)$ can be obtained using barvinok or latte



Quasi-polynomial for $\#(P_N \cap \mathbb{Z}^6)$ can be obtained

using barvinok or latte

```
1/384 * N^5
+ ( -1/64 * {( 1/2 * N + 0 )} + 3/64 ) * N^4
+ ( -19/96 * {( 1/2 * N + 0 )} + 31/96 ) * N^3
+ ( -29/32 * {( 1/2 * N + 0 )} + 17/16 ) * N^2
+ ( -343/192 * {( 1/2 * N + 0 )} + 5/3 ) * N
+ ( -83/64 * {( 1/2 * N + 0 )} + 1 )
```

(Number of voting situations with N voters and candidate a as Condorcet winner)



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(Number of voting situations with N voters and candidate a as Condorcet winner)

Likeliness of Condorcet Paradox

$$1 - 3\frac{\mathsf{q}\text{-poly}}{\binom{N+5}{5}}$$

For large elections $(N \to \infty)$:

$$1 - 3\frac{1/384}{1/120} = \frac{1}{16} = 0.0625$$



Condorcet winner, but Plurality loser

Condorcet winner, but Plurality loser

$$n_{\rm ab} + n_{\rm ac} + n_{\rm ca} > n_{\rm ba} + n_{\rm bc} + n_{\rm cb}$$
 (a beats b) $n_{\rm ab} + n_{\rm ac} + n_{\rm ba} > n_{\rm ca} + n_{\rm cb} + n_{\rm bc}$ (a beats c)

Condorcet winner, but Plurality loser

```
n_{\rm ab}+n_{\rm ac}+n_{\rm ca}>n_{\rm ba}+n_{\rm bc}+n_{\rm cb} (a beats b) n_{\rm ab}+n_{\rm ac}+n_{\rm ba}>n_{\rm ca}+n_{\rm cb}+n_{\rm bc} (a beats c) n_{\rm ba}+n_{\rm bc}>n_{\rm ab}+n_{\rm ac}, n_{\rm ca}+n_{\rm cb} (b wins plurality)
```

Condorcet winner, but Plurality loser

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Likeliness for large elections $(N \to \infty)$: $\frac{16}{135} = 0.1185...$

Condorcet winner, but Plurality loser

$$n_{\rm ab}+n_{\rm ac}+n_{\rm ca}>n_{\rm ba}+n_{\rm bc}+n_{\rm cb}$$
 (a beats b) $n_{\rm ab}+n_{\rm ac}+n_{\rm ba}>n_{\rm ca}+n_{\rm cb}+n_{\rm bc}$ (a beats c) $n_{\rm ba}+n_{\rm bc}>n_{\rm ab}+n_{\rm ac}$, $n_{\rm ca}+n_{\rm cb}$ (b wins plurality)

Likeliness for large elections
$$(N \to \infty)$$
: $\frac{16}{135} = 0.1185...$

• Plurality vs. Plurality Runoff

Condorcet winner, but Plurality loser

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 (a beats b) $n_{\rm ab}+n_{\rm ac}+n_{\rm ba}>n_{\rm ca}+n_{\rm cb}+n_{\rm bc}$ (a beats c) $n_{\rm ba}+n_{\rm bc}>n_{\rm ab}+n_{\rm ac}$, $n_{\rm ca}+n_{\rm cb}$ (b wins plurality)

Likeliness for large elections $(N \to \infty)$: $\frac{16}{135} = 0.1185...$

Plurality vs. Plurality Runoff

$$n_{\rm ab}+n_{\rm ac}>n_{\rm ba}+n_{\rm bc}$$
 (a wins plurality over b) $n_{\rm ba}+n_{\rm bc}>n_{\rm ca}+n_{\rm cb}$ (b wins plurality over c) $n_{\rm ab}+n_{\rm ac}+n_{\rm ca}< n_{\rm ba}+n_{\rm bc}+n_{\rm cb}$ (b beats a)

Condorcet winner, but Plurality loser

$$n_{\rm ab}+n_{\rm ac}+n_{\rm ca}>n_{\rm ba}+n_{\rm bc}+n_{\rm cb}$$
 (a beats b) $n_{\rm ab}+n_{\rm ac}+n_{\rm ba}>n_{\rm ca}+n_{\rm cb}+n_{\rm bc}$ (a beats c) $n_{\rm ba}+n_{\rm bc}>n_{\rm ab}+n_{\rm ac}$, $n_{\rm ca}+n_{\rm cb}$ (b wins plurality)

Likeliness for large elections $(N \to \infty)$: $\frac{16}{135} = 0.1185...$

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Likeliness for large elections $(N \to \infty)$: $\frac{71}{576} = 0.12326...$





















hardly any exact probabilitie









hardly any exact probabilitie

for 4 candidates 24 variables are used in polyhedral model









hardly any exact probabilitie

- for 4 candidates 24 variables are used in polynegral model
 - => polyhedral computations are too difficult









hardly any exact probabilitie

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("most of the time", due to LattE integrale, July 2011)









hardly any exact probabilitie

- for 4 candidates 24 variables are used in polyhedral model
 - => polyhedral computations are too difficult

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IDEA: Reduce dimension by exploiting symmetry!

$$n_{\mathsf{ab}} + n_{\mathsf{ac}} + n_{\mathsf{ca}} > n_{\mathsf{ba}} + n_{\mathsf{bc}} + n_{\mathsf{cb}}$$

$$n_{\mathsf{ab}} + n_{\mathsf{ac}} + n_{\mathsf{ba}} > n_{\mathsf{ca}} + n_{\mathsf{cb}} + n_{\mathsf{bc}}$$

$$N = n_{ab} + n_{ac} + n_{ba} + n_{ca} + n_{bc} + n_{cb}$$

$$n_{\mathsf{ab}} + n_{\mathsf{ac}} + n_{\mathsf{ca}} > n_{\mathsf{ba}} + n_{\mathsf{bc}} + n_{\mathsf{cb}}$$

$$n_{\mathsf{ab}} + n_{\mathsf{ac}} + n_{\mathsf{ba}} > n_{\mathsf{ca}} + n_{\mathsf{cb}} + n_{\mathsf{bc}}$$

$$N = n_{\mathsf{ab}} + n_{\mathsf{ac}} + n_{\mathsf{ba}} + n_{\mathsf{ca}} + n_{\mathsf{bc}} + n_{\mathsf{cb}}$$

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$$n_{\mathsf{ab}} + n_{\mathsf{ac}} + n_{\mathsf{ba}} > n_{\mathsf{ca}} + n_{\mathsf{cb}} + n_{\mathsf{bc}}$$

$$N = n_{ab} + n_{ac} + n_{ba} + n_{ca} + n_{bc} + n_{cb}$$

 n_{a}

 n_{R}

$$n_{\mathsf{a}} + n_{\mathsf{ca}} > n_{\mathsf{ba}} + n_{\mathsf{R}}$$

$$n_{\mathsf{a}} + n_{\mathsf{ba}} > n_{\mathsf{ca}} + n_{\mathsf{R}}$$

$$N = n_{\mathsf{a}} + n_{\mathsf{ba}} + n_{\mathsf{ca}} + n_{\mathsf{R}}$$

 n_{a}

 n_{R}

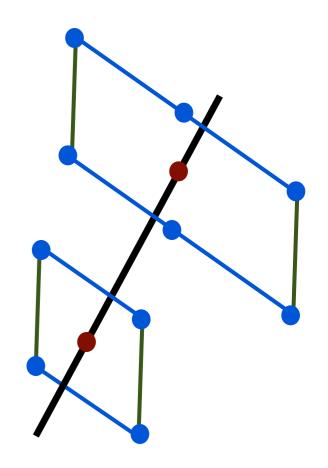
Grouping of variables

$$n_{\mathsf{a}}$$
 $+ n_{\mathsf{ca}} > n_{\mathsf{ba}} + n_{\mathsf{R}}$ n_{R} $+ n_{\mathsf{ba}} > n_{\mathsf{ca}} + n_{\mathsf{R}}$ $N = n_{\mathsf{a}} + n_{\mathsf{ba}} + n_{\mathsf{ca}} + n_{\mathsf{R}}$ n_{R}

$$(n_{\mathsf{a}}, n_{\mathsf{ba}}, n_{\mathsf{ca}}, n_{\mathsf{R}})$$
 describes $(n_{\mathsf{a}} + 1)(n_{\mathsf{R}} + 1)$ voting situations (former lattice points)

Grouping of variables

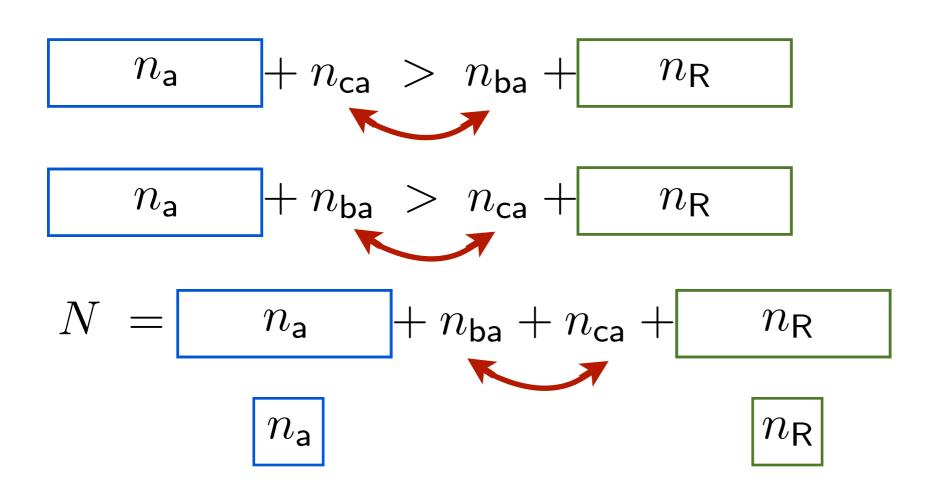
$$n_{\rm a}$$
 $+ n_{\rm ca}$ $> n_{\rm ba}$ $+ n_{\rm R}$ $+ n_{\rm ba}$ $> n_{\rm ca}$ $+ n_{\rm R}$ $N = n_{\rm a}$ $+ n_{\rm ba}$ $+ n_{\rm ba}$ $+ n_{\rm ca}$ $+ n_{\rm R}$ $+ n_{\rm ca}$ $+ n_{\rm R}$

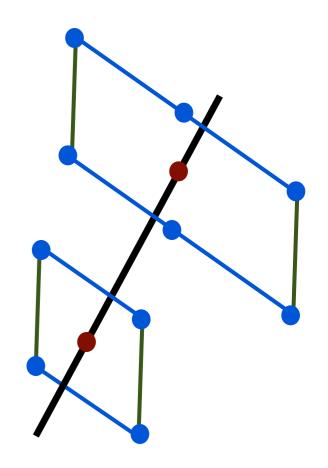


 $(n_{\mathsf{a}}, n_{\mathsf{ba}}, n_{\mathsf{ca}}, n_{\mathsf{R}})$ describes $(n_{\mathsf{a}} + 1)(n_{\mathsf{R}} + 1)$ voting situations (former lattice points)

THUS: the polytope decomposes into fibers of simplotopes (cross products of simplices)

Grouping of variables





 $(n_{\mathsf{a}}, n_{\mathsf{ba}}, n_{\mathsf{ca}}, n_{\mathsf{R}})$ describes $(n_{\mathsf{a}} + 1)(n_{\mathsf{R}} + 1)$ voting situations (former lattice points)

THUS: the polytope decomposes into fibers of simplotopes (cross products of simplices)

$$\mathsf{Prob}(N) \ = \ rac{\left|L_N \cap P \cap \mathbb{Z}^d
ight|}{\left|L_N \cap S \cap \mathbb{Z}^d
ight|}$$

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$$\lim_{N \to \infty} \mathsf{Prob}(N) \ = \ \lim_{N \to \infty} \frac{\left| L_1 \cap P \cap (\mathbb{Z}/N)^d \right|}{\left| L_1 \cap S \cap (\mathbb{Z}/N)^d \right|}$$

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Large elections with four candidates

Large elections with four candidates

No Condorcet winner exists (Condorcet paradox)

$$\lim_{N \to \infty} \mathsf{Prob}(N) \; = \; \frac{331}{2048} = 0.1616\dots$$

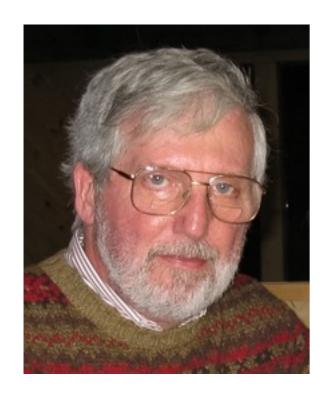
(by integrating polynomial of degree 16 over a 7-dimensional polytope)

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William V. Gehrlein

In an email of Sep. 7th 2011:

Your results particularly got my attention when I finally realized that you had obtained limiting representations for four candidates. This is a significant step forward, and you are not the only person who has been trying to produce such results. However, I believe that you are the first to successfully accomplish this. The only four candidate result that I am aware of is cited in your paper, and I only managed to obtain that by using a trick.

New results with four candidates

Condorcet Efficiency of Plurality

$$\lim_{N\to\infty} \mathsf{Prob}(N) \ = \ \frac{10658098255011916449318509}{14352135440302080000000000} \ = \ 0.74261\dots$$

(by integrating polynomial of degree 11 over a 13-dimensional polytope)

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Plurality vs. Plurality Runoff

$$\lim_{N \to \infty} \mathsf{Prob}(N) \; = \; \frac{2988379676768359}{12173449145352192} = 0.24548 \ldots$$

(by integrating polynomial of degree 18 over a 5-dimensional polytope)

WANT: generalization of Ehrhart theory, counting lattice points with polynomial weights

- Two new methods:
 - via rational generating functions
 - via local Euler-Maclaurin formula



Baldoni, Berline, Vergne, 2009

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 - via rational generating functions
 - via local Euler-Maclaurin formula
- "experimental" implementation available in barvinok



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Want:

Methods exploiting general polyhedral symmetry groups

Exploiting Symmetry in other Polyhedral Computations?

Recent computational successes: (with Mathieu Dutour Sikirić and Frank Vallentin)



• Classification of eight dimensional perfect forms, Electron. Res. Announc. AMS, 13 (2007)

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- The contact polytope of the Leech lattice, preprint at arXiv:0906.1427
 - I orbit with 196,560 vertices in 24 dimensions
 - 1,197,362,269,604,214,277,200 many facets in 232 orbits







helps to compute linear automorphism groups



- helps to compute linear automorphism groups
- converts polyhedral representations using

Recursive Decomposition Methods (Incidence/Adjacency)

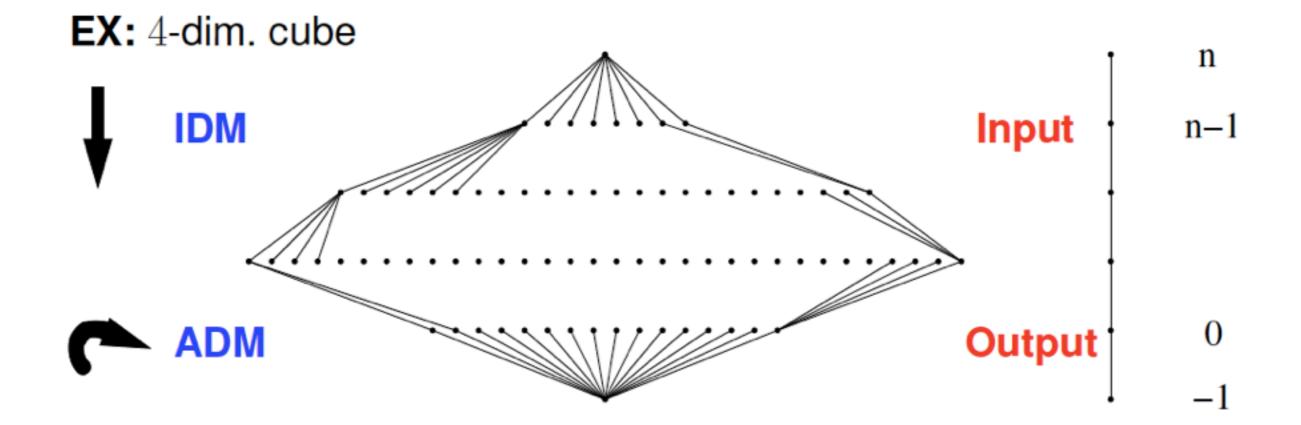
(also used by Christof/Reinelt, Deza/Fukuda/Pasechnik, ...)



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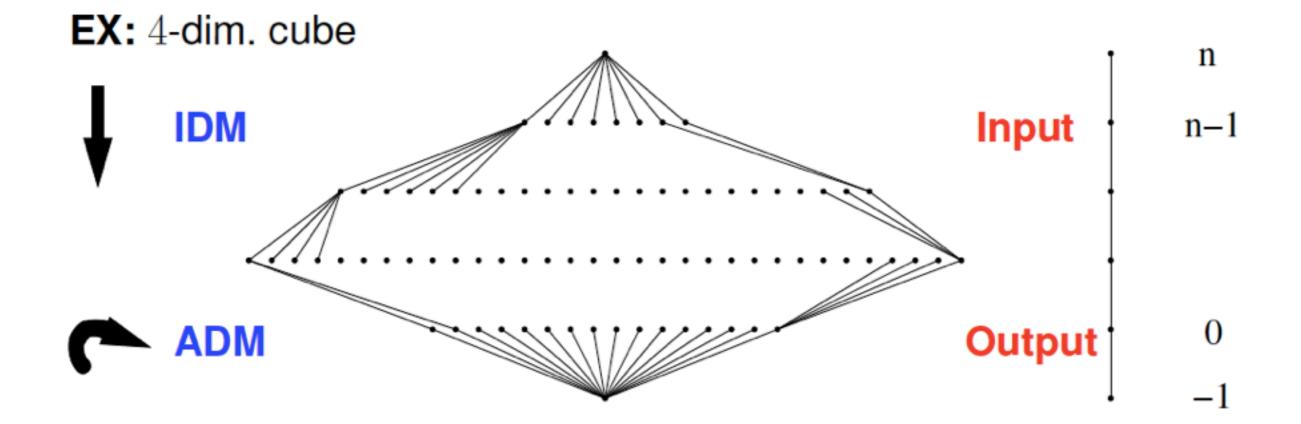
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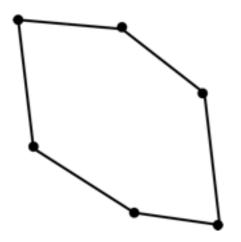
Symmetry Groups

Symmetry Groups

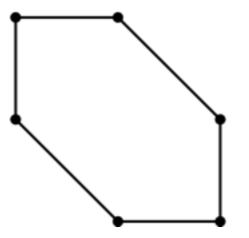
• Combinatorial, Linear, or Geometric Symmetries

Symmetry Groups

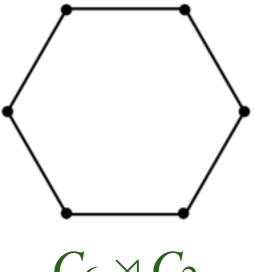
Combinatorial, Linear, or Geometric Symmetries



 $C_6 \rtimes C_2$ trivial
trivial



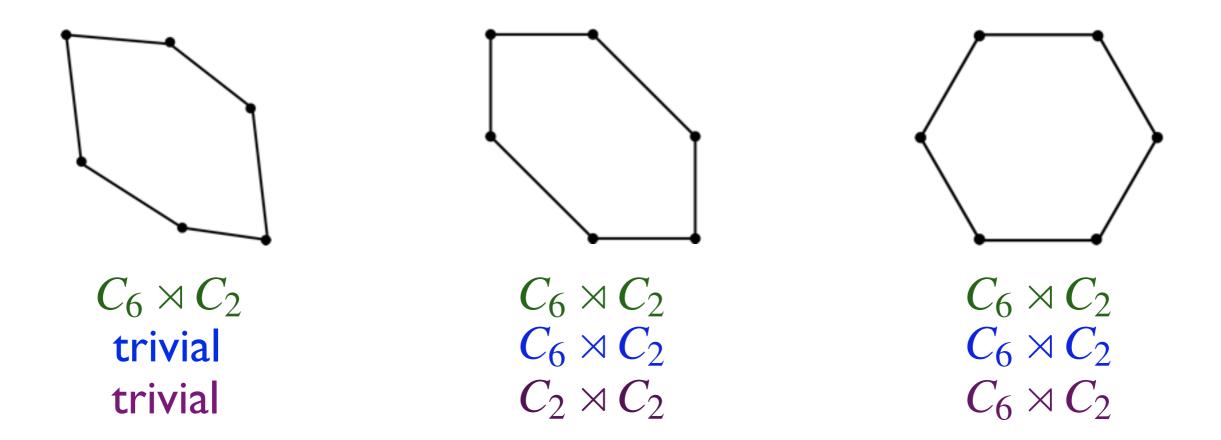
$$C_6 \rtimes C_2$$
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Symmetry Groups

Combinatorial, Linear, or Geometric Symmetries



DEF: A linear automorphism of $\{v_1, \dots, v_m\} \subset \mathbb{R}^n$ is a regular matrix $A \in \mathbb{R}^{n \times n}$ with $Av_i = v_{\sigma(i)}$ for some $\sigma \in S_m$

THM: The group of linear automorphisms is equal to

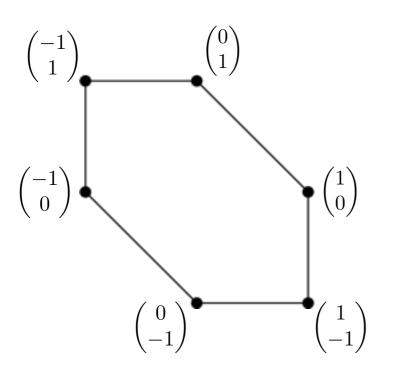
the automorphism group of the complete graph K_m

with edge labels
$$v_i^t Q^{-1} v_j$$
, where $Q = \sum_{i=1}^m v_i v_i^t$

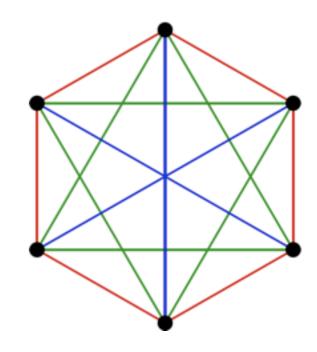
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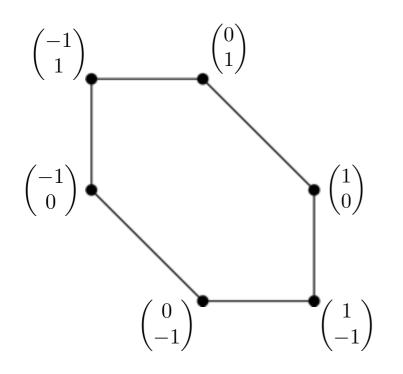
$$Q = \begin{pmatrix} 4 & -2 \\ -2 & 4 \end{pmatrix}$$



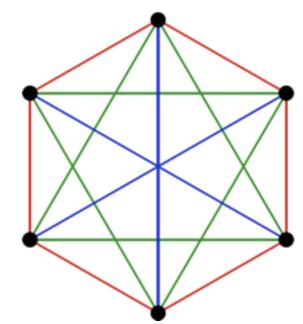
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=> use NAUTY by Brendan McKay



(for vertex enumeration)

(for vertex enumeration)



Find initial orbit(s) / representing vertice(s)

(for vertex enumeration)



• Find initial orbit(s) / representing vertice(s)



- For each new orbit representative
 - enumerate neighboring vertices

(for vertex enumeration)



• Find initial orbit(s) / representing vertice(s)



For each new orbit representative





add as orbit representative if in a new orbit

(for vertex enumeration)



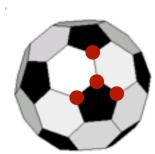
Find initial orbit(s) / representing vertice(s)



For each new orbit representative



enumerate neighboring vertices (up to symmetry)



add as orbit representative if in a new orbit

Representation conversion problem

(for vertex enumeration)

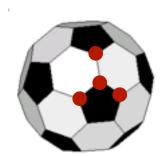


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Representation conversion problem

BOTTLENECK: Stabilizer and In-Orbit computations

(for vertex enumeration)



• Find initial orbit(s) / representing vertice(s)



For each new orbit representative





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Representation conversion problem

BOTTLENECK: Stabilizer and In-Orbit computations

=> Need of efficient data structures and algorithms for permutation groups: BSGS, (partition) backtracking

Ingredient I: Permutation Group Algorithms

 BSGS and (partition) backtrack could be provided by GAP, MAGMA or SAGE

Ingredient I: Permutation Group Algorithms

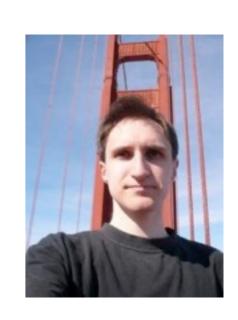
- BSGS and (partition) backtrack
 could be provided by GAP, MAGMA or SAGE
- We use the callable C++ library PermLib
 - open source (new BSD license)
 - with compact API to access core functionality
 - can replace NAUTY





Ingredient I: Permutation Group Algorithms

- Perm Lib
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Vision:

Create "integrated algorithms" combining tools of

Polyhedral Combinatorics and Computational Group Theory

Ingredient II:

Established Representation Conversion Tools

cddlib by Komei Fukuda (Double Description Method)
 incrementally adding inequalities and recomputing vertices at every step



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Irslib by David Avis (Lexicographic Reverse Search)
 pivoting using "Simplex Pivots"



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WHAT ABOUT Symmetry Exploiting Methods ?

cddlib by Komei Fukuda (Double Description Method)
 incrementally adding inequalities and recomputing vertices at every step



Irslib by David Avis (Lexicographic Reverse Search)
 pivoting using "Simplex Pivots"



WHAT ABOUT Symmetry Exploiting Methods ?

- with David Bremner we work(ed) on
 - pivoting methods up to symmetry
 - incremental methods using fundamental domains





[Kum11] Abhinav Kumar, Elliptic fibrations on a generic Jacobian Kummer surface, arxiv:1105.1715



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Getting the group:

```
sympol --automorphisms-only input-file
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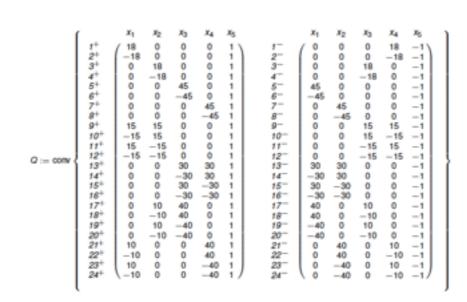
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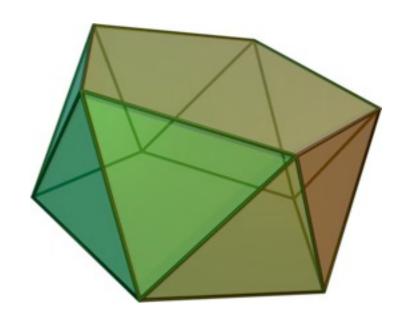
```
sympol --adm 40 input-file
```

```
V-representation
* UP TO SYMMETRY
begin
...
end
permutation group
* order 11520
* w.r.t. to the original inequalities/verti
```

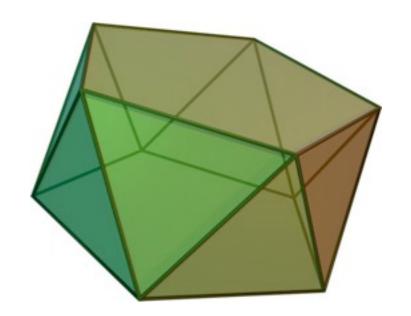






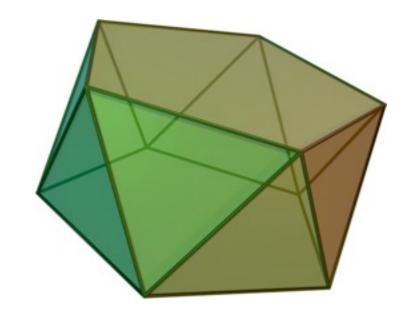






sympol --idm-adm-level 0 1 --adjacencies input-file

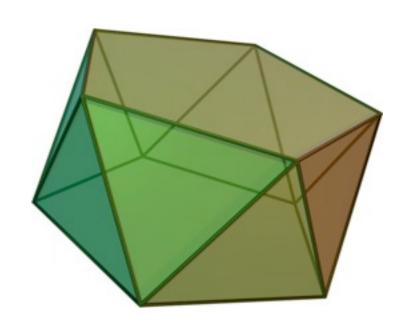




sympol --idm-adm-level 0 1 --adjacencies input-file

```
graph adjacencies {
1 -- 2;
2 -- 4;
2 -- 3;
2 -- 2;
3 -- 10;
3 -- 3;
3 -- 6;
4 -- 5;
4 -- 6;
5 -- 5;
5 -- 7;
5 -- 6;
5 -- 6;
7 -- 7;
7 -- 9;
7 -- 8;
8 -- 8;
9 -- 12:
```

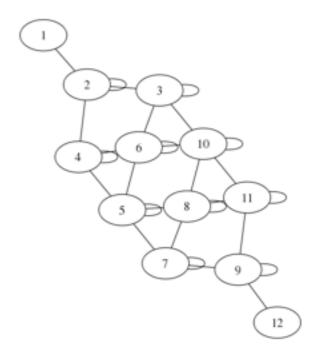




sympol --idm-adm-level 0 1 --adjacencies input-file

```
graph adjacencies {
1 -- 2;
2 -- 4;
2 -- 3;
2 -- 2;
3 -- 10;
3 -- 3;
3 -- 6;
4 -- 5;
4 -- 4;
4 -- 6;
5 -- 5;
5 -- 7;
5 -- 6;
5 -- 8;
6 -- 6;
7 -- 7;
7 -- 9;
7 -- 8;
8 -- 8;
9 -- 12:
```

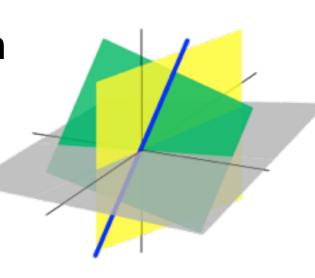
~> neato ~>
 (Graphviz)



What else?

For LPs one can intersect feasible polyhedron

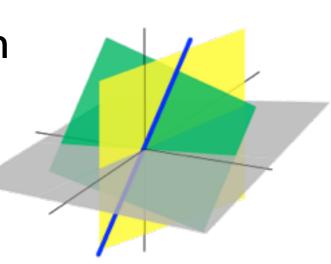
with invariant linear subspace



For LPs one can intersect feasible polyhedron

with invariant linear subspace

(not possible for IPs)



• For LPs one can intersect feasible polyhedron

with invariant linear subspace

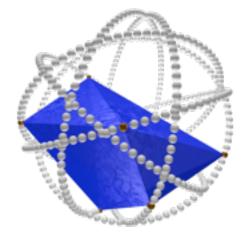
(not possible for IPs)











=> see survey "Symmetry in Integer Linear Programming" by François Margot (2010)



using invariant linear subspace





in Lattice Point Counting





- in Lattice Point Counting
- in Polyhedral Representation Conversions







- in Lattice Point Counting
- in Polyhedral Representation Conversions
- in Integer Programming and MILPs





























- in Lattice Point Counting
- in Polyhedral Representation Conversions
- in Integer Programming and MILPs



























ToDo

- Create efficient computational tools / use more math!
- Integrate tools from Computational Group Theory

Thanks!