

Tameness beyond o-minimality

Philipp Hieronymi

Introductory Workshop on Tame Geometry, Transseries and Applications to Analysis and Geometry, January 2022



Universität Bonn
Mathematisches Institut

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“What might it mean for a first-order expansion of the field of real numbers to be tame or well behaved? In recent years, much attention has been paid by model theorists and real-analytic geometers to the o-minimal setting[...]. But there are expansions of the real field that define sets with infinitely many connected components, yet are tame in some well-defined sense [...]. The analysis of such structures often requires a mixture of model-theoretic, analytic-geometric and descriptive set-theoretic techniques. An underlying idea is that first-order definability, in combination with the field structure, can be used as a tool for determining how complicated is a given set of real numbers.”

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A **structure** on \mathbb{R} is a sequence $\mathfrak{G} := (\mathfrak{G}_m)_{m=1}^{\infty}$ such that for each m :

1. \mathfrak{G}_m is a boolean algebra of subsets of \mathbb{R}^m .
2. If $X \in \mathfrak{G}_m$ and $Y \in \mathfrak{G}_n$, then $X \times Y \in \mathfrak{G}_{m+n}$.
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Let K be a subfield of \mathbb{R} . A **K -geometry** is a structure satisfying

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where $f_1, \dots, f_M, g_1, \dots, g_N$ are K -affine functions and $k, M, N \in \mathbb{N}$, forms a K -geometry.

This follows from the Fourier-Motzkin elimination method. This class contains all K -polyhedra.

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Examples. Semilinear geometry, semialgebraic geometry, subexponential geometry, globally subanalytic geometry,...

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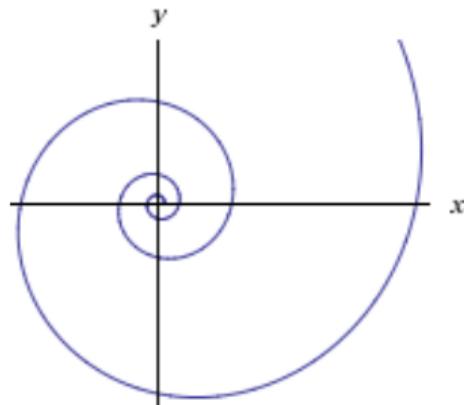
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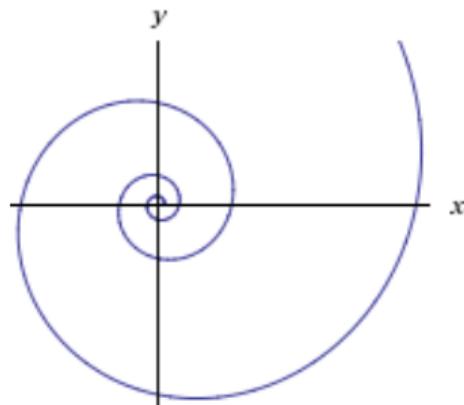
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		dense ω -order				
		no		yes		
field-type	yes	o-minimal	locally o-min.	d-minimal	noiseless	type C
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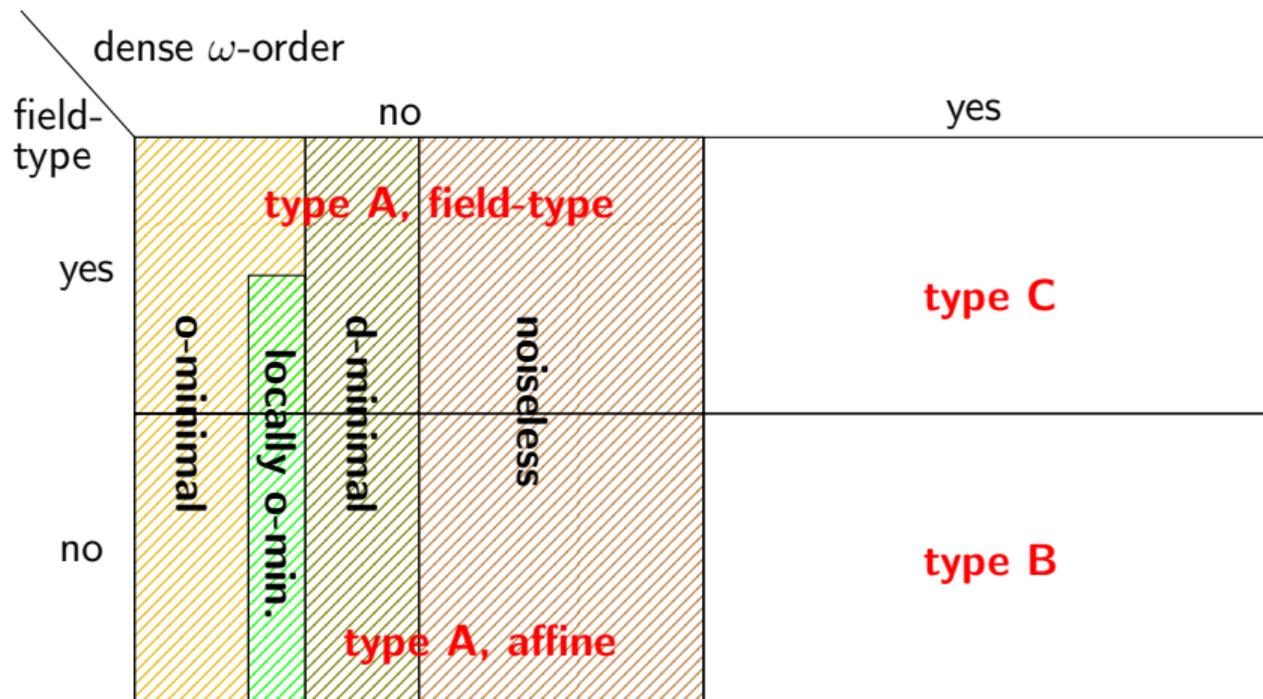
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A few comments about type B geometries.

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