



# **A practical primal-dual interior-point algorithm for nonsymmetric conic optimization**

July 28, 2021

**Erling D. Andersen**

(joint work with Joachim Dahl)

MOSEK ApS,

Email: [e.d.andersen@mosek.com](mailto:e.d.andersen@mosek.com)

Personal WWW: <https://erling.andersens.name>

Company WWW: <https://mosek.com>

[www.mosek.com](http://www.mosek.com)





## Conic optimization

- The problem

- Two nonsymmetric cones

## A primal-dual interior-point algorithm

- Survey of algorithms

- Preliminaries

- Motivation

- The algorithm

## Computational results

## Summary



## Section 1

### Conic optimization





$$\begin{array}{ll} \text{minimize} & \sum (c^k)^T x^k \\ \text{st} & \sum_k A^k x^k = b, \\ & x^k \in \mathcal{K}^k, \quad \forall k, \end{array}$$

where

- $c^k \in \mathcal{R}^{n^k}$ ,
- $A^k \in \mathcal{R}^{m \times n^k}$ ,
- $b \in \mathcal{R}^m$ ,
- $\mathcal{K}^k$  are convex cones.





3 symmetric cones:

- Linear.
- Quadratic.
- Semidefinite.

2 nonsymmetric cones:

- **Exponential.**
- **Power.**

Observation:

- **Almost** all convex optimization problems appearing in practice can be formulated using those 5 cones.





Given  $\alpha_j > 0$  and  $\sum \alpha_j = 1$  then define the power cone:

$$\mathcal{K}_{pow}(\alpha) := \left\{ (x, z) : \prod_{j=1}^n x_j^{\alpha_j} \geq \|z\|, x \geq 0 \right\}.$$

Examples ( $\alpha \in (0, 1)$ ):

$$(t, 1, x) \in \mathcal{K}_{pow}(\alpha, 1 - \alpha) \Leftrightarrow t \geq |x|^{1/\alpha}, t \geq 0,$$

$$(x, 1, t) \in \mathcal{K}_{pow}(\alpha, 1 - \alpha) \Leftrightarrow x^\alpha \geq |t|, x \geq 0,$$

$$(x, t) \in \mathcal{K}_{pow}(e/n) \Leftrightarrow \left( \prod_{j=1}^n x_j \right)^{1/n} \geq |t|, x \geq 0.$$

See also Chares [2] and the Mosek modelling cookbook [6].



## The exponential cone

$$\mathcal{K}_{\text{exp}} := \{(x_1, x_2, x_3) : x_1 \geq x_2 e^{\frac{x_3}{x_2}}, x_2 \geq 0\} \\ \cup \{(x_1, x_2, x_3) : x_1 \geq 0, x_2 = 0, x_3 \leq 0\}$$

## Applications:

$$\begin{aligned} (t, 1, x) \in \mathcal{K}_{\text{exp}} &\Leftrightarrow t \geq e^x, \\ (t, 1, \ln(a)x) \in \mathcal{K}_{\text{exp}} &\Leftrightarrow t \geq a^x, \\ (x, 1, t) \in \mathcal{K}_{\text{exp}} &\Leftrightarrow t \leq \ln(x), \\ (1, x, t) \in \mathcal{K}_{\text{exp}} &\Leftrightarrow t \leq -x \ln(x), \\ (y, x, -t) \in \mathcal{K}_{\text{exp}} &\Leftrightarrow t \geq x \ln(x/y), \text{ (relative entropy)}. \end{aligned}$$

Geometric programming + many more [2, 6].

## Section 2

A primal-dual interior-point algorithm







- Lesson learned from the linear case: Solve the primal and dual problem simultaneously.
- Symmetric cones: Employ the Nesterov-Todd (NT) algorithm [11, 13].
- Nonsymmetric cones: How to generalize the NT algorithm?
  - Nesterov [9, 10], Skajaa and Ye [16], Serrano [15]: Computational results available.
  - Tuncel [17], Myklebust [7], Tuncel and Myklebust [8]: No computational results.

Present work:

- Follows Myklebust and Tuncel.





Primal problem

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{st} & Ax = b, \\ & x \in \mathcal{K} \end{array}$$

and the dual

$$\begin{array}{ll} \text{maximize} & b^T y \\ \text{st} & A^T y + s = c, \\ & s \in (\mathcal{K})^* \end{array}$$

where

$$\mathcal{K} = \mathcal{K}^1 \times \mathcal{K}^2 \times \dots \times \mathcal{K}^k$$

and

$$\mathcal{K}^*$$

is the corresponding dual cone. Known for the 5 cone types.



Define a 3 times differentiable function  $F$  such that

$$F : \text{int}(K) \mapsto \mathcal{R}$$

then it is a  $\nu$ -logarithmically homogeneous self-concordant barrier ( $\nu$ -LHSCB) for  $\text{int}(K)$  if

$$|F'''(x)[u, u, u]| \leq 2(F''(x)[u, u])^{3/2}$$

and

$$F(\tau x) = F(x) - \nu \log \tau.$$

See [11, 13].





If  $F$  is a  $\nu$ -self-concordant barrier for  $K$ , then the Fenchel conjugate

$$F_*(s) = \sup_{x \in \text{int}(K)} \{-s^T x - F(x)\}. \quad (1)$$

is a  $\nu$ -self-concordant barrier for  $K^*$ .

Let

$$\tilde{x} := -F'_*(s), \quad \tilde{s} := -F'(x), \quad \mu := \frac{x^T s}{\nu}, \quad \tilde{\mu} := \frac{\tilde{x}^T \tilde{s}}{\nu}.$$

Then  $\tilde{x} \in \text{int}(K)$ ,  $\tilde{s} \in \text{int}(K^*)$  and

$$\mu \tilde{\mu} \geq 1 \quad (2)$$

with equality iff  $x = -\mu \tilde{x}$  (and  $s = -\mu \tilde{s}$ ).



- An explicit formula for the above barrier may not be known.
- The dual barrier is not unique in general e.g. the primal power cone.



Generalized Goldman-Tucker homogeneous model:

$$\begin{aligned}(H) \quad & Ax - b\tau = 0, \\ & A^T y + s - c\tau = 0, \\ & -c^T x + b^T y - \kappa = 0, \\ & (x; \tau) \in \bar{\mathcal{K}}, (s; \kappa) \in \bar{\mathcal{K}}^*\end{aligned}$$

where

$$\bar{\mathcal{K}} := \mathcal{K} \times \mathcal{R}_+ \quad \text{and} \quad \bar{\mathcal{K}}^* := \mathcal{K}^* \times \mathcal{R}_+.$$

- $\mathcal{K}$  is Cartesian product of  $k + 1$  convex cones.
- The homogeneous model always has a solution.
- Partial list of references:
  - Linear case: [5], [4], [18].
  - Nonlinear case: [12].



## Lemma

Let  $(x^*, \tau^*, y^*, s^*, \kappa^*)$  be any feasible solution to (H), then

i)

$$(x^*)^T s^* + \tau^* \kappa^* = 0.$$

ii) If  $\tau^* > 0$ , then

$$(x^*, y^*, s^*)/\tau^*$$

is an optimal solution.

iii) If  $\kappa^* > 0$ , then at least one of the strict inequalities

$$b^T y^* > 0 \quad (3)$$

and

$$c^T x^* < 0 \quad (4)$$

holds. If the first inequality holds, then (P) is infeasible. If the second inequality holds, then (D) is infeasible.

Remaining case is the illposed case!





The central path:

$$\begin{aligned}Ax - b\tau &= \gamma(A\hat{x} - b\hat{\tau}), \\ A^T y + s - c\tau &= \gamma(A^T \hat{y} + \hat{s} - c\hat{\tau}), \\ -c^T x + b^T y - \kappa &= \gamma(-c^T \hat{x} + b^T \hat{y} - \hat{\kappa}), \\ s + \gamma \hat{\mu} F'(x) &= 0, \\ \tau \kappa - \gamma \hat{\mu} &= 0,\end{aligned}$$

where

$$\hat{\mu} := \frac{(\hat{x})^T \hat{s} + \hat{\tau} \hat{\kappa}}{\nu + 1}$$

and  $(\hat{x}, \hat{\tau}, \hat{y}, \hat{s}, \hat{\kappa})$  is an “interior” solution for  $\gamma = 1$ . The central path is the solutions parameterised by  $\gamma \in [0, 1]$ .





- Idea: Trace the central path using Newton's method.
- Question: Should we use the primal or dual barrier i.e.

$$s - \gamma \hat{\mu} F'(x) = s + \gamma \hat{\mu} \tilde{s} = 0$$

or

$$x - \gamma \hat{\mu} F'_*(s) = x + \gamma \hat{\mu} \tilde{x} = 0$$

where

$$\tilde{x} := -F'_*(s) \text{ and } \tilde{s} := -F'(x).$$



A nonsingular matrix  $W$  is called a primal-dual scaling if it satisfies

$$\begin{aligned}v &:= Wx = W^{-T}s, \\ \tilde{v} &:= W\tilde{x} = W^{-T}\tilde{s}.\end{aligned}$$

Apply the scaling to the primal or the dual centrality conditions:

$$\begin{aligned}W^{-T}(s - \gamma\hat{\mu}F'(x)) &= W^{-T}s + \gamma\hat{\mu}\tilde{s} = v + \gamma\hat{\mu}\tilde{v} = 0, \\ W(x - \gamma\hat{\mu}F'_*(s)) &= Wx + \gamma\hat{\mu}\tilde{x} = v + \gamma\hat{\mu}\tilde{v} = 0.\end{aligned}$$

- Result: The centrality conditions have become symmetric!



Affine direction:

$$\begin{aligned}Ad_x^a - bd_\tau^a &= -(Ax^0 - b\tau^0), \\A^T d_y^a + d_s^a - cd_\tau^a &= -(A^T y^0 + s^0 - c\tau^0), \\-c^T d_x^a + b^T d_y^a - d_\kappa^a &= -(-c^T x^0 + b^T y^0 - \kappa^0), \\Wd_x^a + W^{-T}d_s^a &= -v^0, \\\tau^0 d_\tau^a + \kappa^0 d_\tau^a &= -\tau^0 \kappa^0\end{aligned}$$

where  $v^0 = W^{-T}F'(x^0) = W^T F'_*(s^0)$ .

Centering direction:

$$\begin{aligned}Ad_x^c - bd_\tau^c &= (Ax^0 - b\tau^0), \\A^T d_y^c + d_s^c - cd_\tau^c &= (A^T y^0 + s^0 - c\tau^0), \\-c^T d_x^c + b^T d_y^c - d_\kappa^c &= (-c^T x^0 + b^T y^0 - \kappa^0), \\Wd_x^c + W^{-T}d_s^c &= \mu^0 \tilde{v}^0, \\\tau^0 d_\tau^c + \kappa^0 d_\tau^c &= \mu^0.\end{aligned}$$



For a given  $\gamma \in [0, 1]$  then define

$$\begin{aligned}d_x &:= d_x^a + \gamma d_x^c, \\d_\tau &:= d_\tau^a + \gamma d_\tau^c, \\d_y &:= d_y^a + \gamma d_y^c, \\d_s &:= d_s^a + \gamma d_s^c, \\d_\kappa &:= d_\kappa^a + \gamma d_\kappa^c,\end{aligned}$$

and hence for a step size  $\alpha \in [0, 1]$  we have

$$\begin{aligned}x^+ &:= x^0 + \alpha d_x, \\ \tau^+ &:= \tau^0 + \alpha d_\tau, \\ y^+ &:= y^0 + \alpha d_y, \\ s^+ &:= s^0 + \alpha d_s, \\ \kappa^+ &:= \kappa^0 + \alpha d_\kappa.\end{aligned}$$



$$\begin{aligned}Ax^+ - b\tau^+ &= (1 - \alpha(1 - \gamma))(Ax^0 - b\tau^0), \\A^T y^+ + s^+ - c\tau^+ &= (1 - \alpha(1 - \gamma))(A^T y^0 + s^0 - c\tau^0), \\-c^T x^+ + b^T y^+ - \kappa^+ &= (1 - \alpha(1 - \gamma))(-c^T x^0 + b^T y^0 - \kappa^0), \\(x^+)^T (s^+) + \tau^+ \kappa^+ &= (1 - \alpha(1 - \gamma))((x^0)^T s^0 + \tau^0 \kappa^0).\end{aligned}$$

- **Equal** decrease in infeasibility and complementarity for  $\gamma \in [0, 1)$ .
- If  $\alpha \in ]0, 1]$ , then “convergence”.
- No merit function is needed. Yahoooooo!
- The other primal-dual alg. for nonsymmetric conic optimization lacks this property.



Our method inspired by (Tuncel, Tuncel and Myklebust):

$$\begin{aligned}W^T W &\approx \mu F''(x), \\Wx &= W^{-T} s, \\W\tilde{x} &= W^{-T} \tilde{s}.\end{aligned}$$

- Employ the quasi Newton idea to compute  $W$ . ( $W^T W$  is the variable).
- Has two secant conditions.



## Theorem (Schnabel [14])

Let  $\bar{X}, \bar{S} \in \mathcal{R}^{n \times p}$  have full rank  $p$ . Then there exists  $H \succ 0$  such that  $H\bar{X} = \bar{S}$  if and only if  $\bar{S}^T \bar{X} \succ 0$ .

As a consequence

$$H = \bar{S}(\bar{S}^T \bar{X})^{-1} \bar{S}^T + ZZ^T$$

where  $\bar{X}^T Z = 0$ ,  $\text{rank}(Z) = n - p$ . We have  $n = 3$ ,  $p = 2$  and

$$\bar{X} := (x \quad \tilde{x}), \quad \bar{S} := (s \quad \tilde{s}),$$

with

$$\det(\bar{S}^T \bar{X}) = \nu^2(\mu\tilde{\mu} - 1) \geq 0$$

vanishing only on the central path.



Any scaling with  $n = 3$  satisfies

$$W^T W = \bar{S}(\bar{S}^T \bar{X})^{-1} \bar{S}^T + z z^T$$

where  $(x \quad \tilde{x})^T z = 0$ ,  $z \neq 0$ . Expanding the BFGS update [14]

$$H^+ = H + \bar{S}(\bar{S}^T \bar{X})^{-1} \bar{S}^T - H \bar{X}(\bar{X}^T H \bar{X})^{-1} \bar{X}^T H,$$

for  $H \succ 0$  gives the scaling by Tunçel [17] and Myklebust [8], i.e.,

$$z z^T = H - H \bar{X}(\bar{X}^T H \bar{X})^{-1} \bar{X}^T H,$$

with  $H = \mu F''(x)$ .





A high-order correction:

$$\begin{aligned}Ad_x^{co} - bd_\tau^{co} &= 0, \\A^T d_y^{co} + d_s^{co} - cd_\tau^{co} &= 0, \\-c^T d_x^{co} + b^T d_y^{co} - d_\kappa^{co} &= 0, \\Wd_x^{co} + W^{-T}d_s^{co} &= -\frac{1}{2}W^{-T}F'''(x)[d_x^a, F''(x)^{-1}d_s^a], \\ \tau^0 d_\kappa^{co} + \kappa^0 d_\tau^{co} &= -d_\tau^a d_\kappa^a.\end{aligned}$$

For motivation see [3].

Finally

$$\begin{aligned}d_x &:= d_x^a + \gamma d_x^c + d_x^{co}, \\d_\tau &:= d_\tau^a + \gamma d_\tau^c + d_\tau^{co}, \\d_y &:= d_y^a + \gamma d_y^c + d_y^{co}, \\d_s &:= d_s^a + \gamma d_s^c + d_s^{co}, \\d_\kappa &:= d_\kappa^a + \gamma d_\kappa^c + d_\kappa^{co}.\end{aligned}$$



A 3-self-concordant barrier for the 3 dimensional primal power cone:

$$F(x) = -\log(x_1^{2\alpha} x_2^{2-2\alpha} - x_3^2) - (1 - \alpha) \log x_1 - \alpha \log x_2. \quad (5)$$

suggest by Chares [2]. Generalized in [1].

Is self-dual using redefined inner product.

However,

- The conjugate barrier  $F_*(x)$  or its derivatives cannot be evaluated on closed-form.
- Can be evaluated numerically to high accuracy based of an idea of Nesterov [10].



A 3-self-concordant barrier for the primal exponential cone:

$$F(x) = -\log(x_2 \log(x_1/x_2) - x_3) - \log x_1 - \log x_2. \quad (6)$$

The dual exponential cone is

$$\mathcal{K}_e^* = \text{cl}\{z \in \mathcal{R}^3 \mid e \cdot z_1 \geq -z_3 \exp(z_2/z_3), z_1 > 0, z_3 < 0\}. \quad (7)$$

However,

- The conjugate barrier  $F_*(x)$  or its derivatives cannot be evaluated on closed-form.
- Can be evaluated numerically to high accuracy based of an idea of Akle [15]. (The Lambert function plays a role.)



Easy to compute  $(\hat{x}, \hat{s})$  such that

$$\hat{s} + F'(\hat{x}) = 0.$$

Next compute

$$\rho = \sqrt{\min((1.0 + \|c\|)(1.0 + \|b\|), 1e6)}$$

then let

$$\begin{aligned}x^0 &= \rho \hat{x}, \\ \tau^0 &= \rho, \\ y^0 &= 0, \\ s^0 &= \rho \hat{s}, \\ \kappa^0 &= \rho.\end{aligned}$$



- First compute the affine direction and then compute  $\gamma$ .
- Second compute final direction with centering and corrector.
- Reuse matrix factorization of the Newton equations for the two solves.



Define the one sided neighborhood

$$\mathcal{N}(\beta) = \left\{ (x, s, \tau, \kappa) \mid \tau\kappa \geq \beta\mu, \nu_i \langle F'(x_i), F'_*(s_i) \rangle^{-1} \geq \beta\mu, \forall i \right\},$$

for some  $\beta \in ]0, 1]$ . This bounds

$$\mu\tilde{\mu} \geq 1$$

from above.

We use bisection to compute the largest possible step size  $\alpha \in ]0, 1]$

so

- New point is feasible
- and is in the neighborhood.



Let

$$(x^k, y^k, s^k, \tau^k, \kappa^k)$$

the  $k$ 'th interior-point iterate to the homogeneous model. Since,

$$x^k \in K, s^k \in K^* \text{ and } \tau^k, \kappa^k > 0$$

then compute

$$\rho_p^k = \min \left\{ \rho \mid \left\| A \frac{x^k}{\tau^k} - b \right\|_{\infty} \leq \rho \varepsilon_p (1 + \|b\|_{\infty}) \right\},$$

$$\rho_d^k = \min \left\{ \rho \mid \left\| A^T \frac{y^k}{\tau^k} + \frac{s^k}{\tau^k} - c \right\|_{\infty} \leq \rho \varepsilon_d (1 + \|c\|_{\infty}) \right\},$$

$$\rho_g^k = \min \left\{ \rho \mid \min \left( \frac{(x^k)^T s^k}{(\tau^k)^2}, \left| \frac{c^T x^k}{\tau^k} - \frac{b^T y^k}{\tau^k} \right| \right) \leq \rho \varepsilon_g \max \left( 1, \frac{\min(|c^T x^k|, |b^T y^k|)}{\tau^k} \right) \right\},$$

$$\rho_{pi}^k = \min \left\{ \rho \mid \left\| A^T y^k + s^k \right\|_\infty \leq \rho \varepsilon_i b^T y^k, b^T y^k > 0 \right\},$$

$$\rho_{di}^k = \min \left\{ \rho \mid \left\| Ax^k \right\|_\infty \leq -\rho \varepsilon_i c^T x^k, c^T x^k < 0 \right\}, \text{ and}$$

$$\rho_{ip}^k = \min \left\{ \rho \mid \left\| \begin{array}{c} c^T x^k \\ b^T y^k \\ A^T y^k + s^k \\ Ax^k \end{array} \right\|_\infty \leq \rho \varepsilon_i \left\| \begin{array}{c} y^k \\ s^k \\ x^k \end{array} \right\|_\infty, \left\| \begin{array}{c} y^k \\ s^k \\ x^k \end{array} \right\|_\infty > 0 \right\}.$$





Note

$$\varepsilon_p, \varepsilon_d, \varepsilon_g, \varepsilon_i$$

are nonnegative user specified tolerances.

Observe if for instance

$$\max(\rho_p^k, \rho_d^k, \rho_g^k) \leq 1$$

then

$$\begin{aligned} \left\| A \frac{x^k}{\tau^k} - b \right\|_{\infty} &\leq \varepsilon_p (1 + \|b\|_{\infty}), \\ \left\| A^T \frac{y^k}{\tau^k} + \frac{s^k}{\tau^k} - c \right\|_{\infty} &\leq \varepsilon_d (1 + \|c\|_{\infty}), \\ \min \left( \frac{(x^k)^T s^k}{(\tau^k)^2}, \left| \frac{c^T x^k}{\tau^k} - \frac{b^T y^k}{\tau^k} \right| \right) &\leq \varepsilon_g \max \left( 1, \frac{\min(|c^T x^k|, |b^T y^k|)}{\tau^k} \right) \end{aligned}$$

and hence  $(x^k, y^k, s^k)/\tau^k$  is an almost primal and dual feasible solution with small complementarity gap.



If

$$\rho p_i \leq 1$$

then

$$\left\| A^T y^k + s^k \right\|_{\infty} \leq \varepsilon_i b^T y^k, \quad b^T y^k > 0$$

and define

$$\bar{y} := \frac{y^k}{b^T y^k} \quad \text{and} \quad \bar{s} := \frac{s^k}{b^T y^k}$$

and hence

$$\begin{aligned} b^T \bar{y} &\geq 1, \\ \left\| A^T \bar{y} + \bar{s} \right\| &\leq \varepsilon_i, \\ \bar{s} &\in K^* \end{aligned}$$

i.e. an approximate certificate of primal infeasibility has been computed.



If

$$\rho_{di} \leq 1$$

then an approximate certificate of dual infeasibility has been computed.





If

$$\rho_{ip} \leq 1$$

then

$$\left\| \begin{array}{c} c^T x^k \\ b^T y^k \\ A^T y^k + s^k \\ Ax^k \end{array} \right\|_{\infty} \leq \varepsilon_i \left\| \begin{array}{c} y^k \\ s^k \\ x^k \end{array} \right\|_{\infty}, \quad \left\| \begin{array}{c} y^k \\ s^k \\ x^k \end{array} \right\|_{\infty} > 0$$

i.e. an approximate certificate of ill-posedness.

We have a facial reduction certificate for the primal and/or the dual problem!



The algorithm has been integrated in Mosek v9.

- Problems is dualized if believed beneficial for lin. alg.
- A presolve is performed.
- A scaling is performed.
- Symmetric cones are handled with NT scaling.
- Simple bounds and fixed variables are handled efficiently in the linear algebra.
- Employs highly tuned linear algebra to solve the Newton equation system.

## Section 3

### Computational results





- Hardware: AMD EPYC 7402P 2.80GHz, 24C/48T, 128M Cache (180W) (virum).
- MOSEK: Version 9.2.21.
- Threads: 8 threads is used in test to simulate a typical user environment.
- All timing results  $t$  are in wall clock seconds.
- Test problems: Public (e.g `cbplib.zib.de`) and customer supplied.

# Exponential/power cone optimization



## Optimized problems

Name	# con.	# cone	# var.	# mat. var.
C_Table5_POW_T12	1086456	1328602	3991646	0
WassersteinReg_2img	972	1229312	3687936	0
dump_Prostate_VMAT_308	26872	14649	126320	0
logistic_mip-8-1000-bayes	326	237	1007	0
c-diaz_test_c47	164404	160000	519810	0
x20565-3over2pow	7932	2482	20320	0
z17926	2	75000	225000	0
task_dopt16	1600	26	376	2
entolib_ento2	26	4695	14085	0
entolib_a_56	37	9702	29106	0
exp-ml-scaled-20000	19999	20000	79998	0
exp-ml-20000	19999	20000	79998	0
patil3_conv	418681	413547	1264340	0
c-diaz_test_c47	164404	160000	519810	0
utkarsh_robust_29012019	1228800	819201	2867301	0
varun_conv	333	328	1009	0
z19841	160767	160766	483856	0
z19502	2354679	524286	6546535	0
udomsak	97653	97653	294519	0
relentr25000	1	25000	75000	0
cbf_mra02	3739	3620	11105	0
log-utility-200-5000	10003	5001	25206	0
cbf_cx02-100	5247	5149	15645	0
elmore_delay_16_conv	830	762	2375	0
gp_dave_3_conv	568	373	2052	0
fsparc_6_075_10	942	432	1806	0
c-260209-1	2238	1424	10079	0



# Exponential/power cone optimization



## Result

Name	P. obj.	# sig. fig.	# iter	time(s)
C_Table5_POW_T12	5.3368368447e-02	7	29	172.0
WassersteinReg_2img	-1.9199622575e+01	6	76	255.7
dump_Prostate_VMAT_308	5.6828282708e+04	8	58	752.9
logistic_mip-8-1000-bayes	6.6082723985e+01	9	42	0.1
c-diaz_test_c47	1.8880303425e-02	9	69	78.2
x20565-3over2pow	2.3137374801e+02	8	34	1.7
z17926	1.7942114183e-01	9	45	7.7
task_dopt16	1.3214504598e+01	9	12	0.6
entolib_ento2	-1.1354764143e+01	9	31	0.3
entolib_a_56	-8.2834853287e+00	8	89	4.0
exp-ml-scaled-20000	-3.3125859649e+00	9	139	11.0
exp-ml-20000	-1.9795438202e+04	11	110	4.6
patil3_conv	-1.0539216061e+00	9	88	115.7
c-diaz_test_c47	1.8880303425e-02	9	69	78.2
utkarsh_robust_29012019	1.6172995191e+00	7	83	699.4
varun_conv	-2.3527295782e+01	9	50	0.2
z19841	-2.6100556412e+00	9	85	275.6
z19502	5.1527196039e+06	10	68	398.2
udomsak	7.6466584697e-02	10	206	616.5
relntr25000	6.3511190583e-02	9	22	1.3
cbf_mra02	4.3179836838e+00	9	170	3.4
log-utility-200-5000	-1.8520357724e+03	9	27	1.5
cbf_cx02-100	7.7292742788e+00	8	21	0.7
elmore_delay_16_conv	4.6571224035e+00	8	26	0.2
gp_dave_3_conv	6.1849199940e+00	9	27	0.1
fsparc_6_075_10	4.6989895076e+02	9	20	0.0
c-260209-1	-6.8419351593e-02	9	38	4.1

## Section 4

### Summary





- Presented a primal-dual algorithm for nonsymmetric conic optimization.
- Makes it easy to extend the Nesterov-Todd algorithm to nonsymmetric cones.
- No polynomial complexity proof is provided.
- Good computational results are presented(only 3 dimensional though!).
- Possible improvements:
  - Handle more cone types.
  - Better primal-dual scaling.
  - High accuracy computations in scaling matrix computations.
  - Achieve fast convergence.
  - A multiple corrector.
  - Improve presolve using facial reductions ideas.



- [1] Scott Roy and Lin Xiao.  
On self-concordant barriers for generalized power cones.  
Technical report, Microsoft, 2018.
- [2] Peter Robert Chares.  
*Cones and interior-point algorithms for structured convex optimization involving powers and exponentials.*  
PhD thesis, Ecole polytechnique de Louvain, Universitet catholique de Louvain, 2009.
- [3] J. Dahl and E. D. Andersen.  
A primal-dual interior-point algorithm for nonsymmetric exponential-cone optimization.  
*Math. Programming*, 2021.



- [4] A. J. Goldman and A. W. Tucker.  
Polyhedral convex cones.  
In H. W. Kuhn and A. W. Tucker, editors, *Linear Inequalities and related Systems*, pages 19–40, Princeton, New Jersey, 1956. Princeton University Press.
- [5] A. J. Goldman and A. W. Tucker.  
Theory of linear programming.  
In H. W. Kuhn and A. W. Tucker, editors, *Linear Inequalities and related Systems*, pages 53–97, Princeton, New Jersey, 1956. Princeton University Press.

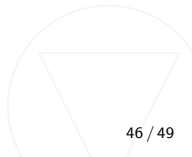




- [6] MOSEK ApS.  
*MOSEK Modeling Cookbook*.  
MOSEK ApS, Fruebjergvej 3, Boks 16, 2100 Copenhagen O,  
2012.  
Last revised September 2018.
- [7] T. Myklebust.  
*On primal-dual interior-point algorithms for convex  
optimisation*.  
PhD thesis, University of Waterloo, 2015.
- [8] T. Myklebust and L. Tunçel.  
Interior-point algorithms for convex optimization based on  
primal-dual metrics.  
Technical report, University of Waterloo, 2014.



- [9] Yu. Nesterov.  
Constructing self-concordant barriers for convex cones.  
Technical report, CORE, Lovain-la-Neuve, 2006.  
Discussion paper 2006/30.
- [10] Yu. Nesterov.  
Towards nonsymmetric conic optimization.  
Technical report, CORE, Lovain-la-Neuve, 2006.  
Discussion paper 2006/38.
- [11] Yu. Nesterov and M. J. Todd.  
Self-scaled barriers and interior-point methods for convex programming.  
*Math. Oper. Res.*, 22(1):1–42, February 1997.





- [12] Yu. Nesterov, M. J. Todd, and Y. Ye.  
Infeasible-start primal-dual methods and infeasibility detectors  
for nonlinear programming problems.  
*Math. Programming*, 84(2):227–267, February 1999.
- [13] Yu. E Nesterov and M. J. Todd.  
Primal-dual interior-point methods for self-scaled cones.  
*SIAM J. on Optim.*, 8:324–364, 1998.
- [14] R. B. Schnabel.  
Quasi-newton methods using multiple secant equations.  
Technical report, Colorado Univ., Boulder, Dept. Comp. Sci.,  
1983.





- [15] Santiago Akle Serrano.  
*Algorithms for unsymmetric cone optimization and an implementation for problems with the exponential cone.*  
PhD thesis, Stanford University, 2015.
- [16] Anders Skajaa and Yinye Ye.  
A homogeneous interior-point algorithm for nonsymmetric convex conic optimization.  
*Math. Programming*, 150:391–422, May 2015.
- [17] L. Tunçel.  
Generalization of primal-dual interior-point methods to convex optimization problems in conic form.  
*Foundations of Computational Mathematics*, 1:229–254, 2001.



- [18] Y. Ye, M. J. Todd, and S. Mizuno.  
An  $O(\sqrt{n}L)$  - iteration homogeneous and self-dual linear programming algorithm.  
*Math. Oper. Res.*, 19:53–67, 1994.

