

# Rigidity theory for Gaussian graphical models: the maximum likelihood threshold of a graph

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<https://dibernstein.github.io>

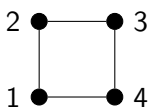
- Broad motivating problem: fit model to relatively small dataset (fewer observations than variables)
- Gaussian graphical models: family of multivariate normal distributions satisfying independence constraints given by a graph
- Goal: use combinatorics of the graph to determine how few observations are needed to be able to fit the graphical model
- Take-home message: rigidity theory offers many tools

# Gaussian graphical models

Let  $\mu \in \mathbb{R}^v$  and  $\Sigma \in \mathbb{R}^{v \times v}$  be positive definite. The ***multivariate normal distribution***  $\mathcal{N}(\mu, \Sigma)$  **with mean  $\mu$  and covariance  $\Sigma$**  has density

$$f_{\mu, \Sigma}(x) := \frac{\exp(-\frac{1}{2}(x - \mu)^T \Sigma^{-1}(x - \mu))}{\sqrt{(2\pi)^v \det(\Sigma)}}$$

Given a graph  $G = (V, E)$ , the ***Gaussian graphical model***  $\mathcal{M}_G$  consists of all multivariate normal distributions  $\mathcal{N}(\mu, \Sigma)$ , with random variables  $V$ , such that  $(\Sigma)_{uv}^{-1} = 0$  whenever  $uv$  is ***not*** an edge of  $G$ .


$$\Sigma^{-1} = \begin{pmatrix} x_{11} & x_{12} & 0 & x_{14} \\ x_{12} & x_{22} & x_{23} & 0 \\ 0 & x_{23} & x_{33} & x_{34} \\ x_{14} & 0 & x_{34} & x_{44} \end{pmatrix}$$

Interpretation:  $uv \notin E$  means  $u \perp\!\!\!\perp v \mid V \setminus \{u, v\}$  for distributions in  $\mathcal{M}_G$

# Maximum likelihood estimation: definition

Suppose we are given:

- A graph  $G = (V, E)$ , and
- datapoints  $x_1, \dots, x_n$ , supposedly iid from some distribution in  $\mathcal{M}_G$ .

The **maximum likelihood estimate (MLE)** is the solution to the following optimization problem, if it exists:

$$\begin{aligned} \max_{\mu, \Sigma} \quad & \prod_{i=1}^n f_{\mu, \Sigma}(x_i) \\ \text{s.t.} \quad & (\Sigma^{-1})_{uv} = 0 \text{ for all } uv \notin E, \\ & \Sigma \succ 0 \end{aligned}$$

This can be found via convex optimization.

# Maximum likelihood estimation: convex optimization

Let  $\hat{\mu}$  and  $S$  be the sample mean and covariance (note:  $\text{rank}(S) = n$  a.s.)

$$\hat{\mu} := \frac{1}{n} \sum_{i=1}^n x_i \quad S := \sum_{i=1}^n (x_i - \hat{\mu})^T (x_i - \hat{\mu}).$$

The MLE in  $\mathcal{M}_G$  exists if and only if the following can be solved:

$$\begin{aligned} \max_K \quad & \text{Tr}(SK) + \log \det K \\ \text{s.t.} \quad & K \succ 0 \quad \text{and} \quad K_{uv} = 0 \text{ for all non-edges } uv \text{ of } G. \end{aligned}$$

## Theorem (Dempster 1972)

*The MLE exists iff there exists  $A \succ 0$  satisfying*

$$A_{ij} = S_{ij} \quad \text{if } i = j \text{ or } ij \text{ is an edge of } G.$$

## Definition (Maximum likelihood threshold)

Given  $G = (V, E)$ ,  $\text{MLT}(G)$  is the minimum  $n$  such that the maximum likelihood estimate in  $\mathcal{M}_G$  exists almost surely given  $n$  datapoints.

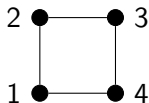
# First bounds on MLT

## Theorem (Dempster 1972)

$MLT(G)$  is the minimum  $r$  such that for almost every  $S \succeq 0$  of rank  $r$ , there exists  $A \succ 0$  such that

$$A_{ij} = S_{ij} \quad \text{if } i = j \text{ or } ij \text{ is an edge of } G.$$

- $MLT(K_n) = n$ .
- If  $G$  has a  $k$ -clique, then  $MLT(G) \geq k$ .



$$\begin{pmatrix} x_{11} & x_{12} & \mathbf{x_{13}} & x_{14} \\ x_{12} & x_{22} & x_{23} & \mathbf{x_{24}} \\ \mathbf{x_{13}} & x_{23} & x_{33} & x_{34} \\ x_{14} & \mathbf{x_{24}} & x_{34} & x_{44} \end{pmatrix}$$

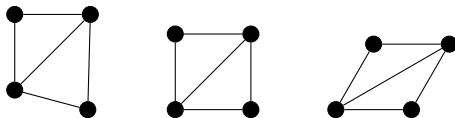
## Theorem (Buhl 1993)

Let  $\omega(G)$  and  $\tau(G)$  denote the clique number and treewidth of  $G$ . Then

$$\omega(G) \leq MLT(G) \leq \tau(G) + 1.$$

## Definition

A **bar and joint framework in  $d$  dimensions** consists of a graph  $G$ , and a map  $p : V(G) \rightarrow \mathbb{R}^d$ . Such a framework is **independent** if the edge-lengths can be independently perturbed.



## Theorem (Asimov and Roth 1978)

Given a graph  $G$ , then if  $p : V(G) \rightarrow \mathbb{R}^d$  is “generic,” then whether the framework  $(G, p)$  is independent in  $\mathbb{R}^d$  does not depend on  $p$ .

One says that  $G$  is (generically) independent in  $\mathbb{R}^d$  if  $(G, p)$  is independent for all generic  $p : V(G) \rightarrow \mathbb{R}^d$ .

# Upper bounds via generic independence

## Definition (Generic completion rank)

The **generic completion rank of  $G$** , denoted  $\text{GCR}(G)$ , is the minimum  $d$  such that  $G$  is generically independent in  $\mathbb{R}^{d-1}$ .

$\text{GCR}(G)$  is also the minimum  $k$  such that every generic partial symmetric matrix whose missing entries correspond to the non-edges of  $G$  can be completed to rank  $k$ .

## Theorem (Uhler 2012, Gross and Sullivant 2018)

$$\text{MLT}(G) \leq \text{GCR}(G).$$

$\text{GCR}(G)$  can be computed in RP time, so it would be great if the above inequality were sharp. However...

## Theorem (Blekherman and Sinn 2019)

$$\text{MLT}(K_{5,5}) = 4 \text{ but } \text{GCR}(K_{5,5}) = 5.$$

$\text{MLT}(K_{n,n})$  grows linearly with  $n$  whereas  $\text{GCR}(K_{n,n})$  grows quadratically.



# MLT in rigidity-theoretic terms

## Definition

Let  $(G, p)$  and  $(G, q)$  be frameworks in  $\mathbb{R}^d$  and  $\mathbb{R}^e$ . Consider the equality

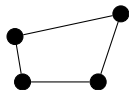
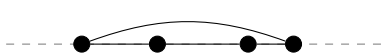
$$\|p(u) - p(v)\| = \|q(u) - q(v)\|.$$

If it holds for all edges  $uv$  of  $G$ , the frameworks are **equivalent**. If it moreover holds for **all pairs** of vertices, the frameworks are **congruent**.

A framework  $(G, p)$  in  $\mathbb{R}^k$  has **full affine span** if  $\{p(v) : v \in V(G)\}$  affinely spans  $\mathbb{R}^k$ .

## Theorem (BDGNST 2021+)

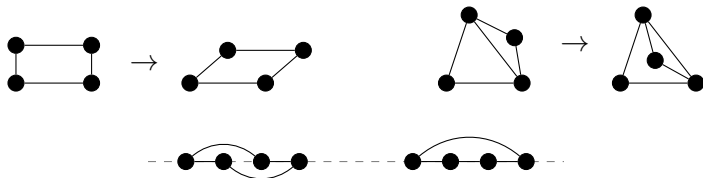
Let  $G$  be a graph with  $n$  vertices.  $\text{MLT}(G)$  is the smallest  $d$  such that every generic framework in  $\mathbb{R}^{d-1}$  is equivalent to a framework in  $\mathbb{R}^{n-1}$  with full affine span.



# More rigidity

One says that  $(G, p)$  is:

- universally rigid if  $(G, p)$  and  $(G, q)$  are congruent when equivalent
- globally rigid if  $(G, p)$  and  $(G, q)$  are congruent when they are equivalent frameworks in the same dimension
- locally rigid if  $(G, p)$  and  $(G, q)$  are congruent when they are equivalent frameworks in the same dimension and sufficiently close



- Local and global rigidity are generic properties (Asimov and Roth 1978; Connelly 2005; Gortler, Healy, and Thurston 2010)
- If  $G$  has an open set of frameworks in  $\mathbb{R}^{d-1}$  that are all universally rigid, then  $\text{MLT}(G) > d$

# Importing results from rigidity

## Theorem (Connelly, Gortler, and Theran 2020)

*$G$  is generically globally rigid in  $\mathbb{R}^{d-1}$  if and only if there exists an open set of configurations on  $G$  in  $\mathbb{R}^{d-1}$  that are all universally rigid.*

## Theorem (BDGNST 2021+)

*If a subgraph of  $G$  on at least  $d + 1$  vertices is generically globally rigid in  $\mathbb{R}^{d-1}$ , then  $\text{MLT}(G) > d$ .*

Implications:

- Lower bounds on MLT generalizing Buhl's result  $\omega(G) \leq \text{MLT}(G)$
- If  $G$  has fewer than 9 vertices, then  $\text{MLT}(G) = \text{GCR}(G)$
- If  $\text{GCR}(G) \leq 4$  or  $\text{MLT}(G) \leq 3$ , then  $\text{MLT}(G) = \text{GCR}(G)$

# Importing results from low-dimensional rigidity

## Proposition (Folklore)

Let  $G$  be a graph with  $n$  vertices. Then

- $G$  is independent in  $\mathbb{R}^1$  iff  $G$  has no cycles
- $G$  is globally rigid in  $\mathbb{R}^1$  iff  $G$  is 2-connected

If  $\text{GCR}(G) = 3$ , then  $\text{MLT}(G) = 3$ :

- $\text{GCR}(G) = 3$  implies  $G$  has a cycle
- cycles are globally rigid in  $\mathbb{R}^1$ , so  $\text{MLT}(G) > 2$
- $\text{MLT}(G) \leq \text{GCR}(G)$ , so  $\text{MLT}(G) = 3$ .

## Theorem (Berg and Jordán 2003)

If  $G$  is 3-connected and minimally dependent in  $\mathbb{R}^2$ , then  $G$  is globally rigid in  $\mathbb{R}^2$ .

## Theorem (BDGNST 2021+)

If  $\text{GCR}(G) = 4$  then  $\text{MLT}(G) = 4$ .

# Weak maximum likelihood thresholds

## Definition (Weak maximum likelihood threshold)

Given a graph  $G$ ,  $\text{WMLT}(G)$  denotes the minimum  $n$  such that the maximum likelihood estimate in  $\mathcal{M}_G$  given  $n$  datapoints exists with positive probability.

## Proposition (Folklore)

$\text{WMLT}(G) = 1$  iff  $G$  has no edges.

## Proposition (BDGNST)

*If  $\text{WMLT}(G) = 2$ , then there exists an orientation of the edges of  $G$  yielding the order diagram of a partially ordered set.*

- Conjecture: the converse is true too
- If the above conjecture holds, then computing  $\text{WMLT}(G)$  is NP-hard

- Find an algorithm for computing  $\text{MLT}(G)$  and  $\text{WMLT}(G)$
- Find new examples of graphs where  $\text{MLT}(G) < \text{GCR}(G)$
- Determine if there exists an efficient algorithm for finding the largest  $d$  such that  $G$  contains a subgraph that is globally rigid in  $\mathbb{R}^d$
- Bound  $\text{MLT}(G)$  in terms of the genus of  $G$  (Dewar 2021+)
- Can **coordinated** rigidity (Schulze, Serocold, and Theran 2018) be used to understand MLTs of **colored** Gaussian graphical models?
- Are there any subfields of rigidity theory that could be used to understand MLTs of **directed** Gaussian graphical models?

Thank you for your attention!

(Preprint coming soon)