

# Mobile Facility Location

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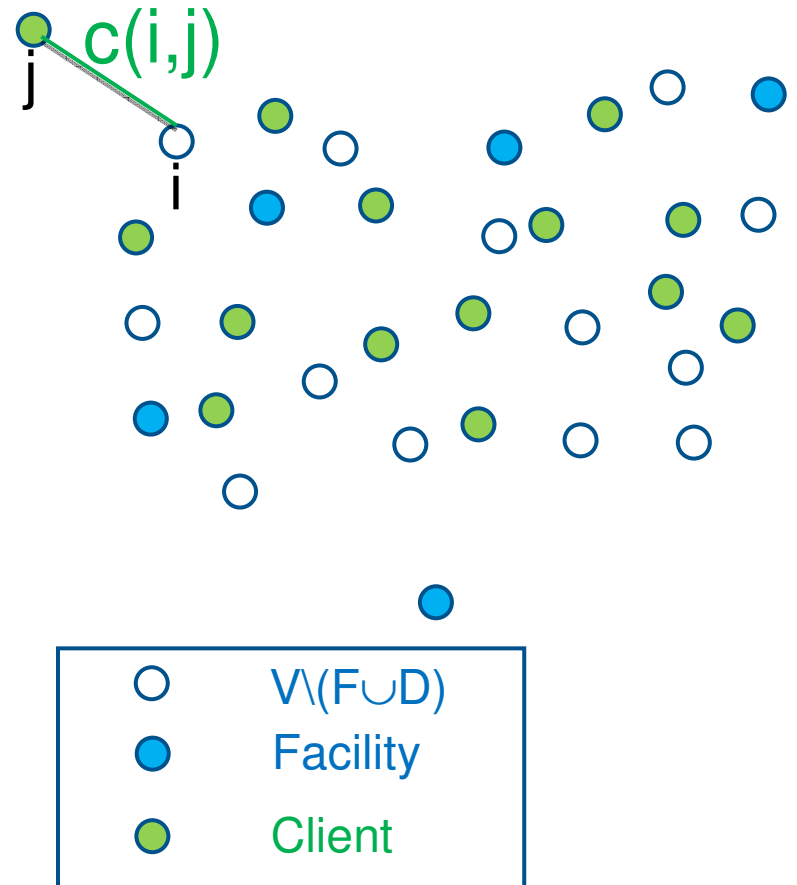
# Mobile Facility Location(MFL)

- ▶ We are given:
  - ▶  $F \subseteq V$ : k initial facility locations.
  - ▶  $D \subseteq V$ : set of clients.

Located in common metric space  $G=(V,c)$ .

- ▶  $c(i,j)$  : cost of moving from point i to point j. We assume that  $c(i,j)$ s form a metric.

Sample Example:



# Goal

- ▶ Find solution  $\mathbf{S}=\{\mathbf{s}_1,\mathbf{s}_2,\dots,\mathbf{s}_k\}$ 
  - ▶ Moves each facility  $i$  to location  $s_i$  in  $S$  incurring movement cost of  $c(i,s_i)$ :

$$\sum_{s_i \in S} c(i, s_i) = \sum_{i \in F} f_i$$

- ▶ Assign each client to the closest location in  $S$ :

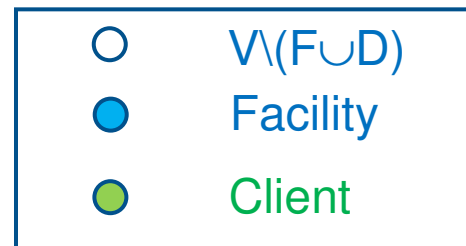
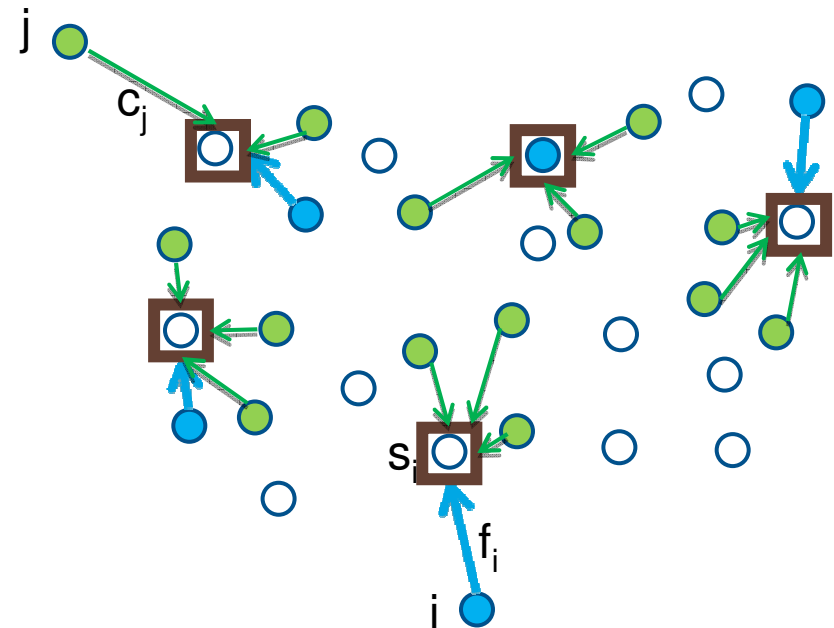
$$\sum_{j \in D} c(j, S) = \sum_{j \in D} c_j$$

- ▶ Goal: minimize total cost

$$\sum_{i \in F} f_i + \sum_{j \in D} c_j$$

$$OPT = \sum_{i \in F} f_i^* + \sum_{j \in D} c_j^*$$

## Sample Example:



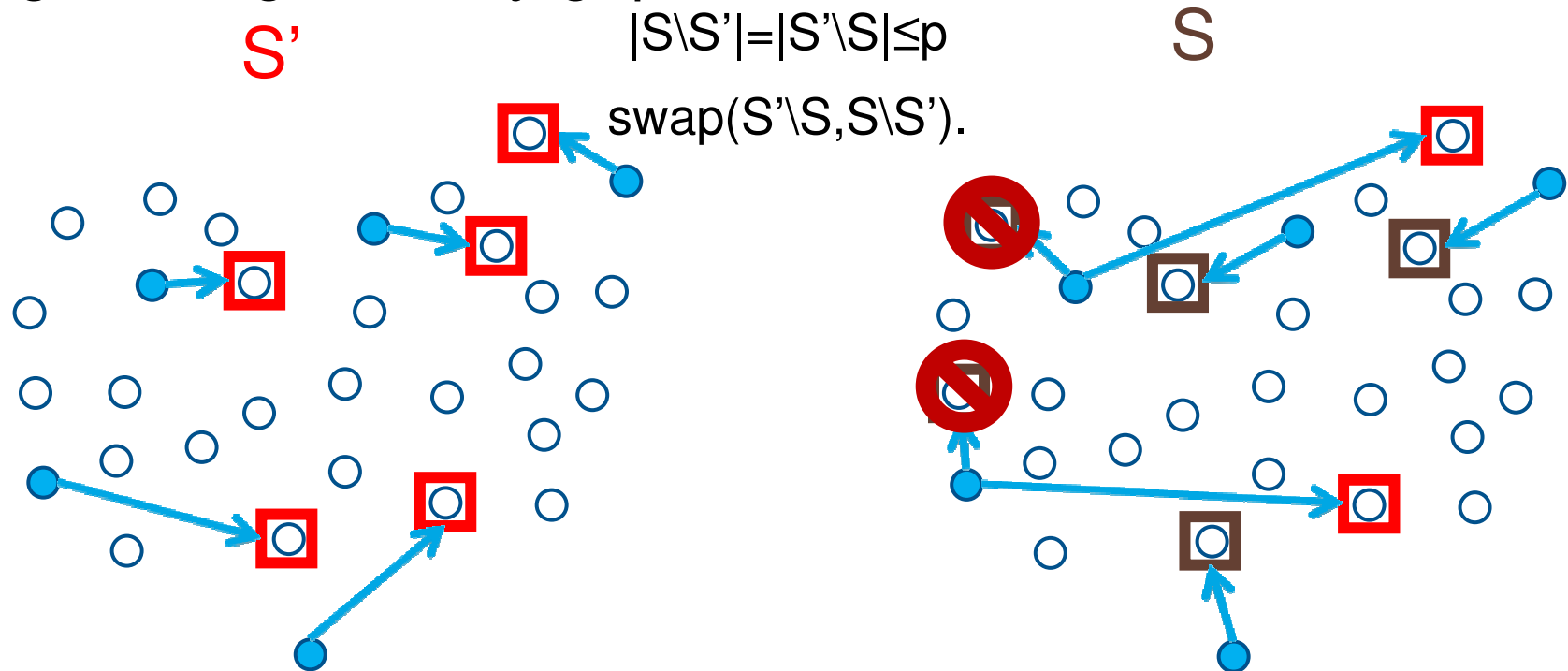
# Our work

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- ▶ We give  $(3+\varepsilon)$ -approximation for any constant  $\varepsilon>0$  based on local-search (first combinatorial algorithm).
  - ▶ Previous best: 8-approximation due to Friggstad and Salavatipour [FS] based on LP-rounding.
- ▶ Extension to weighted generalization: the movement cost for facility  $i$  is  $w_i$  times the distance travelled by  $i$ .

# Local Search

- ▶ At each step, swap in and swap out a fixed number (say  $p$ ) of locations and rematch **all** initial locations to new set of locations. Without rematching, we may get a large locality gap.



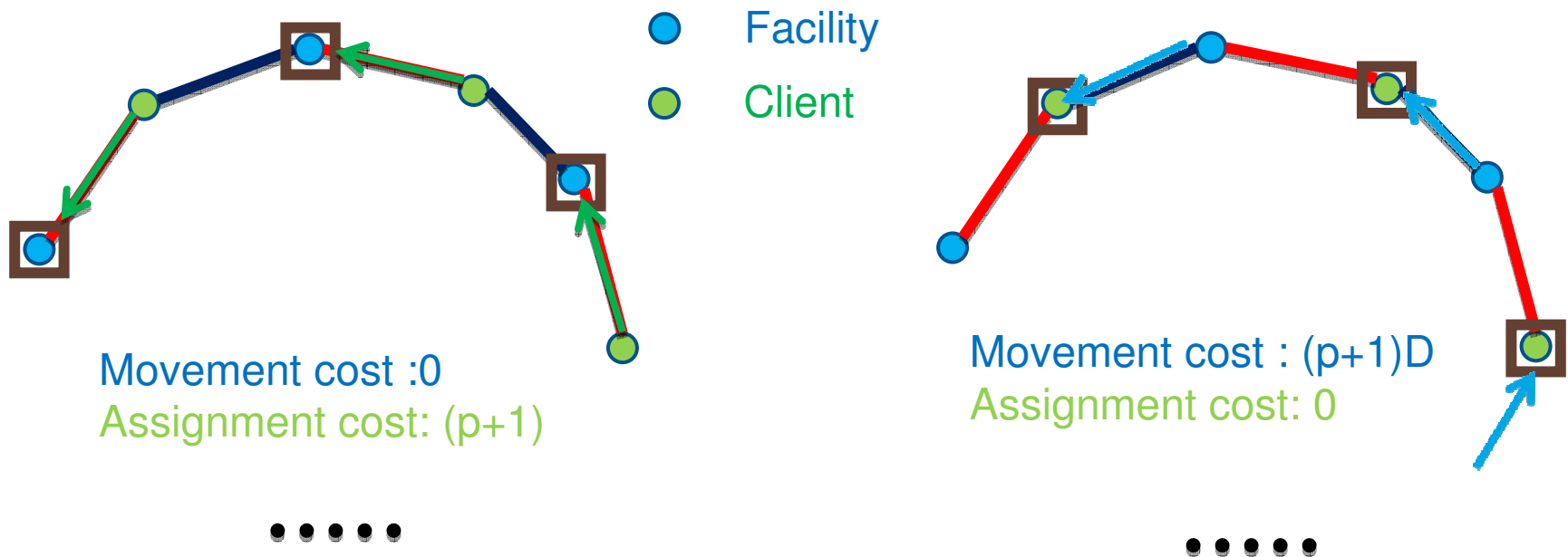
# Analysis

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- ▶ Generate inequalities using test swaps. Each individual inequality may have expensive terms but can amortize the cost by suitably combining the inequalities.
- ▶ Use “recursion” tree to identify a suitable collection of path swaps.

# Different metrics for facilities and clients

- ▶ Local search has big locality gap if facilities and clients move in different metrics.



—  $(1, \infty)$     **Red** edge: assignment cost: 1 and movement cost:  $\infty$ .  
—  $(\infty, D)$     **Blue** edge: assignment cost:  $\infty$  and movement cost:  $D$ .

# Open Questions

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- ▶ Better approximation ratio?
  - ▶  $1+\sqrt{3}$ -approximation for k-median by Shi Li and Ola Svensson.
- ▶ Combinatorial algorithm for matroid median.
- ▶ What if we use approximate matching between initial locations and final locations instead of optimal matching?
- ▶ Reducing approximation ratio for single swap case.



Thank You

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QUESTIONS?