Preprocessing and Reduction for Degenerate Semidefinite Programs

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13th Midwest Optimization Meeting

Motivation: Loss of Slater's CQ

- Primal-dual interior-point methods assume Slater's CQ to hold.
- <u>However</u>, surprisingly many conic opt / SDP relaxations, instances arising from applications (QAP, GP, strengthened MC, SNL, Molecular Conformation) do not satisfy Slater's CQ
- Lack of Slater's CQ results in: unbounded dual solutions; theoretical and numerical difficulties

• Solution:

- theoretical facial reduction (Borwein, Wolkowicz'81[1])
- preprocess for regularized smaller problem (C.,Schurr, Wolkowicz'11[4])
- backward stable when # facial reduction step = 1



- Minimal Faces for Preprocessing
- Theorem of Alternative to Slater's CQ

2 Algorithm

- Backward Stability
- Numerical Results

Minimal Faces for Preprocessing Theorem of Alternative to Slater's CQ

Semidefinite Program, SDP

(P)
$$v_{\mathrm{P}} = \sup_{y \in \mathbb{R}^m} b^\top y$$
 s.t. $\mathcal{A}^* y - C \leq 0$
(D) $v_{\mathrm{D}} = \inf_{X \in \mathbb{S}^n} \langle C, X \rangle$ s.t. $\mathcal{A}(X) = b, X \succeq 0$

where

Minimal Faces for Preprocessing Theorem of Alternative to Slater's CQ

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o primal feasible set:

$$\mathcal{F}_{\mathrm{P}}^{Z} := \{Z \succeq 0 : Z = C - \mathcal{A}^{*}y\}$$

• dual feasible set:

$$\mathcal{F}^{X}_{\mathrm{D}} := \{X \succeq 0 : \mathcal{A}(X) = b\}$$

Minimal Faces for Preprocessing Theorem of Alternative to Slater's CQ

Semidefinite Program, SDP

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Assumptions

- A is onto.
- (P) is feasible; i.e., wlog, $C \succeq 0$.

Minimal Faces for Preprocessing Theorem of Alternative to Slater's CQ

Semidefinite Program, SDP

(P)
$$v_{\mathrm{P}} = \sup_{y \in \mathbb{R}^{m}} b^{\top}y$$
 s.t. $\mathcal{A}^{*}y - C \preceq 0$
(D) $v_{\mathrm{D}} = \inf_{X \in \mathbb{S}^{n}} \langle C, X \rangle$ s.t. $\mathcal{A}(X) = b, X \succeq 0$

Slater's CQ for (P)

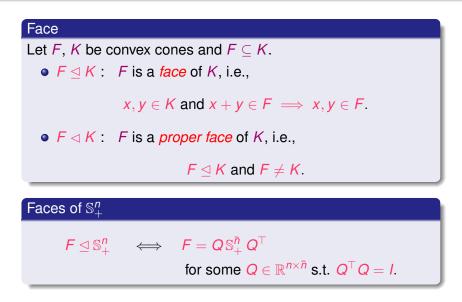
$$\exists \hat{y} \text{ s.t. } Z = C - \mathcal{A}^* \hat{y} \succ 0$$

Strong Duality

If Slater's CQ holds for (P) and opt(P) is bounded above, then opt(P) = opt(D) and opt(D) is attained.
 Theory
 Minimal Faces for Preprocessing

 Algorithm
 Theorem of Alternative to Slater's CQ

Faces of Cones - for Characterization of Optimality



Minimal Faces

Primal and Dual Minimal Faces

$$\begin{split} f_{\mathrm{P}} &:= \mathsf{face}(\mathcal{F}_{\mathrm{P}}^{Z}) \trianglelefteq \mathbb{S}_{+}^{n} \\ f_{\mathrm{D}} &:= \mathsf{face}(\mathcal{F}_{\mathrm{D}}^{X}) \trianglelefteq \mathbb{S}_{+}^{n} \end{split}$$

- $(\mathcal{F}_{P}^{Z} : \text{primal feasible set})$ $(\mathcal{F}_{D}^{X} : \text{dual feasible set})$
- Why Are We Interested in Minimal Faces?
 - Let $Z = C A^* y$. Then

$$Z \in \mathbb{S}^n_+ \iff Z \in f_{\mathrm{P}}$$

• $f_{\mathrm{P}} = \mathbf{Q}_{\mathrm{fin}} \mathbb{S}_{+}^{\bar{n}} \mathbf{Q}_{\mathrm{fin}}^{\top}$ for some $\mathbf{Q}_{\mathrm{fin}} \in \mathbb{R}^{n \times \bar{n}}$ satisfying $\mathbf{Q}_{\mathrm{fin}}^{\top} \mathbf{Q}_{\mathrm{fin}} = I$.

Slater's CQ holds iff

$$f_{\mathrm{P}} = \mathbb{S}^n_+$$
.

Regularizing (P) Using Minimal Face

Borwein-Wolkowicz'81 [1], $f_{\rm P} := \operatorname{face}(\mathcal{F}_{\rm P}^Z)$

(P) is equivalent to the regularized SDP

$$(\mathbf{P}_{\mathrm{reg}}) \qquad \mathbf{v}_{\mathrm{P}} = \mathbf{v}_{\mathrm{RP}} := \sup_{\mathbf{y}} \left\{ \mathbf{b}^{\top} \mathbf{y} : \mathcal{A}^* \mathbf{y} - \mathbf{C} \preceq_{\mathbf{f}_{\mathrm{P}}} \mathbf{0} \right\}$$

Theory

Algorithm

Lagrangian Dual of (P_{reg}) Satisfies Strong Duality:

Let

$$(\mathbf{D}_{\operatorname{reg}}) \quad \mathbf{V}_{\operatorname{DRP}} := \inf_{X} \big\{ \langle C, X \rangle : \mathcal{A}(X) = b, \ X \succeq_{(f_{P})^{*}} \mathbf{0} \big\}.$$

Then

$$V_{\rm DRP} = V_{\rm RP} = V_{\rm P},$$

and *v*_{DRP} is <u>attained</u>.

Regularizing (P) Using Minimal Face

Theory

Algorithm

(P) is equivalent to (P_{reg}), given by

$$(\mathbf{P}_{\mathrm{reg}}) \qquad \sup_{y} \ b^{\top}y \ \text{s.t.} \ g^{\prec}(y) \leq 0, \ g^{=}(y) = 0,$$

where

$$egin{aligned} g^{\prec}(y) &:= & \mathcal{Q}_{ ext{fin}}^{ op}(\mathcal{A}^*y - \mathcal{C})\mathcal{Q}_{ ext{fin}}, \ g^{=}(y) &:= egin{bmatrix} P_{ ext{fin}}^{ op}(\mathcal{A}^*y - \mathcal{C})P_{ ext{fin}}\ P_{ ext{fin}}^{ op}(\mathcal{A}^*y - \mathcal{C})\mathcal{Q}_{ ext{fin}} \end{bmatrix}, \ egin{bmatrix} P_{ ext{fin}} & \mathcal{Q}_{ ext{fin}} \end{bmatrix} \in \mathbb{R}^{n imes n} ext{ is orthogonal.} \end{aligned}$$

Generalized Slater's CQ holds for (P_{reg}) : $\exists \hat{y} \text{ s.t. } g^{\prec}(\hat{y}) \prec 0 \text{ and } g^{=}(\hat{y}) = 0.$

Slater's CQ and Theorem of Alternative

Theory

Algorithm

Slater's CQ for (P)

$$\exists \hat{y} \text{ s.t. } Z = C - \mathcal{A}^* \hat{y} \succ 0$$

Theorem of the Alternative

Assume that $\exists \tilde{y} \text{ s.t. } C - \mathcal{A}^* \tilde{y} \succeq 0$. Then Slater's CQ holds iff

 $\mathcal{A}(D) = 0, \ \langle C, D \rangle = 0, \ D \succeq 0 \implies D = 0.$ (*)

Theorem of Alternative and Primal Minimal Face

Theory

Algorithm

Alternative to Slater's CQ

$$\mathcal{A}(D) = 0, \ \langle C, D \rangle = 0, \ 0 \neq D \succeq 0$$
 (*)

Determining a proper face $f \triangleleft \mathbb{S}^n_+$ containing f_P

$$C - \mathcal{A}^* y \succeq 0 \implies \langle C - \mathcal{A}^* y, D^* \rangle = 0,$$

so $\mathcal{F}_{\mathrm{P}}^{Z} \subseteq \mathbb{S}_{+}^{n} \cap \{D^{*}\}^{\perp}$.

• Wlog suppose rank $(D^*) = n - \bar{n} < n$. Write $D^* = PD_+P^\top$, with $D_+ \succ 0$, and $[P \ Q] \in \mathbb{R}^{n \times n}$ orthogonal.

Then

$$Z = C - \mathcal{A}^* y \succeq 0 \implies \begin{bmatrix} P^\top Z P & P^\top Z Q \\ Q^\top Z P & Q^\top Z Q \end{bmatrix} \in \begin{bmatrix} 0 & 0 \\ 0 & \mathbb{S}_+^{\bar{n}} \end{bmatrix}$$

so $\mathcal{F}_{\mathrm{P}}^{Z} \subseteq \mathcal{Q}\mathbb{S}_{+}^{\bar{n}} \mathcal{Q}^{\top} \lhd \mathbb{S}_{+}^{n}$.

Theorem of Alternative and Primal Minimal Face

Theory

Algorithm

Alternative to Slater's CQ

$$\mathcal{A}(D) = 0, \ \langle C, D \rangle = 0, \ 0 \neq D \succeq 0$$
 (*)

Two Forms of Reduced Problem

Let $D^* = PD_+P^{\top}$ solve (*), with $D_+ \in \mathbb{S}^{n-\bar{n}}_+$ ($\bar{n} > 0$), and [P Q] $\in \mathbb{R}^{n \times n}$ orthogonal. Suppose $\mathcal{R}(Q \cdot Q^{\top}) \cap \mathcal{R}(\mathcal{A}^*)$ is of dim $\bar{m} > 0$. Then (P) is equivalent to $\sup_{y} \left\{ b^{\top}y : Z = C - \mathcal{A}^*y, \ Q^{\top}ZQ \succeq 0, P^{\top}ZP = 0, P^{\top}ZQ = 0 \right\},$ or $\sup_{y} \left\{ b^{\top}(\mathcal{P}v) : \bar{C} - Q^{\top}(\mathcal{A}^*\mathcal{P}v)Q \succeq_{\mathbb{S}^{\bar{n}}_+} 0 \right\},$ where $\mathcal{P} : \mathbb{R}^{\bar{m}} \to \mathbb{R}^m$ is a one-one map satisfying $\mathcal{R}(\mathcal{A}^*\mathcal{P}) = \mathcal{R}(Q \cdot Q^{\top}) \cap \mathcal{R}(\mathcal{A}^*).$

In particular, the linear map $Q^{\top}(\mathcal{A}^*\mathcal{P}(\cdot))Q$ is one-one.

Minimal Faces for Preprocessing Theorem of Alternative to Slater's CQ

Auxiliary Problem

Alternative to Slater's CQ

$$\mathcal{A}_{C}(D) := \begin{pmatrix} \mathcal{A}(D) \\ \langle C, D \rangle \end{pmatrix} = 0, \ 0 \neq D \succeq 0$$
 (*)

How to find a solution D^* of (*)?

• use the auxiliary problem

(AP)
$$\min_{\delta,D} \delta \text{ s.t. } \|\mathcal{A}_{C}(D)\|_{2} \leq \delta,$$
$$\operatorname{trace}(D) = \sqrt{n},$$
$$D \succeq 0.$$

- Both (AP) and its dual satisfy Slater's CQ.
- Suppose (δ*, D*) is an optimal solution to (AP). If δ* = 0, then D* solves (*). If δ* > 0, then Slater's CQ holds for (P).

Theory Minimal Faces for Preprocessing Algorithm Theorem of Alternative to Slater's CQ

Auxiliary Problem and Strict Complementarity

Auxiliary Problem and Reduced Problem

If $(0, PD_+P^{\top})$ (with $D_+ \in \mathbb{S}_+^{n-\bar{n}}$ and $[P \ Q]$ orthogonal) solves

$$AP) \qquad \min_{\delta,D} \ \delta \ \text{ s.t. } \|\mathcal{A}_{\mathcal{C}}(D)\|_{2} \leq \delta,$$

trace(D) =
$$\sqrt{n}$$
,
D > 0.

then

$$\mathbf{V}_{\mathbf{P}} = \sup_{\mathbf{v}} \left\{ b^{\top}(\mathcal{P}\mathbf{v}) : \bar{\mathbf{C}} - \mathbf{Q}^{\top}(\mathcal{A}^*\mathcal{P}\mathbf{v})\mathbf{Q} \succeq_{\mathbb{S}^{\bar{n}}_{+}} \mathbf{0} \right\}.$$

Strict Complementarity of (AP) with soln. $(0, PD_+P^\top)$

Slater's CQ holds for
$$\sup_{v} \left\{ b^{\top}(\mathcal{P}v) : \overline{C} - Q^{\top}(\mathcal{A}^*\mathcal{P}v)Q \succeq_{\mathbb{S}^{\overline{p}}_{+}} 0 \right\}$$

if and only if

(AP) has a strictly complementary optimal p-d solution pair.



Facial Reduction Algorithm

- iteratively reduces (P) to a smaller equivalent problem
- requires at most n 1 iterations

One iteration of facial reduction

```
Input( \mathcal{A} : \mathbb{S}^n \to \mathbb{R}^m, b \in \mathbb{R}^m, C \in \mathbb{S}^n):
Obtain an optimal solution (\delta^*, D^*) of
                 (AP) \min_{\delta | D|} \delta s.t. \|\mathcal{A}_{C}(D)\|_{2} < \delta, trace(D) = \sqrt{n}, D \succ 0;
if \delta^* > 0, then
        STOP; Slater's CQ holds for (A, b, c);
else
        if D^* \succ 0, then
                  STOP; the unique solution y of the equation C = A^* y is optimal;
        else
                 using D^* = \begin{bmatrix} P & Q \end{bmatrix} \begin{bmatrix} D_+ & 0 \\ 0 & 0 \end{bmatrix} \begin{bmatrix} P^\top \\ Q^\top \end{bmatrix}, form the equivalent problem
                                  \sup_{v} \left\{ b^{\top}(\mathcal{P}v) : \overline{C} - Q^{\top}(\mathcal{A}^*\mathcal{P}v)Q \succeq_{\mathbb{S}^{\overline{n}}} 0 \right\};
                  update: \mathcal{A}^* \leftarrow \bar{\mathcal{A}}^* := Q^\top (\mathcal{A}^* \mathcal{P}(\cdot))Q, \quad C \leftarrow \bar{C}, \quad b \leftarrow \bar{b} := \mathcal{P}^* b.
        end if
end if
```

Theory Backward Stability Algorithm Numerical Results

Nearby Solutions for Reduced Problem

Assumptions

Let
$$(\delta^*, D^*)$$
 solve (AP),
 $D^* = \begin{bmatrix} P & Q \end{bmatrix} \begin{bmatrix} D_+ & 0 \\ 0 & D_e \end{bmatrix} \begin{bmatrix} P^\top \\ Q^\top \end{bmatrix}$ with $D_+ \succ 0$,
and $\mathcal{R}(Q \cdot Q^\top) \cap \mathcal{R}(\mathcal{A}^*) \neq \{0\}$.

Theory Backward Stability Algorithm Numerical Results

Nearby Solutions for Reduced Problem

v feas \implies dist(v, original fea. region) is small

Given any v s.t. $\overline{C} - \overline{A}^* v \succeq 0$, there exists y s.t. $C - A^* y \succeq 0$ and

 $\|\boldsymbol{y} - \mathcal{P}\boldsymbol{v}\| \leq \xi \cdot \|\boldsymbol{C} - \mathcal{A}^*\boldsymbol{y}\|,$

where $\xi > 0$ depends on \mathcal{A} , C, and D^* :

$$\begin{split} \xi &:= \frac{3\sqrt{2}}{\sigma_{\min}(\mathcal{A})} \left[\alpha(\mathcal{A}, \mathcal{C}) \frac{\|D^*\|}{\lambda_{\min}(D_+)} \right]^{1/2},\\ \text{and} \quad \alpha(\mathcal{A}, \mathcal{C}) &:= \begin{cases} \frac{\delta^*}{\sigma_{\min}(\mathcal{A})} & \text{if } \mathcal{C} \in \mathcal{R}(\mathcal{A}^*),\\ \frac{\delta^*}{\sigma_{\min}(\mathcal{A}_{\mathcal{C}})} & \text{if } \mathcal{C} \notin \mathcal{R}(\mathcal{A}^*). \end{cases} \end{split}$$

Theory Backward Stability Algorithm Numerical Results

Nearby Solutions for Original Problem

facial reduction steps = 1: y feas \implies dist(y, reduced fea. region) is small

Suppose $\exists \hat{v} \text{ s.t. } \delta_2^* := \lambda_{\min}(\bar{C} - \bar{\mathcal{A}}^* \hat{v}) > 0$. Given any *y* such that

 $Z = C - \mathcal{A}^*(y + y_Q) \succeq 0,$ there exists *v* such that $\overline{C} - \overline{\mathcal{A}}^* v \succeq 0$ and

$$\|\boldsymbol{y} - \mathcal{P}\boldsymbol{v}\| \leq \zeta \cdot \|\boldsymbol{Z}\| \cdot \frac{\delta_2^* + \sigma_{\max}(\mathcal{A}^*) \|\boldsymbol{y} - \mathcal{P}\hat{\boldsymbol{v}}\|}{\delta_2^* + \zeta \sigma_{\max}(\mathcal{A}^*) \|\boldsymbol{Z}\|}$$

where $\zeta > 0$ depends on \mathcal{A} , C, and D^* :

and

$$\zeta := \frac{2\sqrt{2}}{\sigma_{\min}(\mathcal{A}_{PQ}^{*})} \left[\alpha(\mathcal{A}, \mathcal{C}) \frac{\|D^{*}\|}{\lambda_{\min}(D_{+})} \right]^{1/2},$$

$$\mathcal{A}_{PQ}^{*} u := \sum_{i=1}^{m-\bar{m}} u_{i} (PP^{\top} \widehat{A}_{\bar{m}+i} PP^{\top} + PP^{\top} \widehat{A}_{\bar{m}+i} QQ^{\top} + QQ^{\top} \widehat{A}_{\bar{m}+i} PP^{\top}),$$

$$\mathcal{R} (\mathcal{A}^{*}) = \operatorname{span}(\widehat{A}_{1}, \dots, \widehat{A}_{m}) \text{ with } \mathcal{R} (\mathcal{A}^{*}) \cap \mathcal{R} (Q \cdot Q^{\top}) = \operatorname{span}(\widehat{A}_{1}, \dots, \widehat{A}_{\bar{m}}).$$

TheoryBackward StabilityAlgorithmNumerical Results

Numerics With/Without Facial Reduction

Computational results

- obtained using SeDuMi on MATLAB 7.11,
- performed on a machine with Intel Duo Core and 4GB RAM.
- First set of results are from specially generated test problems.
- Second set of results are from randomly generated instances where
 - there is a positive duality gap

• $v_{\rm P} = 0$

Theory	Ba
Algorithm	Nu

Backward Stability Numerical Results

Numerics With/Without Facial Reduction

Name	n	m	Optval	Optval
			with facial reduction	without facial reduction
Example 1	3	2	0	-6.30238e-016
Example 2	3	2	0	+0.570395
Example 3 $(v_P = 0, v_D = 1)$	3	4	0	+6.91452e-005
Example 4 (infea. dual)	3	3	0	+Inf
Example 5 (Slater's CQ holds)	10	5	+5.02950e+02 +5.02950e+02	
Example 6	6	8	+1	+1
Example 7	5	3	0	-2.76307e-012
Example 9a	20	20	0	Inf
Example 9b	100	100	0	Inf

[Solved using SeDuMi on MATLAB]

Theory	
Algorithm	

Numerics With/Without Facial Reduction

Name	n	т	Optval	Optval
			with facial reduction	without facial reduction
RandGen1	10	5	+1.5914e-015	+1.16729e-012
RandGen2	100	67	+1.1056e-010	NaN
RandGen3	200	140	+5.0557e-010	NaN
RandGen4	200	140	+1.02803e-009	NaN
RandGen5	120	45	-5.47393e-015	-1.63758e-015
RandGen6	320	140	+5.9077e-025	NaN
RandGen7	40	27	-5.2203e-029	+5.64118e-011
RandGen8	60	40	-2.03227e-029	NaN
RandGen9	60	40	+5.61602e-015	-3.52291e-012
RandGen10	180	100	+2.47204e-010	NaN
RandGen11	255	150	+7.71685e-010	NaN

[Solved using SeDuMi on MATLAB]

Backward Stability Numerical Results

- Minimal representations of the data regularize (P); use min. face f_P (and/or implicit rank reduction)
- goal: a backwards stable preprocessing algorithm to handle (feasible) conic problems for which Slater's CQ (almost) fails

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Backward Stability Numerical Results

Thanks for your attention!

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Theory

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