

Is asymptotic extremal graph theory of dense graphs trivial?

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joint work with Sergey Norin

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Introduction

- Asymptotic extremal graph theory has been studied for more than a century ([Mantel 1908](#), [Turán](#), [Erdős](#), [Lovász](#), [Szemerédi](#), [Bollobás](#), [Simonovits](#), ...).

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- Discovery of rich algebraic structure underlying many of these techniques.
- Neater proofs with no low-order terms.
- Methods for applying these techniques in semi-automatic ways.

algebraic inequalities

Many fundamental theorems in extremal graph theory can be expressed as **algebraic inequalities** between **subgraph densities**.

Theorem (Freedman, Lovász, Schrijver 2007)

*Every such inequality follows from the positive semi-definiteness of a certain **infinite** matrix.*

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A formal calculus capturing many standard arguments (induction, Cauchy-Schwarz,...) in the area.

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Observation (Lovász-Szegedy and Razborov)

*Every algebraic inequality between subgraph densities can be **approximated** by a **finite** number of applications of the Cauchy-Schwarz inequality.*

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- **Razborov**: Significant improvement over the known bounds on Turán's Hypergraph Problem.
- **Hladky, Kral, Norin**: Improved bounds for Caccetta-Haggkvist conjecture.
- **HH, Hladky, Kral, Norin, Razborov**: A question of Sidorenko and Jagger, Šťovíček and Thomason.
- **HH, Hladky, Kral, Norin, Razborov**: A conjecture of Erdős.

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Question (Razborov)

Can every true algebraic inequality between subgraph densities be proved using a finite amount of manipulation with subgraph densities of finitely many graphs?

Homomorphism Densities

Extremal graph theory

Studies the relations between the number of occurrences of different subgraphs in a graph G .

Extremal graph theory

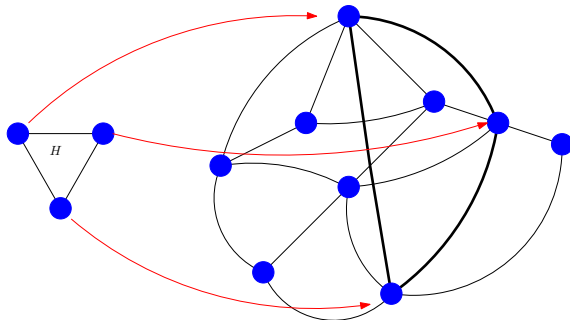
Studies the relations between the number of occurrences of different subgraphs in a graph G .

Equivalently one can study the relations between the “homomorphism densities”.

Homomorphism Density

Definition

- Throw the vertices of H on the vertices of G at random.

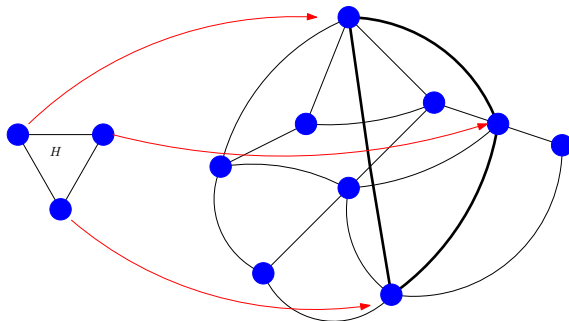


Homomorphism Density

Definition

- Throw the vertices of H on the vertices of G at random.

$$t_H(G) := \Pr[\text{edges go to edges}].$$



Definition

A map $f : H \rightarrow G$ is called a homomorphism if it maps edges to edges.

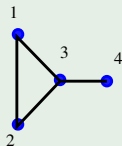
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A map $f : H \rightarrow G$ is called a homomorphism if it maps edges to edges.

$$t_H(G) = \Pr[f : H \rightarrow G \text{ is a homomorphism}].$$

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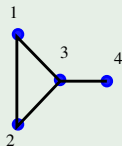
Example



$$t_{K_2}(G) = \frac{2 \times 4}{4^2} = \frac{1}{2}.$$

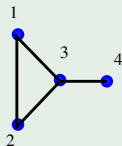
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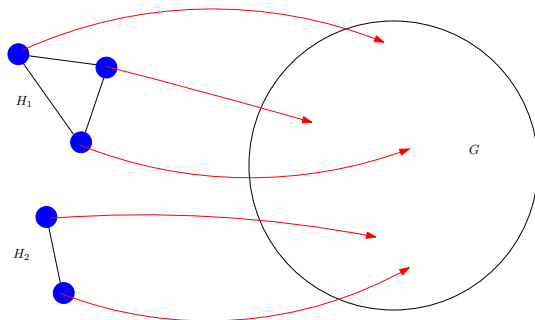


$$t_{K_3}(G) = \frac{3!}{4^3}.$$

- Asymptotically $t_H(\cdot)$ and subgraph densities are equivalent.

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- The functions t_H have nice algebraic structures:

$$t_{H_1 \sqcup H_2}(G) = t_{H_1}(G)t_{H_2}(G).$$



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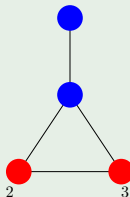
$$t_{K_3}(G) - 2t_{K_2 \sqcup K_2}(G) + t_{K_2}(G) \geq 0.$$

Algebra of Partially labeled graphs

Definition

A **partially labeled graph** is a graph in which some vertices are labeled by *distinct* natural numbers.

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Let H be partially labeled with labels L . For $\phi : L \rightarrow G$, define

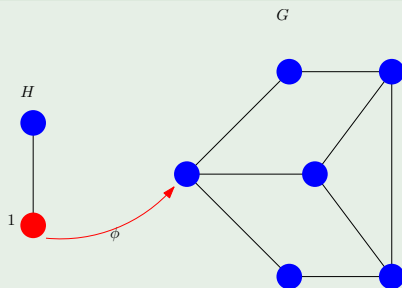
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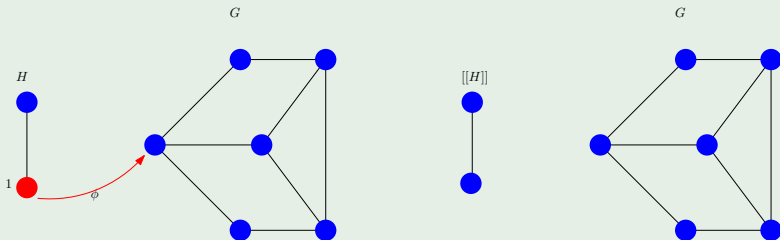
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$$t_{H,\phi}(G) = \frac{3}{6} = \frac{1}{2}.$$

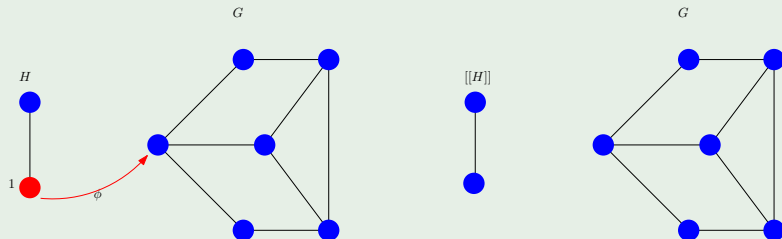
Example



Definition

Let $\llbracket H \rrbracket$ be H with no labels.

Example



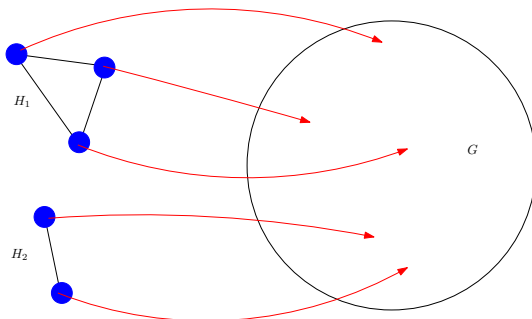
Definition

Let $[[H]]$ be H with no labels.

$$\mathbb{E}_{\phi} [t_{H,\phi}(G)] = t_{[[H]]}(G)$$

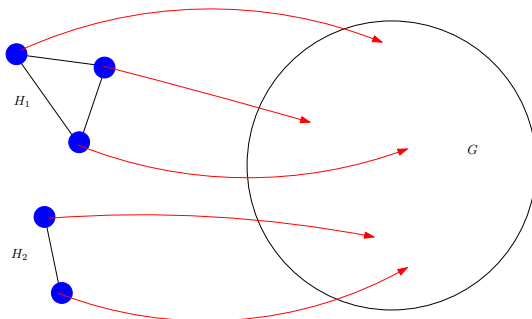
- Recall that:

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- This motivates us to define $H_1 \times H_2 := H_1 \sqcup H_2$.

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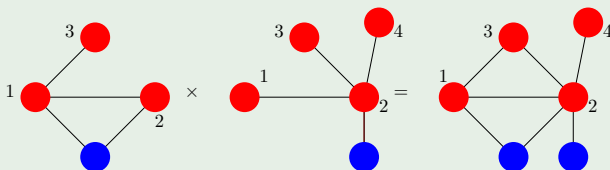
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The product $H_1 \cdot H_2$ of partially labeled graphs H_1 and H_2 :

- Take their disjoint union, and then identify vertices with the same label.
- If multiple edges arise, only one copy is kept.

Example

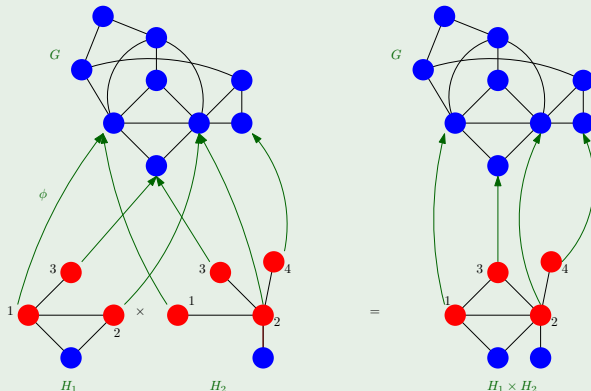


- Let H_1 and H_2 be partially labeled with labels L_1 and L_2 .

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- We have $t_{H_1, \phi}(G) t_{H_2, \phi}(G) = t_{H_1 \times H_2, \phi}(G)$.

Example



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$$\begin{aligned}
 0 &\leq \left(\sum b_i t_{H_i, \phi}(G) \right)^2 = \sum b_i b_j t_{H_i, \phi}(G) t_{H_j, \phi}(G) \\
 &= \sum b_i b_j t_{H_i \times H_j, \phi}(G)
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Abstraction

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- We want to understand the set of all positive quantum graphs.

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- Partially labeled quantum graphs form an algebra:

$$(a_1 H_1 + \dots + a_k H_k) \cdot (b_1 L_1 + \dots + b_\ell L_\ell) = \sum a_i b_j H_i \cdot L_j.$$

Unlabeling operator

$\llbracket \cdot \rrbracket : \text{partially labeled quantum graph} \mapsto \text{quantum graph}$

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Recall

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Corollary

Always

$$\llbracket g_1^2 + \dots + g_k^2 \rrbracket \geq 0.$$

Question (Lovász's 17th Problem, Lovász-Szegedy)

Is it true that every $f \geq 0$ is of the form

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Theorem (HH and Norin)

*The answer to the above questions is **negative**.*

Reflection positivity

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Observation (Lovasz-Szegedy and Razborov)

If $f \geq 0$ and $\epsilon > 0$, there exists a positive integer k and quantum labeled graphs g_1, g_2, \dots, g_k such that

$$-\epsilon \leq f - \llbracket g_1^2 + g_2^2 + \dots + g_k^2 \rrbracket \leq \epsilon.$$

positive polynomials

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Example (Motzkin's polynomial)

$$x^4y^2 + y^4z^2 + z^4x^2 - 6x^2y^2z^2 \geq 0.$$

Extending to quantum graphs

Theorem (HH and Norin)

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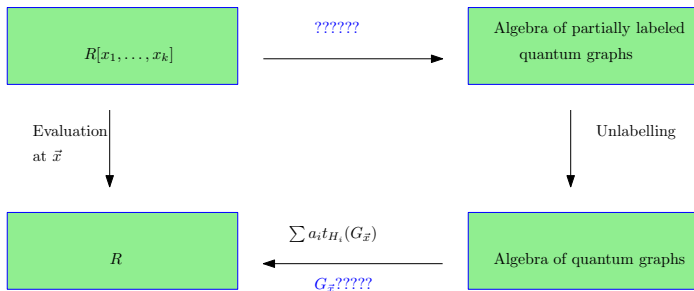
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- **Co-recursively enumerable** : Try to find a point that makes p negative.

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- **Co-recursively enumerable** : Try to find a point that makes p negative.
- **recursively enumerable** : Try to write $p = \sum (p_i/q_i)^2$.

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- Our solution to Lovász's 17th problem was based on an analogy to polynomials.
- Since there are polynomials which are positive but not sums of squares, our theorem was expected.
- **Lovász**: Does Artin's theorem (sums of rational functions) hold for graph homomorphisms?
- Maybe at least the decidability? (**A 10th problem**)

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Checking the positivity of $p \in \mathbb{R}[x_1, \dots, x_k]$ on \mathbb{Z}^k is undecidable.
- $\{t_H(G) : G\}$ is dense in $[0, 1]$, so maybe we should expect decidability for inequalities on homomorphism densities.

Theorem (HH and Norin)

The following problem is undecidable.

- **QUESTION:** *Does the inequality*
$$a_1 t_{H_1}(G) + \dots + a_k t_{H_k}(G) \geq 0 \text{ hold for every graph } G?$$

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Corollary

Analogue of Artin's theorem does not hold for homomorphism densities.

Proof

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Equivalently

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- **INSTANCE:** A polynomial $p(x_1, \dots, x_k)$ and graphs H_1, \dots, H_k .
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Instead I will prove the following theorem:

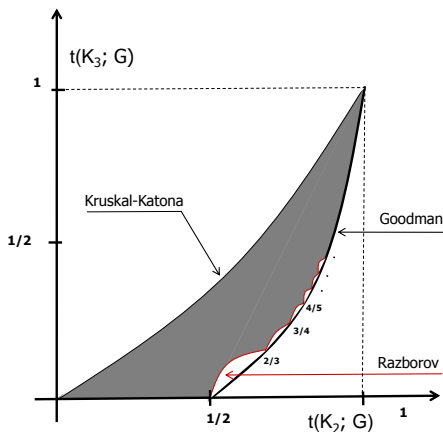
Theorem

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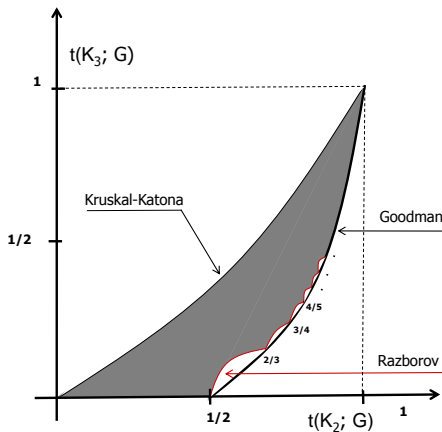
- **INSTANCE:** A polynomial $p(x_1, \dots, x_k, y_1, \dots, y_k)$.
- **QUESTION:** Does the inequality $p(t_{K_2}(G_1), \dots, t_{K_2}(G_k), t_{K_3}(G_1), \dots, t_{K_3}(G_k)) \geq 0$ hold for every G_1, \dots, G_k ?

- **Hilbert 10th:** Checking the positivity of $p \in \mathbb{R}[x_1, \dots, x_k]$ on $\{1 - \frac{1}{n} : n \in \mathbb{Z}\}^k$ is undecidable.

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- **Bollobás, Razborov:** Goodman's bound is achieved only when $t_{K_2}(G) \in \{1 - \frac{1}{n} : n \in \mathbb{Z}\}$.



Let S be the grey area and $g(x) = 2x^2 - x$. (Goodman:
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$$q := p \prod_{i=1}^k (1 - x_i)^6 + C_p \times \left(\sum_{i=1}^k y_i - g(x_i) \right).$$

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- $q < 0$ for some $x_i = t_{K_2}(G_i)$ and $y_i = t_{K_3}(G_i)$. (reduction)

Proof by computers!

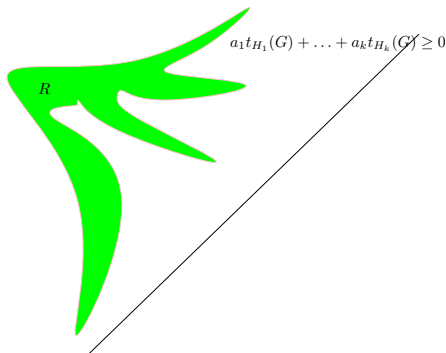


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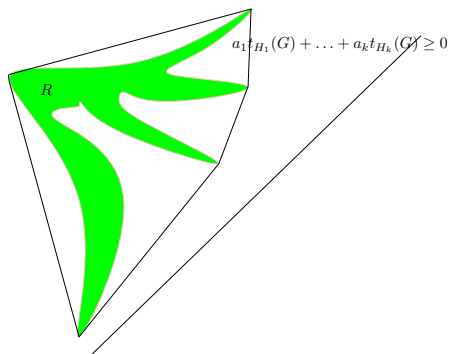
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Theorem (Freedman, Lovász, Schriver, Szegedy)

$(x_{H_1}, x_{H_2}, \dots) \in [0, 1]^{\mathbb{N}}$ belongs to the convex hull of $R \iff$

- **Condition I:** $x_{K_1} = 1$.
- **Condition II:** $x_{H \sqcup K_1} = x_H$ for all graph H .
- **Condition III:** The infinite matrix whose ij -th entry is $x_{\llbracket F_i \times F_j \rrbracket}$ is positive semi-definite.

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PSD method

Minimize

$$\sum a_i x_{H_i}$$

Subject to

(i) $x_{K_1} = 1$

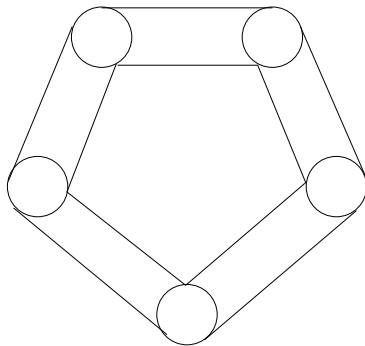
(ii) $x_{H \sqcup K_1} = x_H$ for all graphs H

(iii) The matrix $[x_{\llbracket F_i \times F_j \rrbracket}]_{ij}$ is p.s.d.

Application I: 5 cycles in triangle-free graphs

Theorem (HH, Hladky, Kral, Norin, Razborov)

Erdős's Conjecture 1982: The maximum number of cycles of length five in a triangle-free graph is $(n/5)^5$.



Application II: common graphs

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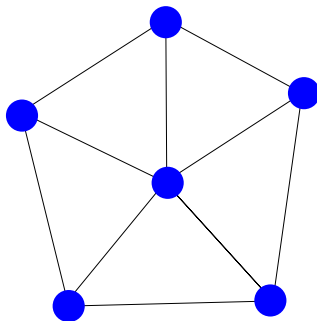
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- **Thomason 1989:** Any graph containing K_4 is not common.
- **Jagger, Stovicek, Thomason 1994 and Sidorenko 1994:**
Are there common non-3-colorable graph? Is W_5 common?

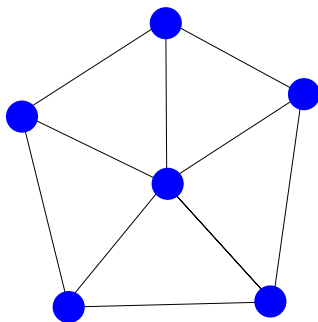
Theorem (HH, Hladky, Kral, Norin, Razborov)

The wheel W_5 is common.



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The wheel W_5 is common.



- The computer generated proof uses 150 Cauchy-Schwarz's.

Open problems

Recall

R is the closure of $\{(t_{H_1}(G), t_{H_2}(G), \dots) : G\} \subset [0, 1]^{\mathbb{N}}$.



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- Lovász's collection of open problems.

Thank you!