

# Conservative Dichotomy Revisited

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# Conservative CSP

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## Definition 1:

*Instance:* A triple  $(V, \mathcal{L}, \mathcal{C})$ , where  $V$  is a set of variables,  $\mathcal{C}$  is a set of constraints, and  $\mathcal{L}$  is a set of lists  $L_v$  for each  $v \in V$

*Question:* Is there a solution  $\varphi$  such that  $\varphi(v) \in L_v$  for every  $v \in V$

## Definition 2:

$\text{CSP}(\Gamma)$  where  $\Gamma$  contains all the unary relations

## Definition 3:

$\text{CSP}(\mathbb{A})$  where every operation of  $\mathbb{A}$  is conservative

$(f(x_1, \dots, x_n) \in \{x_1, \dots, x_n\})$

# Old Dichotomy

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2003/2011 paper, about 80 pages

What has happened over the last 8 years:

- Generalized Majority-Minority algorithm (Dalmau, 2006)
- Maroti's retraction (Maroti, this morning)
- Absorbing sets (Barto/Kozik, 2010)

# Conservative Algebras: Edges

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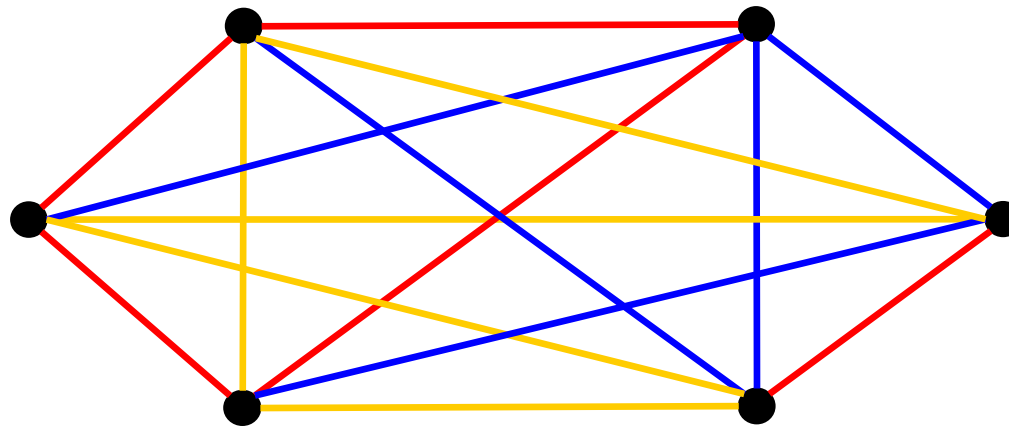
If  $\text{CSP}(\mathbb{A})$  is poly time, every 2-element subalgebra of  $\mathbb{A}$  must have one of the Schaefer's operations: semilattice, majority, or affine

- semilattice operation
- majority operation, no semilattice
- affine operation, no semilattice or majority

# Edge Coloured Graphs

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$G(\mathbb{A})$ :



## Theorem (B. 2003)

$\text{CSP}(\mathbb{A})$  for a conservative  $\mathbb{A}$  is poly time iff for any 2-element  $B \subseteq A$  there is  $f \in \text{Term}(\mathbb{A})$ , which is affine, majority, or semilattice; otherwise  $\text{CSP}(\mathbb{A})$  is NP-complete.

# Edge Coloured Graphs II

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## Lemma

There are  $f, g, h \in \text{Term}(\mathbb{A})$  such that  $f$  is semilattice on every semilattice edge,  $g$  is majority on every majority edge, and  $h$  is affine on every affine edge.

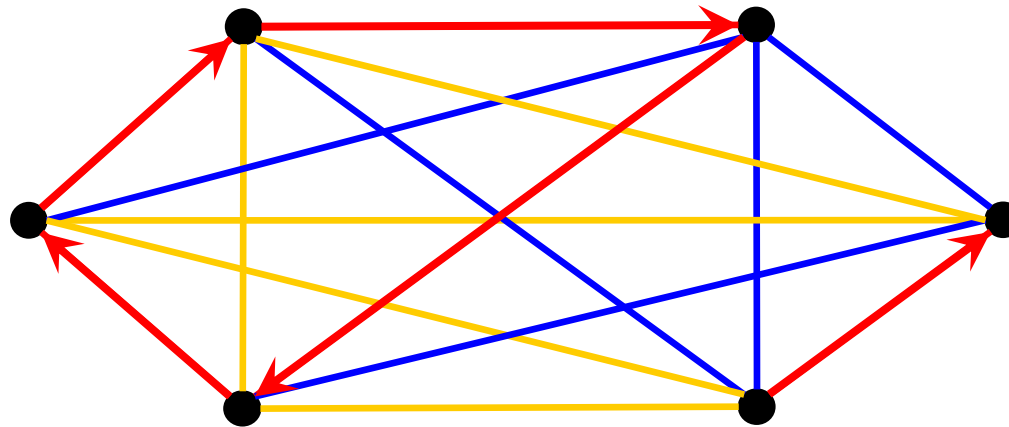
Extra conditions:

- $f(x, y) = x$  on every majority edge
- $g(x, y, z) = x$  on every affine edge
- $h(x, y, z) = x$  on every majority edge

# Edge Coloured Graphs III

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$G(A)$ :



As semilattice operation induces an order, red edges are directed  
 Will use  $\cdot$  instead of  $f$

# Graphs of Relations

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Sometimes we use more general graphs:  $G(\mathbb{A})$  where  $\mathbb{A}$  is not conservative, but a subdirect product of conservative algebras.

- $G(\mathbb{A})$  is constructed in the same way, edges are 2-element subalgebras.
- It is not complete, but any subalgebra induces a connected subgraph.
- All results and proofs remain the same with only minor tweaks



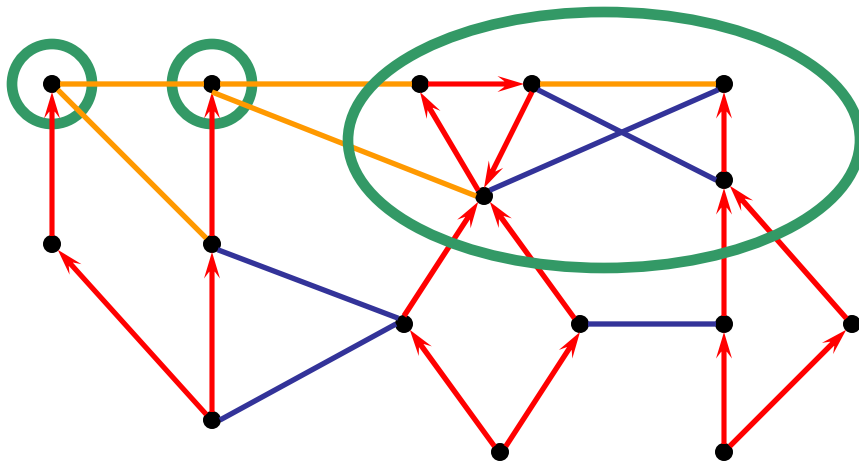
# AS-Components

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Let  $\mathbb{A}$  be a conservative algebra

$B \subseteq \mathbb{A}$  is called an **as-component** (affine-semilattice) if it is minimal with respect to the property:

there is no affine or semilattice (directed) edge in  $G(\mathbb{A})$  sticking out of  $B$



The remaining edges  
are majority

# Reductions

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We use two types of reductions:

- As-components exclusion
- Maroti's retractions

# Linked Relations

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Let  $R \subseteq A \times B$ .

$\lambda_A$  is a congruence on  $A$  defined as the transitive closure of the relation  $\{(a,b) \mid (a,c),(b,c) \in R \text{ for some } c\}$ ;

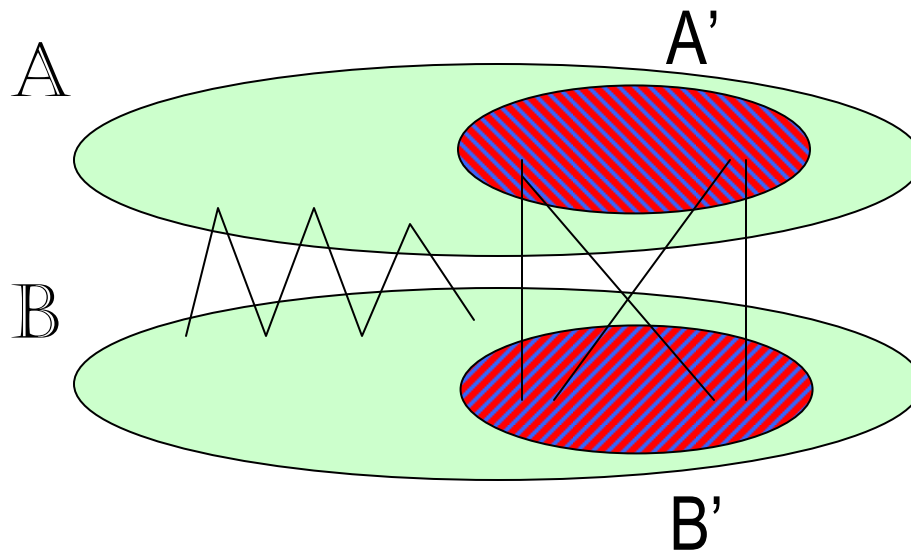
$\lambda_B$  is defined on  $B$  in the same way

$R$  is linked if  $\lambda_A, \lambda_B$  are total relations

# Binary Rectangularity

## Binary Rectangularity Lemma

Let  $R \leq A \times B$  be linked, and let  $A', B'$  be as-components of  $A$  and  $B$ , respectively, such that  $(a,b) \in R$  for some  $a \in A'$  and  $b \in B'$ . Then  $A' \times B' \subseteq R$ .



# The Connectivity Lemma

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## Connectivity Lemma

Let  $R \leq A_1 \times \dots \times A_n$ , and let  $A'_i$  be an as-component of  $A_i$ ,  $i \in [n]$ , such that  $(a_1, \dots, a_n) \in R$  for some  $a_i \in A'_i$ ,  $i \in [n]$ . Then  $R' = R \cap (A'_1 \times \dots \times A'_n)$  is a subdirect product of the  $A'_i$ , and  $R'$  is an as-component of  $R$ .

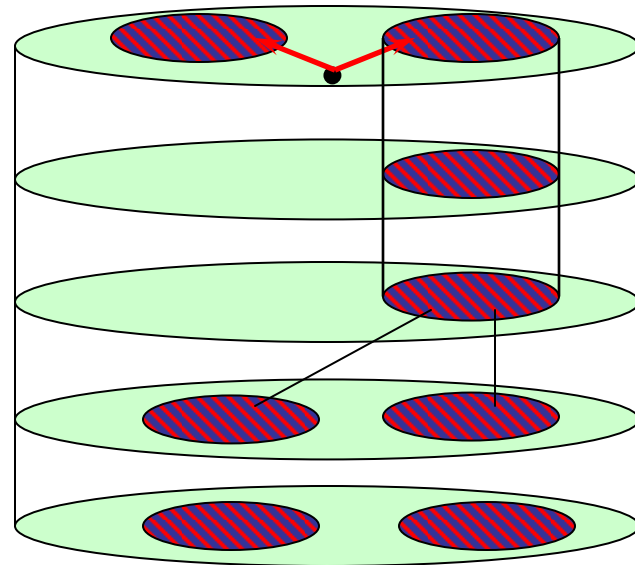
# Strands

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Let  $R \leq A_1 \times \dots \times A_n$  let  $A_i \subseteq A_i, A_j \subseteq A_j$  be as-components

Positions  $i$  and  $j$  are  $A_i, A_j$ -**related** if for any  $(a_1, \dots, a_k) \in R$   
 $a_i \in A_i$  iff  $a_j \in A_j$

$I \in \{1, \dots, k\}$  is a **strand**  
 w.r.t. as-components  
 $A_1, \dots, A_k$  if any  $i, j \in I$  are  
 $A_i, A_j$ -related



# Rectangularity

## Rectangularity Lemma

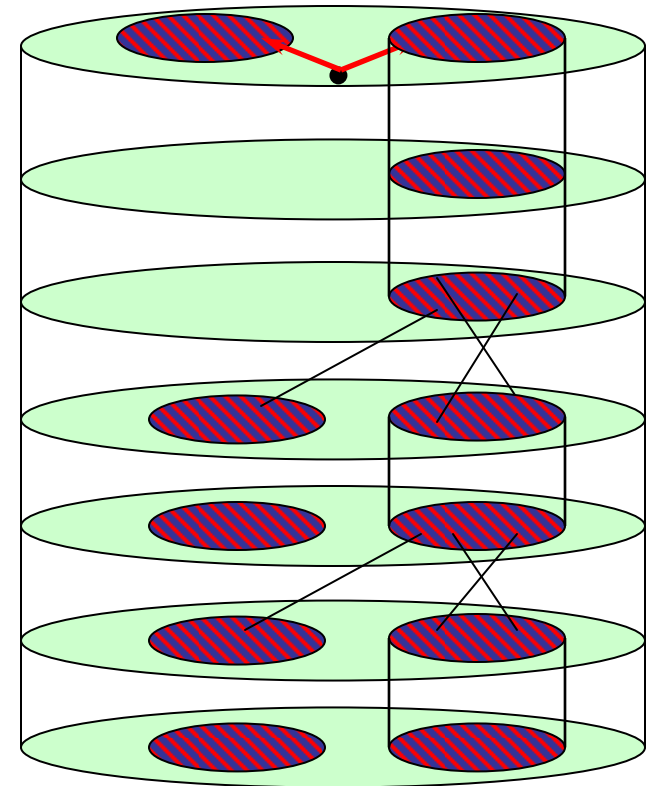
Let  $R \leq \mathbb{A}_1 \times \dots \times \mathbb{A}_n$  and  $A_1, \dots, A_k$  as-components such that

$R \cap (A_1 \times \dots \times A_k) \neq \emptyset$ . Let also

$I_1, \dots, I_k$  be the partition of  $\{1, \dots, n\}$  into strands w.r.t.

$A_1, \dots, A_s$  and  $R_i = \text{pr}_{I_i} R \cap \prod_{j \in I_i} A_j$ .

Then  $R_1 \times \dots \times R_s \subseteq R$ .



# Rectangularity: to Binary Relations

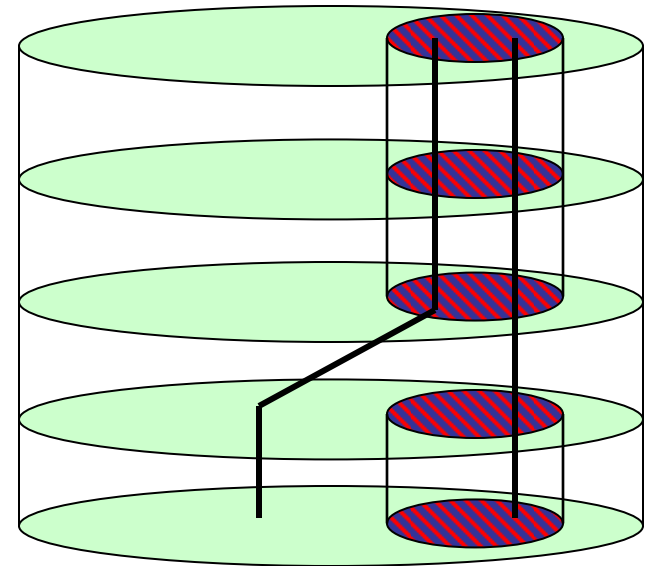
Let  $R \leq A_1 \times \dots \times A_n$  and  $A_1, \dots, A_n$  be as-components such that there is  $(a_1, \dots, a_n) \in R$  with  $a_i \in A'_i$

Suppose 1 and  $n$  belong to different strands. Then there are  $(a, b), (a', b') \in \text{pr}_{1,n} R$  with  $a, a' \in A'_1$ ,  $b \in A'_n$ ,  $b' \in A_n - A'_n$ .

Take  $(a_1, \dots, a_n), (b_1, \dots, b_n) \in R$   
with  $a_1 = a, b_1 = a', a_n = b, b_n = b'$

Let  $a_i \in A'_i$  for  $i \in [m]$  and  
 $a_i \in A_i - A'_i$  otherwise

Consider  $R$  as a binary relation on  
 $\text{pr}_{[m]} R$  and  $\text{pr}_{\{m+1, \dots, n\}} R$





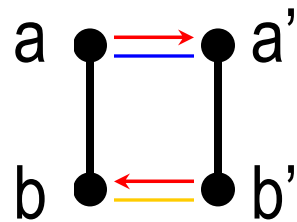
# Rectangularity: Binary Relations

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Let  $R \leq A \times B$ , and  $A', B'$  as components of  $A, B$  such that there are  $(a, b), (a', b') \in R$  with  $a, a' \in A'$ ,  $b \in B'$ ,  $b' \in B - B'$ .

We prove that the pairs can be chosen such that  $(a', b) \in R$ .

There is a sequence in  $R$



Since  $bb'$  is a majority or semilattice edge, either

$$\frac{a'}{b} = \frac{a}{b} \cdot \frac{a'}{b'} \in R \quad \text{or} \quad \frac{a'}{b} = h\left(\frac{a}{b}, \frac{a}{b}, \frac{a'}{b'}\right) \in R$$

# Rectangularity: Congruences of AS-Components

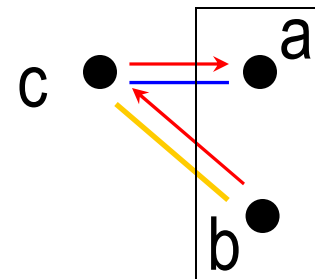
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Let  $A'$  be an as-component of  $\mathbb{A}$ , and  $\lambda$  a congruence such that  $(a,b) \in \lambda$  for some  $a \in A'$  and  $b \in A - A'$ . Then  $A'$  is in a  $\lambda$ -block.

If  $\lambda$  is nontrivial on  $A'$ , choose  $c \in A'$  with  $(a,c)$  not in  $\lambda$  and such that  $ca$  is either semilattice or affine. Since  $b \in A - A'$ ,  $bc$  is either semilattice or majority.

Then  $(b \cdot c, b) \in \lambda$ , so  $b \cdot c = b$ , and  $bc$  is majority.

$$\text{Then } \frac{a}{c} = g\left(\begin{matrix} a & c & c \\ b & c & c \end{matrix}\right) \in \lambda$$



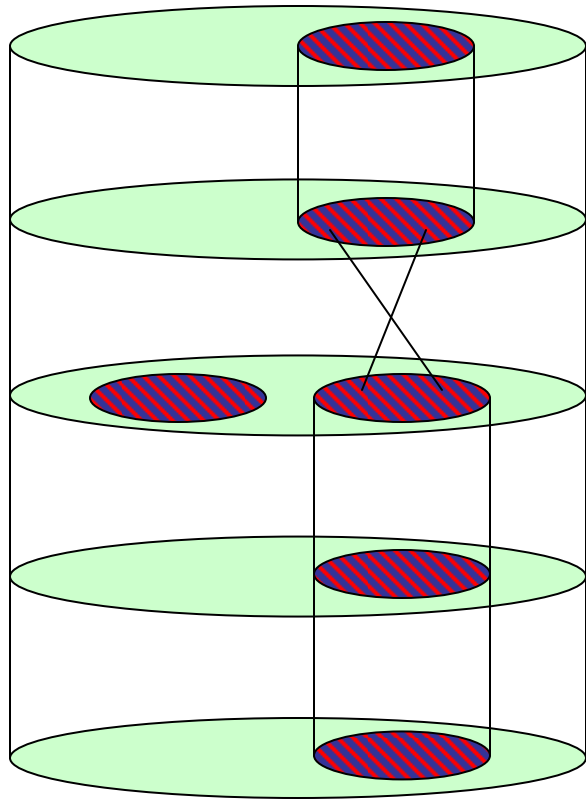
## Rectangularity: Binary Relations II

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- Let  $R \leq A \times B$ , and  $A', B'$  as-components of  $A, B$  such that the coordinate positions of  $R$  are not  $A', B'$ -related. W.l.o.g. there are  $(a, b), (a, b') \in R$  with  $a \in A'$ ,  $b \in B'$ ,  $b' \in B - B'$ .
- Let  $\lambda$  be the link congruence on  $B$ .  $B'$  is inside a  $\lambda$ -block.
- Therefore if  $R'$  denotes the restriction of  $R$  on the link congruences blocks containing  $A'$  and  $B'$ ,  $R'$  is linked.
- By the Binary Rectangularity Lemma  $A' \times B' \subseteq R$

# AS-Components Exclusion

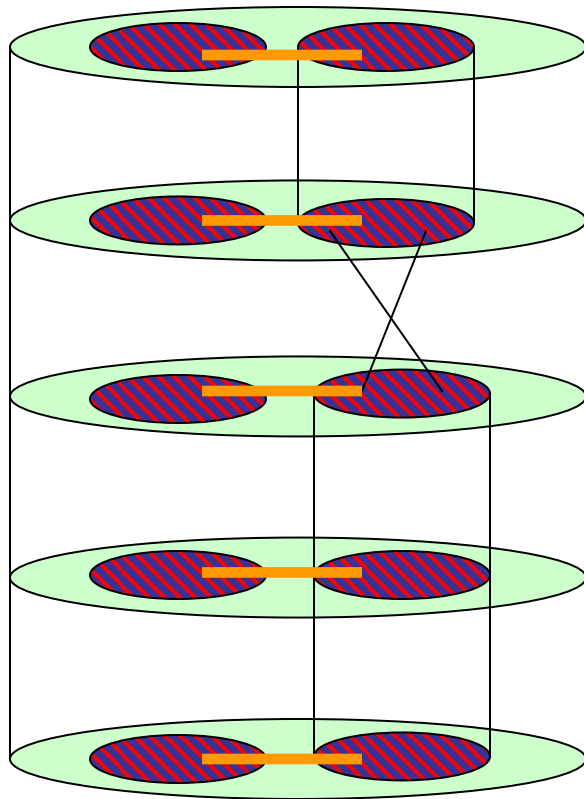
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1. find as-components  $A_v, v \in V$ , such that for any  $v, w \in V$  as-components  $A_v, A_w$  are consistent
2. find the strands
3. for each strand  $W$  solve the problem restricted to  $W$  and  $A_w, w \in W$
4. if every such problem has a solution, by the Rectangularity Lemma any combination of such solutions gives a solution to the problem
5. otherwise remove elements of the failed as-components

# Finding Strands and Components

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2. To find a strand take a variable  $v$  and an as-component  $A$  of  $A_v$  and find all variables  $w$  and as-components  $B$  of  $A_w$  such that  $v, w$  are  $A, B$ -related

1.

## Chinese Remainder Theorem

Let  $R \leq A_1 \times \dots \times A_n$  and let  $A'_1, \dots, A'_n$  be as-components of  $A_1, \dots, A_n$  such that for any  $i, j \in \{1, \dots, n\}$  there is a tuple  $(a_1, \dots, a_k) \in R$  such that  $a_i \in A'_i, a_j \in A'_j$ .  
 Then  $R \cap (A'_1 \times \dots \times A'_n)$  is a sudirect product of  $A'_1, \dots, A'_n$

# Maroti's Retractions

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Let  $\mathcal{A}$  be a class of algebras of similar type closed under taking subalgebras and retracts via idempotent unary polynomials.

Let  $A \in \mathcal{A}$  and  $f$  a binary term operation of  $A$  such that

- a.  $f(x, f(x, y)) = f(x, y)$ ;
- b. for each  $a \in A$  the map  $x \mapsto f(x, a)$  is not surjective
- c. the set  $C$  of  $a \in A$  such that  $x \mapsto f(x, a)$  is surjective generates a proper subalgebra of  $A$

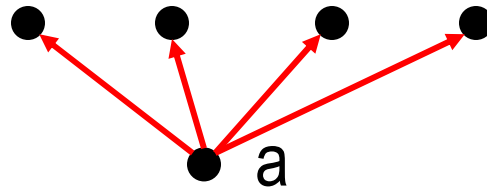
Then  $\text{CSP}(\mathcal{A})$  is poly time reducible to  $\text{CSP}(\mathcal{A} - \{A\})$ .

# Reductions II

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a.  $x \cdot (x \cdot y) = x \cdot y$

b. If  $a \in A$  is such that  $x \rightarrow a \cdot x$  is surjective, then  $a \cdot x = x$

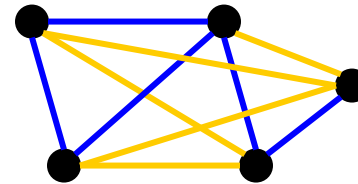


$a$  does not belong to any as-component, and we can use as-components exclusion

## Reductions III

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- c. If  $x \rightarrow x \cdot a$  is surjective for each  $a$ , then  $x \cdot a = x$ , and  $\mathbb{A}$  has no semilattice edges.



An operation  $m(x,y,z)$  on  $A$  is called Generalized Majority Minority if for any  $a,b \in A$ ,

either  $m(x,x,y) = m(x,y,x) = m(y,x,x) = x$  for  $x,y \in \{a,b\}$ ,  
or  $m(x,x,y) = m(y,x,x) = y$

Operation  $g(h(x,y,z),y,z)$  is majority on each majority edge,  
and is affine on each affine edge

Then there is a ternary GMM operation on  $\mathbb{A}$ .



# Question

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Are as-components and minimal absorbing subuniverses the same??