# Computational Speedup in Spatial Bayesian Image Modeling via GPU Computing

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#### **Outline**

- Introduction
- GPU Computing
- Results
- Concluding Remarks

# Problem 1: Group fMRI Analysis

- Hierarchical Bayesian model <sup>1</sup>
- Level 1 Likelihood
  - separate model for each subject (typically 10 20)
  - image intensities finite mixture with spatially correlated weights
  - several components required to model each activation region
  - component means finite mixture about activation centers
- Level 2
  - Activation centers finite mixture model about population centers
- Level 3
  - Population centers modeled as a homogeneous Poisson process

I will concentrate on Level 1 as it is the most computationally intense

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<sup>&</sup>lt;sup>1</sup>Xu, Johnson, Nichols, Nee. Biometrics 2009

# Problem 1: Group fMRI Analysis

- $y_v$  intensity at voxel v. Approximately 250K voxels
- $x_v (x, y, z)$  spatial location
- $\phi(y; m, v)$  Gaussian density w/mean m and variance v

#### Likelihood and weights

All voxels assumed to be conditionally independent

$$\pi(y_{\nu} \mid \cdot) = w_0 \phi(y_{\nu}; 0, 1) + \sum_{j=1}^{J} w_j \phi(y_{\nu}; \theta_j, \sigma_j^2)$$

$$w_0 \propto m$$

$$w_j = \frac{\phi_3(x_v; \eta_j, \Psi_j)}{m + \sum_{k=1}^J \phi_3(x_v; \eta_k, \Psi_k)}$$

where  $\phi_3(x; \eta, \Psi)$  is the Gaussian density in  $\mathbb{R}^3$  at x

#### Definition

Suppose that  $Y = \{Y(\xi) : \xi \in S\}$  is a non-negative random field such that with probability one,  $\xi \to Y(\xi)$  is a locally integrable function. If  $[X \mid Y] \sim PP(S, Y)$ , then X is said to be a *Cox process* driven by Y.

#### Density

The density of  $X_B = X \cap B$ ,  $B \subset S$ , w.r.t. the dist. of a unit-rate Poisson process, is given by

$$\pi(x) = \mathbb{E}\left[ \exp\left(|B| - \int_B Y(\xi) d\xi 
ight) \prod_{\xi \in x} Y(\xi) 
ight].$$

Furthermore, the intensity function is given by  $\rho(\xi) = \mathbb{E}[Y(\xi)]$ .

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#### Definition

If  $Y(\xi)$  is a Gaussian Process, then the Cox process driven by  $\exp[Y(\xi)]$  is said to be a Log Gaussian Cox Process.

The dist of (X, Y) completely determined by

$$m(\xi) = \mathbb{E}(Y(\xi))$$
  $c(\xi, \eta) = Cov(Y(\xi), Y(\eta)).$ 

• We will assume  $c(\xi, \eta)$  is translation invariant and isotropic

$$c(||\xi - \eta||) = \sigma^2 r(\alpha ||\xi - \eta||).$$

A useful family of corr functions is the power exponential family:

$$r(\alpha||\xi - \eta||) = \exp(-\alpha||\xi - \eta||^{\delta}).$$

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- obviously, this matrix is too large to handle computationally

 The structure of most stationary correlation matrices is Toeplitz in 1D, block Toeplitz in 2D and nested block Toeplitz in 3D.

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- *F* turns out to be the DFT matrix and  $\Lambda = \sqrt{n} \operatorname{diag}(Fc)$ .
  - c is of length  $n = 2.91 \times 2.109 \times 2.91 (= 7,221,032)$
  - for any vector v,  $Fv \leftrightarrow DFT(v)$  and  $F^Hv \leftrightarrow IDFT(v)$ .

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- Thus, we can use the FFT which is extremely fast:  $O(n \log_2 n)$ .

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#### For example

$$C^{1/2}v = F\Lambda^{1/2}F^{H}v$$

$$= F\sqrt{n}\operatorname{diag}(\sqrt{Fc})F^{H}v$$

$$= \sqrt{n}\operatorname{DFT}(\sqrt{\operatorname{DFT}(c)} \odot \operatorname{IDFT}(v))$$

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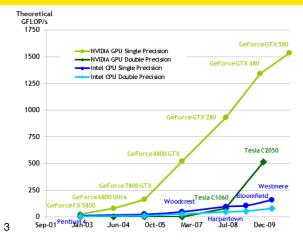
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$$= F\Lambda^{-1}F^{H}b$$

$$= \frac{1}{\sqrt{n}}DFT(IDFT(b) \oslash DFT(c))$$

Although the FFT algorithm is fast, the base c in this problem is very large. Hence, overall the MCMC algorithm is quite slow.

# General Purpose Graphical Processing Unit Computing



Tesla C2050 GPU Single Precision — 1 TFLOP/s

Tesla M2090 GPU Single Precision — 1.33 TFLOP/s

<sup>&</sup>lt;sup>3</sup>Source: NVIDIA CUDA C Programming Guide, Version 4.0

### General Purpose Graphical Processing Unit Computing

How do GPUs achieve such stunning performance gains over CPUs?

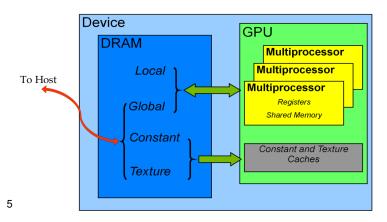


More transistors are dedicated to data processing rather than data caching and control flow. GPUs typically have hundreds of cores. (Tesla C2050 — 448 cores)

<sup>&</sup>lt;sup>4</sup>Source: NVIDIA CUDA C Programming Guide, Version 4.0

# **GPU Memory Bandwidth Limitations**

Key to performance with HDD: reduce memory transfers  $\leftarrow$  More memory Faster transfer  $\rightarrow$ 



<sup>&</sup>lt;sup>5</sup>Source: NVIDIA CUDA C Best Practices Guide, Version 4.0

#### How is the GPU exposed in C/C++?

#### C/C++ interface (CUDA and OpenCL)

- allocate memory on the GPU special malloc function
- copy data from CPU to GPU memory special function
- call kernel function—executes on GPU, N times in parallel
- copy GPU output to CPU special function
- free GPU memory

 kernels—C function, executed on GPU, N times in parallel by N threads

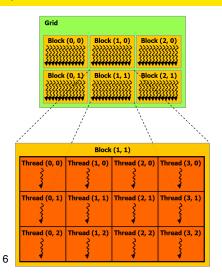
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- The host interface (the C program) calls the kernel
  - specifies the number of threads/block and the number of blocks/grid
  - passes memory pointers where data have been loaded onto the GPU.

#### How is the GPU exposed in C/C++?



<sup>&</sup>lt;sup>6</sup>Source: NVIDIA CUDA C Programming Guide, Version 4.0

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### Example

#### Typical C code

- loop over voxels, i
  - loop over mixture comp, j
    - calculate jth comp contrib. to likelihood at voxel i
  - end mixture comp loop
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# Example

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#### Parallel code

- allocate GPU memory
- data copy CPU → GPU
- call kernel()
- data copy GPU → CPU
- free GPU memory

kernel code (voxels processed in parallel)

- load shared memory
- loop over mixture comp, j
  - calculate jth comp contrib.
     to likelihood at voxel i
- end mixture comp loop

# Gain in Computational Efficiancy

Model	Processor	Iterations	Memory	Time	Factor
		(x 1000)	(Mb)	(hours)	
groupFMRI	CPU	5		229.52	
	GPU	5	175	1.20	191
LGCP	CPU	120		595.00	
	GPU	120	311	4.75	125

#### Caveats

- reduce data transfers
- beware of conditional branching (e.g. if-else)
- bank conflicts (memory)
  - for shared memory, a bank conflict occurs when two or more threads within a warp try to access different bytes of memory within the bank
    - a warp is a group of 32 threads—minimum size of data processed. A hardware constraint. Blocks of threads are further divided into warps.
  - if two or more threads try to access the same byte, no bank conflict occurs
  - may need to pad arrays to avoid bank conflicts
  - in my opinion, most difficult aspect of GPGPU computing

#### Conclusions

- GPGPU computing can substantially improve performance
- Likelihood based approaches can also gain from GPGPU computing
- In general, the larger the problem, the greater the practical gains.