Large Deviations

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- Mathematical Model with noise or randomness in it.
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- Instead we have probabilities for "yes" and "no".

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- Asymptotic methods are needed.

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- How small?

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- Let us look at some examples.

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- At what rate?

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- Is there a recipe for computing H(p)?

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- $P \simeq e^{-nH + o(n)}$

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- Change the model, but do it frugally!

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- $c(A) = \inf_{x \in A} I(x)$

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 $If \inf_{x \in A^o} I(x) = \inf_{x \in \bar{A}} I(x)$

$$\lim_{n \to \infty} \frac{1}{n} \log P_n(A) = -\inf_{x \in A} I(x)$$

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$$\exp[nF(x)]dP_n = \exp[n\sup_x [F(x) - I(x)] + o(n)]$$

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Exponential

$$\exp\left[-\frac{1}{\epsilon}I(f(\cdot))\right]$$

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$$\frac{dQ_{\epsilon}}{dP_{\epsilon}} = \exp\left[\frac{1}{\epsilon} \int_{0}^{T} (f'(t) - b(x(t))(dx(t) - b(x(t))dt)\right]$$
$$-\frac{1}{2\epsilon} \int_{0}^{T} ||f'(t) - b(x(t))||^{2} dt$$

Exit problem

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 $\overline{Y} = \mathcal{M}(X)$

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Now $\nu_n \to \overline{\mu}$

$$\frac{1}{n}\log\frac{dQ_n}{dP_n} = \frac{1}{n}\sum_{i}\log\frac{\pi'(X_i, X_{i+1})}{\pi(X_i, X_{i+1})}$$

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$$\lambda(V) \Leftrightarrow I(\mu)$$
$$I(\mu) = \sup_{u > 0} -\int \frac{\mathcal{A}u}{u} d\mu$$

 $\blacksquare \mathcal{A}$ is self adjoint with respect to α .

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$$I(\mu) = \mathcal{D}(\sqrt{\frac{d\mu}{d\alpha}})$$

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$$\lambda(V) = \sup_{g:||g||_{2,\alpha}=1} \left[\int Vg^2 d\alpha - \mathcal{D}(g) \right]$$

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- d=1, Brownian motion, L_1 or similar.

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- Random graphs, Random matrices

THE END