

Non-commutative polynomial optimization and the variational RDM method

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RDM method:

N-representability of the 2-RDM \rightarrow positivity constraints.

\Rightarrow Semidefinite programming can be used to compute electronic energy of atoms and molecules.

This talk: The SDP-based RDM method is a special case of a more general and abstract approach (not motivated by *N*-representability).

- We have introduced a method to solve *non-commutative polynomial optimization* problems.
- Computing the energy of a system of N electrons is an instance of these optimization problems.
- When applying our method to this particular instance, we recover the RDM method.

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- 1 The RDM method based on semidefinite programming
- 2 Non-commutative polynomial optimization (NCPO)
- 3 The RDM method as a special case of NCPO
- 4 Other applications of NCPO

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Variational approach

Introduce finite set of R basis functions (orbitals)

$$\phi_1(\mathbf{r}), \phi_2(\mathbf{r}), \dots, \phi_R(\mathbf{r})$$

and expand the N -electron wave-function $\Psi(\mathbf{r}_1, \dots, \mathbf{r}_N)$ as

$$\Psi(\mathbf{r}_1, \dots, \mathbf{r}_N) = \sum_{j_1 \dots j_N} c_{j_1 \dots j_N} \phi_{j_1}(\mathbf{r}_1) \dots \phi_{j_N}(\mathbf{r}_N).$$

The corresponding Hamiltonian is

$$H = \sum_{ij} H_{ij} a_i^\dagger a_j + \sum_{ijkl} H_{ijkl} a_i^\dagger a_j^\dagger a_k a_l$$

where a_i, a_i^\dagger are the annihilation and creation operators associated to orbital ϕ_i .

Energy is determined by the 2-RDM

$$\begin{aligned} E = \langle H \rangle &= \sum_{ij} H_{ij} \langle a_i^\dagger a_j \rangle + \sum_{ijkl} H_{ijkl} \langle a_i^\dagger a_j^\dagger a_k a_l \rangle \\ &= \sum_{ij} H_{ij} {}^1D_{ij} + \sum_{ijkl} H_{ijkl} {}^2D_{ijkl} \end{aligned}$$

where 1D and 2D are the one-particle and two-particle reduced density matrices defined by

$${}^1D_{ij} = \langle a_i^\dagger a_j \rangle$$

$${}^2D_{ijkl} = \langle a_i^\dagger a_j^\dagger a_k a_l \rangle$$

Note that ${}^1D_{ij} = (N - 1) \sum_k {}^2D_{ikkj}$.

N -representability problem

Ground state energy:

$$E_g = \min_{^2D} \left\{ \sum_{ij} H_{ij} {}^1D_{ij} + \sum_{ijkl} H_{ijkl} {}^2D_{ijkl} \right\}$$

N -representability problem

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Problem: this yields an energy far lower than the exact ones because not every 2-particle density matrix 2D originates from a N -particle wavefunction $|\psi\rangle$.

We must impose N -representability conditions on 2D

Necessary N-representability conditions

① Positivity conditions [Coleman 63, Garrod and Percus 64]:

- ${}^1D \succeq 0$, ${}^1Q \succeq 0$, where

$${}^1D_{ij} = \langle a_i^\dagger a_j \rangle$$

$${}^1Q_{ij} = \langle a_i a_j^\dagger \rangle$$

- D, Q, G conditions: ${}^2D \succeq 0$, ${}^2Q \succeq 0$, ${}^2G \succeq 0$, where

$${}^2D_{ijkl} = \langle a_i^\dagger a_j^\dagger a_k a_l \rangle$$

$${}^2Q_{ijkl} = \langle a_i a_j a_k^\dagger a_l^\dagger \rangle$$

$${}^2G_{ijkl} = \langle a_i^\dagger a_j a_k^\dagger a_l \rangle$$

② Linear conditions that relate all these matrices to 2D

For instance:

$${}^1Q_{ij} = \delta_{ij} - {}^1D_{ij} = (N-1) \sum_k {}^2D_{ikkj}$$

$${}^2Q_{ijkl} = \delta_{jk} {}^1D_{il} - {}^2D_{ikjl}$$

Ground-state energy from semidefinite programming

Minimization of

$$E = \sum_{ij} H_{ij} {}^1D_{ij} + \sum_{ijkl} H_{ijkl} {}^2D_{ijkl}$$

subject to the previous positivity and linear constraints is a typical instance of **semidefinite programming**.

This minimization problem can be solved **exactly**.

It provides a **lower-bound** on the ground-state energy.

Higher-order constraints

A whole hierarchy of additional SDP constraints can be added to increase accuracy [Erdahl, Jin 00], [Mazziotti, Erdahl 01].

E.g.: Positivity conditions on the 3-RDMs:

$${}^3D_{ijklmn} = \langle a_i^\dagger a_j^\dagger a_k^\dagger a_l a_m a_n \rangle$$

$${}^3E_{ijklmn} = \langle a_i^\dagger a_j^\dagger a_k a_l^\dagger a_m a_n \rangle$$

$${}^3F_{ijklmn} = \langle a_i^\dagger a_j a_k a_l^\dagger a_m^\dagger a_n \rangle$$

$${}^3Q_{ijklmn} = \langle a_i a_j a_k a_l^\dagger a_m^\dagger a_n^\dagger \rangle$$

Why are all these matrices positive?

- All previous matrices are of the form

$$\begin{aligned} M_{ij} &= \langle \Psi | C_i^\dagger C_j | \Psi \rangle \\ &= \langle v_i | v_j \rangle \end{aligned}$$

where $|v_i\rangle = C_i|\Psi\rangle$.

For instance: ${}^2G_{ijkl} = \langle a_i^\dagger a_j a_k^\dagger a_l \rangle = \langle \Psi | C_{ji}^\dagger C_{kl} | \Psi \rangle$
with $C_{kl} = a_k^\dagger a_l$.

- M is positive semidefinite if and only if $M_{ij} = \langle v_i | v_j \rangle$.

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Non-commutative polynomial optimization

Let $X = \{X_1, \dots, X_n\}$ be n variables that we view as operators acting on some unspecified Hilbert space H .

Let $p(X), q(X), r(X)$ be polynomials in X .

E.g.: $p(X) = X_1 + 3X_1X_2 - 4X_4X_2$.

We want to solve

$$\begin{array}{ll} \min_{H, X, |\Psi\rangle} & \langle \Psi | p(X) | \Psi \rangle \\ \text{s.t.} & q(X) = 0 \\ & r(X) | \Psi \rangle = 0 \end{array}$$

Note: H is not fixed in advance, $\dim(H)$ is not bounded.

Why *non-commutative* optimization?

If we add the commutativity constraints $X_i X_j + X_j X_i = 0$, the scalar representation $X_i = x_i \in \mathbb{R}$ is always a solution.

The optimization problem is then equivalent to a standard polynomial optimization over \mathbb{R}^n

$$\begin{array}{ll} \min_{x \in \mathbb{R}^n} & p(x) \\ \text{s.t.} & q(x) = 0 \\ & r(x) = 0 \end{array}$$

Solving NCPO through semidefinite programming

In [arXiv:0903.4368](#), we introduce a [sequence of relaxations](#) R_i that provide lower-bounds on the global solution of the original NCPO problem.

Each of these relaxations R_i is a [semidefinite program](#).

In the limit $R_i \rightarrow R_\infty$, the lower-bounds [converge](#) to the global solution of NCPO.

Notation

Monomials

- Given the n operators X_1, \dots, X_n ,
a monomial X_α of degree k is a product of k operators X_i :

$$X_\alpha = X_{\alpha_1} X_{\alpha_2} \dots X_{\alpha_k}$$

- We denote the identity operator I as the monomial $X_0 = I$.
- The product of two monomials $X_\alpha X_\beta$ is itself a monomial that we denote $X_{\alpha\beta} = X_\alpha X_\beta$.

Polynomials

A polynomial $p(X)$ is a linear combination of monomials

$$p(X) = \sum_{\alpha} p_{\alpha} X_{\alpha}$$

Construction of the relaxation of degree k

- Consider the set of all vectors of the form

$$\{X_\alpha|\Psi\rangle\} = \{|\Psi\rangle, X_i|\Psi\rangle, X_iX_j|\Psi\rangle, \dots, (X_{i_1} \dots X_{i_k})|\Psi\rangle\}$$

where X_α is at most of degree k .

Then the matrix kM with entries

$${}^kM_{\alpha\beta} = \langle\Psi|X_\alpha^\dagger X_\beta|\Psi\rangle$$

is positive definite: ${}^kM \succeq 0$.

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- The objective function

$$\begin{aligned}\langle\Psi|p(X)|\Psi\rangle &= \sum_{\alpha} p_{\alpha} \langle\Psi|X_{\alpha}|\Psi\rangle \\ &= \sum_{\alpha} p_{\alpha} \langle\Psi|X_0^\dagger X_{\alpha}|\Psi\rangle = \sum_{\alpha} p_{\alpha} {}^kM_{0\alpha}\end{aligned}$$

is a linear function of the entries of kM

Construction of the relaxation of degree k

- The conditions $q(X) = 0$ and $r(X)|\Psi\rangle = 0$ induce **linear constraints** on kM :

$$\begin{aligned} q(X) = \sum_{\nu} q_{\nu} X_{\nu} = 0 &\Rightarrow \sum_{\nu} q_{\nu} \langle \Psi | X_{\alpha}^{\dagger} X_{\nu} X_{\beta} | \Psi \rangle = 0 \quad \forall \alpha, \beta \\ &\sum_{\nu} q_{\nu} \langle \Psi | X_{\alpha}^{\dagger} X_{\nu} | \Psi \rangle = 0 \quad \forall \alpha, \beta \\ &\sum_{\nu} q_{\nu} {}^kM_{\alpha, \nu \beta} = 0 \quad \forall \alpha, \beta \end{aligned}$$

$$\begin{aligned} r(X)|\Psi\rangle = \sum_{\nu} r_{\nu} X_{\nu} |\Psi\rangle = 0 &\Rightarrow \sum_{\nu} r_{\nu} \langle \Psi | X_{\alpha}^{\dagger} X_{\nu} | \Psi \rangle = 0 \quad \forall \alpha \\ &\sum_{\nu} r_{\nu} {}^kM_{\alpha, \nu} = 0 \quad \forall \alpha \end{aligned}$$

Relaxation of degree k

We define the relaxation R_k of degree k as the following SDP:

$$\begin{array}{ll}\min_{^kM} & \sum_{\alpha} p_{\alpha} \, ^kM_{0\alpha} \\ \text{s.t.} & ^kM \succeq 0 \\ & \sum_{\nu} q_{\nu} \, ^kM_{\alpha,\nu\beta} = 0 \quad \forall \alpha, \beta \\ & \sum_{\nu} r_{\nu} \, ^kM_{\alpha,\nu} = 0 \quad \forall \alpha\end{array}$$

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Let p_k be the solution of R_k
and p_{\star} be the solution of the original NCPO problem, then

$$p_k \leq p_{k+1} \leq \dots \leq p_{\star}$$

Results

- If $q(X) = 0$ implies that $\|X_i\|_2 \leq C$:

$$\lim_{k \rightarrow \infty} p_k = p_*$$

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Results

- If $q(X) = 0$ implies that $\|X_i\|_2 \leq C$:

$$\lim_{k \rightarrow \infty} p_k = p_\star$$

- In practice, we observe very often that R_2 , R_3 , or R_4 already yield the optimal solution p_\star
- Stopping criterion: if $\text{rank } {}^k M = \text{rank } {}^{k-1}({}^k M)$:
 - $p_k = p_\star$
 - the optimal $|\Psi\rangle$ and X live in a Hilbert space H of $\dim(H) = \text{rank } {}^k M$.
 - We have a procedure to reconstruct the optimal $|\Psi\rangle$ and X .

Method is related to other mathematical techniques

- $$\begin{array}{ll} \min_{H, X, |\Psi\rangle} & \langle \Psi | p(X) | \Psi \rangle \\ \text{s.t.} & q(X) = 0 \\ & X_i X_j - X_j X_i = 0 \end{array} \quad \Leftrightarrow \quad \begin{array}{ll} \min_{x \in \mathbb{R}^n} & p(x) \\ \text{s.t.} & q(x) = 0 \end{array}$$

→ Recover the SDP method for scalar polynomial optimization of [Parrilo 00, Lasserre 01].

- Dual formulation of the SDP relaxations linked to the theory of SOS decompositions of positive polynomials [Putinar 93, Helton and McCullough 04].

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Electronic energy as NCPO

The ground-state energy of N e^- that can occupy R orbitals is the solution to the following NCPO

with operator variables $\{a_1, \dots, a_r, a_1^\dagger, \dots, a_r^\dagger\}$

$$\begin{aligned} \min \quad & \sum_{ij} H_{ij} \langle \Psi | a_i^\dagger a_j | \Psi \rangle + \sum_{ijkl} H_{ijkl} \langle \Psi | a_i^\dagger a_j^\dagger a_k a_l | \Psi \rangle \\ \text{s.t.} \quad & \{a_i, a_j\} = 0 \\ & \{a_i^\dagger, a_j^\dagger\} = 0 \\ & \{a_i^\dagger, a_j\} - \delta_{ij} = 0 \\ & \sum_i (a_i^\dagger a_i - N) |\Psi\rangle = 0 \end{aligned}$$

Example: relaxation of degree 2

Remember: matrix kM is built as overlap matrix of set of vectors $\{X_\alpha|\Psi\rangle\}$ where X_α is at most of degree k .

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In our case 2M is thus built on set of vectors

$$\{|\Psi\rangle, a_i|\Psi\rangle, a_i^\dagger|\Psi\rangle, a_i a_j|\Psi\rangle, a_i^\dagger a_j^\dagger|\Psi\rangle, a_i^\dagger a_j|\Psi\rangle, a_i a_j^\dagger|\Psi\rangle\}$$

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General property

It is sufficient to consider set of vectors $\{X_\alpha|\Psi\rangle\}$ that are linearly independent under the constraints $q(X) = 0, r(X)|\Psi\rangle = 0$

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General property

It is sufficient to consider set of vectors $\{X_\alpha|\Psi\rangle\}$ that are linearly independent under the constraints $q(X) = 0$, $r(X)|\Psi\rangle = 0$

Thus since $a_i^\dagger a_j + a_j a_i^\dagger - \delta_{ij} = 0$ and $\sum_i (a_i^\dagger a_i - N) |\Psi\rangle = 0$, it is sufficient to consider set

$$\{a_i|\Psi\rangle, a_i^\dagger|\Psi\rangle, a_i a_j|\Psi\rangle, a_i^\dagger a_j^\dagger|\Psi\rangle, a_i^\dagger a_j|\Psi\rangle\}$$

Example: relaxation of degree 2

$${}^2M = \begin{array}{c|ccccc} & a_i \rangle & a_i^\dagger \rangle & a_i a_j \rangle & a_i^\dagger a_j^\dagger \rangle & a_i^\dagger a_j \rangle \\ \hline \langle a_k^\dagger & \langle a_k^\dagger a_i \rangle & \langle a_k^\dagger a_i^\dagger \rangle & \langle a_k^\dagger a_i a_j \rangle & \langle a_k^\dagger a_i^\dagger a_j^\dagger \rangle & \langle a_k^\dagger a_i^\dagger a_j \rangle \\ \langle a_k & & \langle a_k a_i^\dagger \rangle & \langle a_k a_i a_j \rangle & \langle a_k a_i^\dagger a_j^\dagger \rangle & \langle a_k a_i^\dagger a_j \rangle \\ \langle a_k^\dagger a_l^\dagger & & & \langle a_k^\dagger a_l^\dagger a_i a_j \rangle & \langle a_k^\dagger a_l^\dagger a_i^\dagger a_j^\dagger \rangle & \langle a_k^\dagger a_l^\dagger a_i^\dagger a_j \rangle \\ \langle a_k a_l & & & & \langle a_k a_l a_i^\dagger a_j^\dagger a_j \rangle & \langle a_k a_l a_i^\dagger a_j \rangle \\ \langle a_k^\dagger a_l & & & & & \langle a_k^\dagger a_l a_i^\dagger a_j \rangle \end{array}$$

Example: relaxation of degree 2

Linear constraints on kM which follow from

- the conditions $q(X) = \sum_{\nu} q_{\nu} X_{\nu} = 0$:

$$\sum_{\nu} q_{\nu} \langle X_{\alpha}^{\dagger} X_{\nu} X_{\beta} \rangle = 0 \quad \forall \alpha, \beta$$

For instance: $q(X) = \{a_i^{\dagger}, a_j\} - \delta_{ij}$, $X_{\alpha} = a_k^{\dagger}$, $X_{\beta} = a_l$

$$\Rightarrow \langle a_k^{\dagger} a_i^{\dagger} a_j a_l \rangle + \langle a_k^{\dagger} a_j a_i^{\dagger} a_l \rangle - \delta_{ij} \langle a_k^{\dagger} a_l \rangle = 0$$

- and the conditions $r(X)|\Psi\rangle = \sum_{\nu} r_{\nu} X_{\nu} |\Psi\rangle = 0$:

$$\sum_{\nu} q_{\nu} \langle X_{\alpha}^{\dagger} X_{\nu} \rangle = 0 \quad \forall \alpha$$

For instance: $r(X) = \sum_i (a_i^{\dagger} a_j - N)$, $X_{\alpha} = a_k a_l$:

$$\Rightarrow \sum_i \langle a_k a_l a_i^{\dagger} a_i \rangle - N \langle a_k a_l \rangle = 0$$

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Example: relaxation of degree 2

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+ Linear relations between 1D , 1Q , 2D , 2Q , 2G

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+ Linear relations between 1D , 1Q , 2D , 2Q , 2G

We recover the same SDP than in the RDM method

Electronic energy as NCPO

Similarly, relaxations of higher-degree of the problem

$$\begin{aligned} \min \quad & \sum_{ij} H_{ij} \langle \Psi | a_i^\dagger a_j | \Psi \rangle + \sum_{ijkl} H_{ijkl} \langle \Psi | a_i^\dagger a_j^\dagger a_k a_l | \Psi \rangle \\ \text{s.t.} \quad & \{a_i, a_j\} = 0 \\ & \{a_i^\dagger, a_j^\dagger\} = 0 \\ & \{a_i^\dagger, a_j\} - \delta_{ij} = 0 \\ & \sum_i (a_i^\dagger a_i - N) |\Psi\rangle = 0 \end{aligned}$$

correspond to implementing higher-order positivity
 N -representability constraints of the RDM method.

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correspond to implementing higher-order positivity
 N -representability constraints of the RDM method.

Note that here

$$p_2 \leq p_3 \leq \dots p_N = p_\star$$

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Quantum violation of Bell inequalities

Original motivation for our method [*PRL 07, NJP 08*]

$$\begin{array}{ll} \min & \sum_{abxy} c_{abxy} \langle \Psi | E_a^x E_b^y | \Psi \rangle \\ \text{s.t.} & E_a^x E_{a'}^x = \delta_{aa'} E_a^x \quad \text{and} \quad \sum_a E_a^x = I \\ & E_b^y E_{b'}^y = \delta_{bb'} E_b^y \quad \text{and} \quad \sum_b E_b^y = I \\ & [E_a^x, E_b^y] = 0 \end{array}$$

The dimension of H is not bounded

Quantum violation of Bell inequalities

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The dimension of H is not bounded

Has been applied to 241 Bell inequalities in [*Pal, Vertesi 08*]

Optimal solution obtained for 221 Bell inequalities

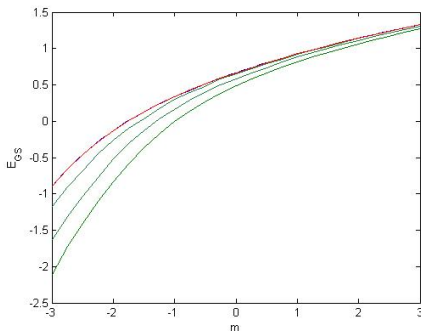
For the remaining 20 ones: gap between our LB and best known UB is of order 10^{-4} .

Other applications

- Security of device-independent cryptography ✓
- Continuous variable system ✓

Particle in a double-well potential:

$$\begin{array}{ll} \min & \frac{p^2}{2} + \frac{x^2}{2} + mx^4 \\ \text{s.t.} & [x, p] = i \end{array}$$



Conclusion

Similar techniques than the ones in the SDP RDM method allow to solve a broad class of non-commutative polynomial optimization problems.

Interest of our method for solving NCPO:

- Flexible and works on different problems
- Yields lower-bounds
- Does not rely on symmetries (good for quantum chemistry)
- Allows to deal with infinite Hilbert space without truncating the Hilbert space