

# Reduction theory and coding of geodesics on the modular surface

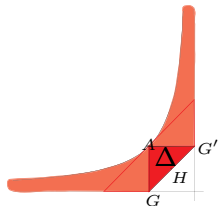
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# A family of maps $f_{a,b}$

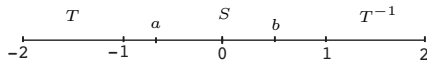
suggested for consideration by Don Zagier.

Let  $\mathcal{P} = \{(a, b) \mid a \leq 0 \leq b, b - a \geq 1, -ab \leq 1\} \supset$   
 $\Delta = \{-1 \leq a \leq 0 \leq b \leq 1, b - a \geq 1\}$ ;



Let  $f_{a,b} : \mathbb{R} \rightarrow \mathbb{R}$  be defined as

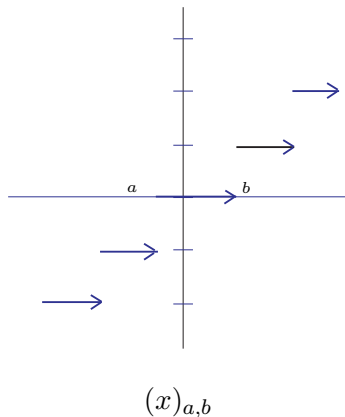
$$f_{a,b}(x) = \begin{cases} x + 1 & \text{if } x < a \\ -\frac{1}{x} & \text{if } a \leq x < b \\ x - 1 & \text{if } x \geq b. \end{cases}$$



$Tx = x + 1, Sx = -\frac{1}{x},$   
 $T^{-1}x = x - 1$   
generators of  $SL(2, \mathbb{Z})$

The map  $f_{a,b}$  defines what we call  **$(a, b)$ -continued fractions** using a generalized **integral part** function  $(x)_{a,b}$ :

# $(a, b)$ -continued fractions: joint work with I. Ugarcovici



## Theorem

→ If  $(a, b) \in \Delta$ , then any  $x$  can be expressed uniquely in the form

$$x = n_0 - \frac{1}{n_1 - \frac{1}{\ddots}}$$

where  $n_0 = (x)_{a,b}$ ,  $x_1 = -\frac{1}{x-n_0}$  and  $n_{i+1} = (x_{i+1})_{a,b}$ ,  $x_{i+1} = -\frac{1}{x_i - n_i}$ , i.e.  
 $r_k = (n_0, n_1, \dots, n_k)_{a,b} = \frac{p_k}{q_k} \rightarrow x$ .

# The natural extension map

$$\text{Let } F_{a,b} : \bar{\mathbb{R}}^2 \rightarrow \bar{\mathbb{R}}^2, F_{a,b}(x, y) = \begin{cases} (x + 1, y + 1) & \text{if } y < a \\ (-\frac{1}{x}, -\frac{1}{y}) & \text{if } a \leq y < b \\ (x - 1, y - 1) & \text{if } y \geq b \end{cases}$$

be the (natural) extension map of  $f_{a,b}$ .

Numerical experiments led Zagier to conjecture that  $F_{a,b}$  possesses a global attractor set  $D_{a,b} = \cap_{n=0}^{\infty} F_{a,b}^n(\bar{\mathbb{R}}^2)$  with finite rectangular structure on which it is essentially bijective, and every point of the plane is mapped to  $D_{a,b}$  after finitely many iterations of  $F_{a,b}$ .

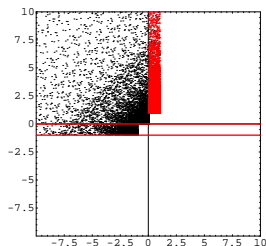
If one identifies a geodesic of the upper half-plane with a pair of real numbers  $(u, w) \in \bar{\mathbb{R}}^2$ ,  $u \neq w$  — its endpoints, then  $F_{a,b}$  maps geodesics to geodesics, and the existence of an attractor for  $F_{a,b}$  corresponds to a reduction of geodesics, hence we perceive  $F_{a,b}$  as a reduction map.

# Zagier's Reduction theory conjecture

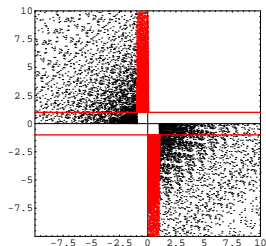
**RTC:** For any  $(a, b) \in \mathcal{P}$   $F_{a,b}$  has an attractor  $D_{a,b}$  with finite rectangular structure with the following additional property:

- for every  $(u, w) \in \bar{\mathbb{R}}^2 \exists N \geq 0$  s.t.  $F_{a,b}^N(u, w) \in D_{a,b}$ .

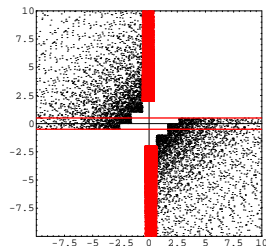
RTC holds for three classical cases:



$A_{-1,0}$   
minus c.f.

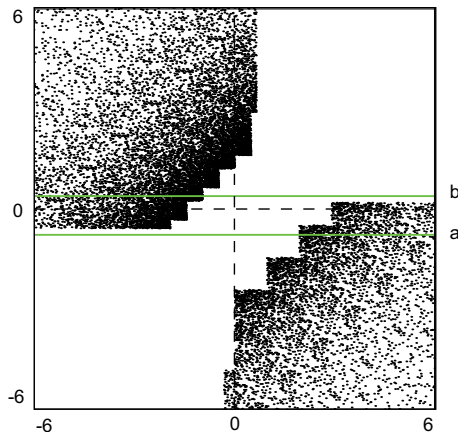


$A_{-1,1}$   
alternating c.f.



$A_{-1/2,1/2}$   
nearest integer c.f.

# Mathematica experimentation: Stage 1



A generic example  $a = -\frac{4}{5}, b = \frac{2}{5}$   
attractor obtained by iterating  
random points

# Definition of attractor via trapping region

## Definition

$\Theta_{a,b} \subset \bar{\mathbb{R}}^2$  is a **trapping region** for the reduction map  $F_{a,b}$  if

- for every pair  $(x, y) \in \bar{\mathbb{R}}^2$ ,  $\exists N > 0$  such that  $F_{a,b}^N(x, y) \in \Theta_{a,b}$ ;
- $F_{a,b}(\Theta_{a,b}) \subset \Theta_{a,b}$ .

## Definition

We define the **attractor** starting with the trapping region:

$$D_{a,b} = \bigcap_{n=0}^{\infty} D_n, \text{ where } D_n = \bigcap_{i=0}^n F_{a,b}^i(\Theta_{a,b}).$$

# Existence of a trapping region

## Theorem

The region  $\Theta_{a,b}$  is bounded by step-functions

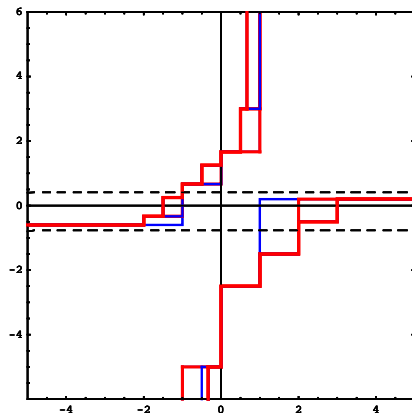
$$u(x) = \begin{cases} b-1 & \text{if } x \leq -1 \\ \min(-1/(b-1) - 1, -1/a) & \text{if } -1 \leq x \leq 0 \text{ and} \\ -1/(b-1) & \text{if } 0 \leq x \leq 1, \end{cases}$$
$$\ell(x) = \begin{cases} -1/(a+1) & \text{if } -1 \leq x \leq 0 \\ \max(-1/(a+1) + 1, -1/b) & \text{if } 0 \leq x \leq 1 \\ a+1 & \text{if } x \geq 1 \end{cases}$$

is the trapping region for the reduction map  $F_{a,b}$ .

**Observation:**  $\ell(x)$  and  $u(x)$  take the initial values of the orbits of  $a$  and  $b$ .



## Mathematica experimentation: Stage 2



Obtaining attractor from the trapping region

**Observation:** horizontal boundary levels belong to the orbits of  $a$  and  $b$ .

## Orbits of $a$ and $b$

The map  $f_{a,b} = f$  is discontinuous for  $x = a, b$ , however, we can look at two maps, one **on the left** and one **on the right** of  $x = a, b$ , and the corresponding **split orbits**:

$$\mathcal{O}_\ell(a) = \{Ta, fTa, f^2Ta, \dots\}, \quad \mathcal{O}_u(a) = \{Sa, fSa, f^2Sa, \dots\}$$

and

$$\mathcal{O}_\ell(b) = \{Sb, fSb, f^2Sb, \dots\}, \quad \mathcal{O}_u(b) = \{T^{-1}b, fT^{-1}b, f^2T^{-1}b, \dots\}.$$

**Observation:** horizontal segments of the upper boundary of the attractor belong to  $\mathcal{O}_u(a)$  and  $\mathcal{O}_u(b)$ , and of the lower boundary - to  $\mathcal{O}_\ell(a)$  and  $\mathcal{O}_\ell(b)$ , hence we will refer to  $\mathcal{O}_\ell(a)$  and  $\mathcal{O}_\ell(b)$  as **lower**, and to  $\mathcal{O}_u(a)$  and  $\mathcal{O}_u(b)$  as **upper** orbits of  $a$  and  $b$ .

# The cycle property

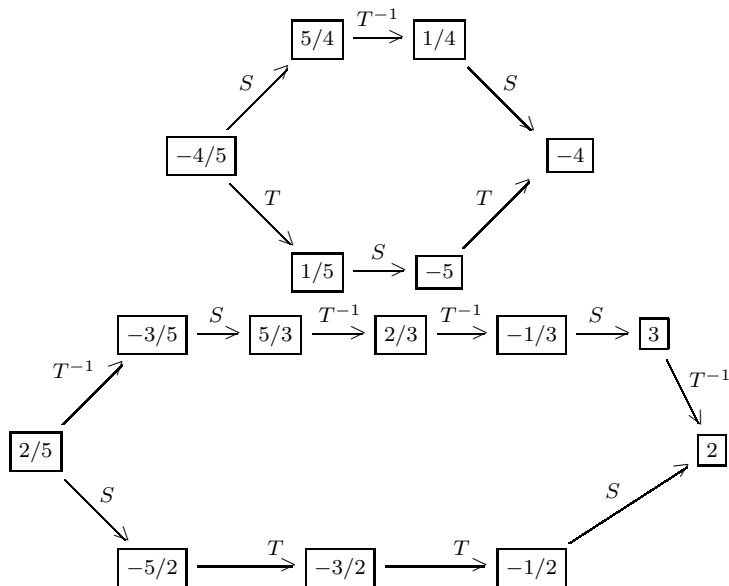
Our initial experiments showed that the following pattern was prevalent and generic:

## Definition

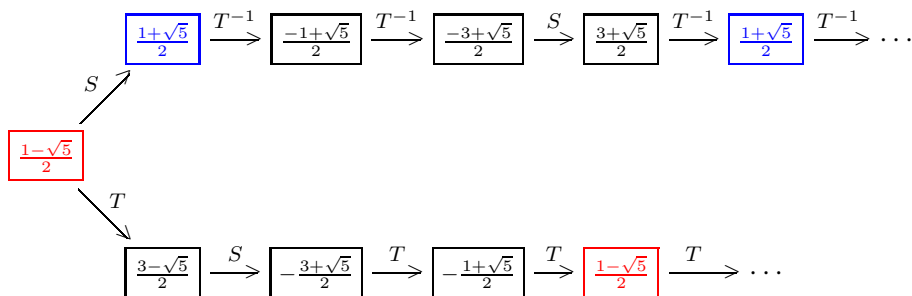
- We say that  $a$  (resp.,  $b$ ) has the **cycle property** if the upper and lower orbits meet forming a cycle.
- If the product over the cycle equals the identity we say that the cycle property is **strong**, otherwise, the cycle property is **weak**.

Another pattern that we noticed was periodicity of the orbits.

A generic example:  $a = -\frac{4}{5}, b = \frac{2}{5}$  - the cycle property!



Periodic expansion:  $a = \frac{1-\sqrt{5}}{2}, b = \frac{-1+\sqrt{5}}{2}$



# The finiteness condition

## Definition

We say that  $(a, b)$  satisfies the **finiteness condition** if  $a$  and  $b$  either have the cycle property or their expansions are eventually periodic.

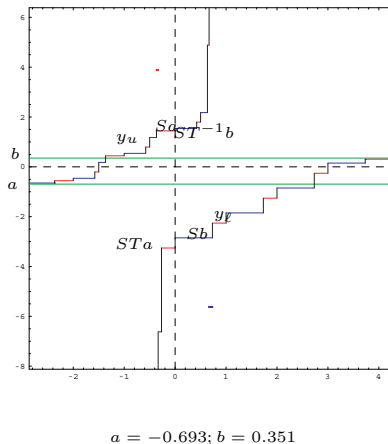
Thus the sets  $\mathcal{L}_a, \mathcal{U}_a, \mathcal{L}_b$ , and  $\mathcal{U}_b$ , called the **truncated orbits**, are finite.

$$\mathcal{L}_a = \begin{cases} \mathcal{O}_\ell(a) & \text{if } a \text{ has periodic expansion} \\ \text{lower part of } a\text{-cycle} & \text{if } a \text{ has the cycle property} \\ \text{lower part of } a\text{-cycle} \cup \{0\} & \text{if } a \text{ has weak cycle property,} \end{cases}$$
$$\mathcal{U}_a = \begin{cases} \mathcal{O}_u(a) & \text{if } a \text{ has periodic expansion} \\ \text{upper part of } a\text{-cycle} & \text{if } a \text{ has the cycle property} \\ \text{lower part of } a\text{-cycle} \cup \{0\} & \text{if } a \text{ has weak cycle property,} \end{cases} \quad \text{and,}$$

similarly,  $\mathcal{L}_b$  and  $\mathcal{U}_b$ .

**Observation:** Only truncated orbits appear as horizontal levels of the attractor, ends of the cycles appear only if  $= 0$ , i.e. if the cycle is weak.

# Mathematica experimentation: Stage 3



The boundary of the attractor is **computed** from the data  $(a, b)$ : the lower boundary consists of all levels in  $\mathcal{L}_a \cup \mathcal{L}_b$  and the upper - of all levels in  $\mathcal{U}_a \cup \mathcal{U}_b$ . The  $x$ -levels are solutions of the (overdetermined) system of fractional-linear equations that is consistent and equivalent to the system of two equations at consecutive levels  $y_u < S_a$  and  $y_l > S_b$ .

The system was solved and picture drawn by computer.

# Finiteness condition implies finite rectangular structure

## Definition

We say that a proper subset of  $\bar{\mathbb{R}}^2$  has **finite rectangular structure** if it consists of two (or one, in degenerate cases) connected components bounded by non-decreasing step-functions with finitely many steps.

The main result of our work is the following:

## Theorem (FRS)

*If  $(a, b) \in \Delta$  satisfies the **finiteness condition**, the attractor set  $D_{a,b} \subsetneq \bar{\mathbb{R}}^2$  has finite rectangular structure, and  $F_{a,b} : D_{a,b} \rightarrow D_{a,b}$  is a bijection except for some images of the boundary of  $D_{a,b}$ .*

- **Step 1:** construction of a set  $A_{a,b}$  where  $F_{a,b}$  is bijective by starting with **3 connected levels**  $STa < Sb < y_\ell$  and proving that all levels in  $\mathcal{L}_a \cup \mathcal{L}_b$  are **connected**, and similarly of upper levels.
- **Step 2:**  $D_{a,b} = A_{a,b}$ .  $A_{a,b} \subset D_n$  for all  $\forall n$  used for connectedness.



# The set of exceptions $\mathcal{E}$ to the finiteness condition

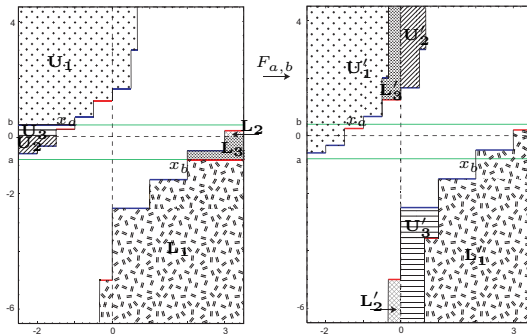
## Theorem

- For all  $(a, b) \in \Delta \setminus \{ \text{the diagonal } b = a + 1 \}$  the finiteness condition is satisfied, hence Theorem FRS holds.
- The exceptional set  $\mathcal{E} \in \Delta$  is a uncountable Lebesgue measure 0 on the diagonal  $b = a + 1$ .
- The finiteness condition is necessary for finite rectangular structure of the attractor, i.e. if  $(a, b) \in \mathcal{E}$  s.t.  $\mathcal{U}_a \cup \mathcal{U}_b \cup \mathcal{L}_a \cup \mathcal{L}_b$  is infinite, either the attractor  $D_{a,b}$  is disconnected, or it consists of two connected components whose boundary functions take all values of  $\mathcal{U}_a \cup \mathcal{U}_b \cup \mathcal{L}_a \cup \mathcal{L}_b$ , hence are not step-functions with finitely many steps.

# Reduction theory conjecture

## Theorem

If  $a, b \in \Delta$  and both have the **strong cycle property**, the RTC holds.

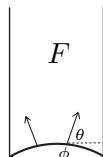
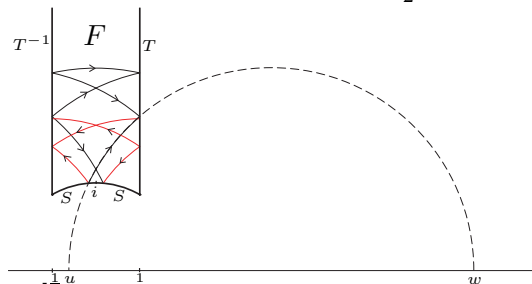


Observation: levels corresponding to the ends of the cycles are inside the attractor, hence any boundary component is mapped inside under some iteration of  $F_{a,b}$ .

# Geometric code for the modular surface

$M = PSL(2, \mathbb{Z}) \backslash \mathcal{H}$  – modular surface

$F = \{z \in \mathcal{H} \mid |z| \geq 1, |\operatorname{Re} z| \leq \frac{1}{2}\}$



The cross-section  $B$

$$T(z) = z + 1$$

$$S(z) = -1/z$$

$$[T, T, T, T, S, T^{-1}, T^{-1}, T^{-1}, S]$$

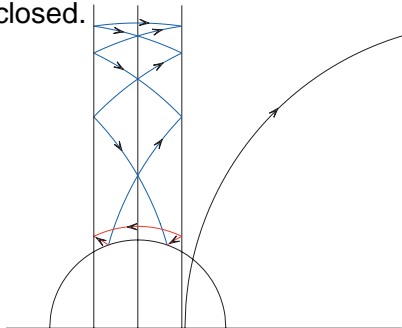
$$[4, -3] \Leftarrow \text{geometric code}$$

- Any geodesic not going to the cusp in either direction has a **geometric code**  $[\dots n_{-2}, n_{-1}, n_0, n_1, n_2, \dots]$ ,  $n_i \neq 0$ .
- Any closed geodesic passing through  $B$  - axis of  $A \in SL(2, \mathbb{Z})$  - has a periodic code  $[n_1, n_2, \dots, n_m]$  and  $A = T^{n_1} S T^{n_2} S \dots T^{n_m} S$ .
- Left shift  $\sigma$  of the sequence corresponds to the first return to the cross-section  $B$ .

# Which geometric codes are realized?

Denote the set of all admissible geometric codes by  $X$ .

$X \subset \mathcal{N}^{\mathbb{Z}}$ , where  $\mathcal{N} = \{n \in \mathbb{Z}, |n| \geq 1\}$  – alphabet, is  $\sigma$ –invariant and closed.



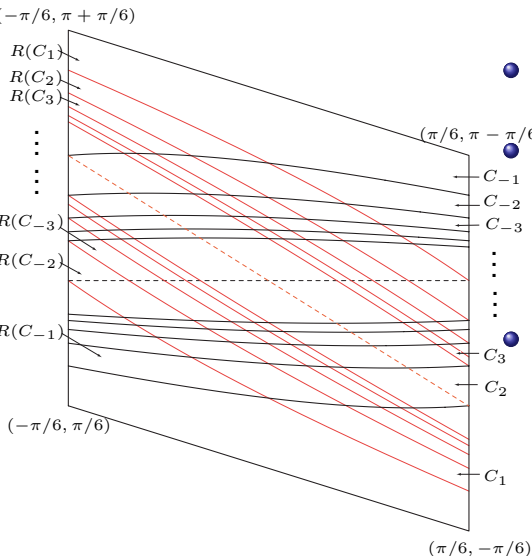
Not all sequences of non-zero integers are realized as geometric codes of geodesics on  $M$ .

## Example

$[8, 2]$  is not a geometric code since the geometric code of the axis of  $T^8ST^2S$  is  $[6, -2]$  [K].

Mathematica experimentation: drawing geodesics on the modular surface.

# Mathematica experimentation: the geometric partition



Geometric partition and its image under the return map  $R$

- Partition:**  $B = \sqcup_{n \in \mathcal{N}} C_n$ ,  
 $C_n = \{v \in B \mid n_1(v) = n\}$ .
- If  $n_1(v) = n$ ,  $n_2(v) = m$  for some vector  $v \in B$ , then  $R(C_n) \cap C_m \neq \emptyset$ . Therefore, 2 cannot be followed by 1, 2, 3, 4, 5.
- $R(C_n)$  and  $C_m$  intersect **transversally** if and only if  $|n|, |m| \geq 2$ , and  $|1/n + 1/m| \leq 1/2$ :  
**geometrically Markov geodesics**

# A class of admissible geometric codes

## Theorem ([KU])

Any bi-infinite sequence  $[\dots, n_{-1}, n_0, n_1, n_2, \dots]$  such that  $\left| \frac{1}{n_i} + \frac{1}{n_{i+1}} \right| \leq \frac{1}{2}$  for  $i \in \mathbb{Z}$ , is realized as a geometric code of a geodesic on  $M$ .

We denote this set by  $X_M$  – **geometrically Markov codes** – it is

- a maximal, 1-step countable topological Markov chain in  $X$ ;
- the maximal if  $X_M$  is symmetric (i.e. given by a symmetric transition matrix).

## Theorem ([KU, KU1])

Complexity of the geometric code: the space  $X$  of geometric codes is not a finite-step topological Markov chain.

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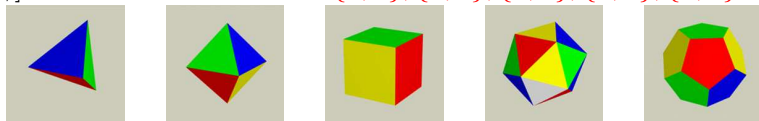
[KU] S. Katok and I. Ugarcovici, *Geometrically Markov geodesics on the modular surface*, Moscow Math. J., **1** (2005), 135-155.

[KU1] S. Katok and I. Ugarcovici, *Symbolic dynamics for the modular surface and beyond*, Bull. Am. Math. Soc., **44** (2007) 87-132

# Positive geodesics

This class includes a class of codes of **positive geodesics** found earlier in [GK]:  $[\gamma] = [\dots, n_{-1}, n_0, n_1, \dots]$ , where  $\frac{1}{n_i} + \frac{1}{n_{i+1}} \leq \frac{1}{2}$  for all  $i$ , i.e.

$[\gamma]$  does not contain **2** and  **$\{3, 3\}, \{3, 4\}, \{4, 3\}, \{3, 5\}, \{5, 3\}$**



- all segments comprising  $\gamma$  in  $F$  are **positively** (clockwise) oriented.
- the geometric code of  $\gamma$  coincides with its **arithmetic code** given by minus continued fractions ( $a = -1, b = 0$ ). It is obtained by juxtaposing expansions of  $w = [n_0, n_1, \dots]$  and  $1/u = [n_{-1}, n_{-2}, \dots]$ :

$$[\gamma] = [\dots n_{-2}, n_{-1}, n_0, n_1, n_2, \dots],$$

where the geodesic from  $u$  to  $w$  is reduced, i.e.  $0 < u < 1, w > 1$ .

[GK] B. Gurevich and S. Katok, *Arithmetic coding and entropy for the positive geodesic flow on the modular surface*, Moscow Math. J., 1, no. 4 (2001), 569–582.

# Coding via $(a, b)$ -continued fractions

$(a, b)$ -continued fraction expansion can be used for coding if

- it satisfies the **RTC**, and
- has a **dual** (or is self-dual).

## Definition

The  $(a, b)$ -expansion has a **dual** if the reflection of  $A_{a,b}$  in the line  $w = -u$  is an attractor for some  $(a', b')$ -expansion.

If  $(a', b') = (a, b)$ , the  $(a, b)$ -expansion is called **self-dual**.

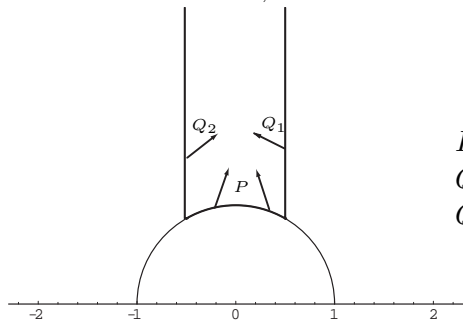
## Definition

A geodesic in  $\mathcal{H}$  from  $u$  to  $w$  is called  $(a, b)$ -*reduced* if  $(u, w) \in \Lambda_{a,b} = F_{a,b}(D_{a,b} \cap \{a \leq y \leq b\})$ .



# Coding via $(a, b)$ -continued fractions

If  $(a, b) \in \Delta$ , the  $(a, b)$ -reduced geodesic from  $u$  to  $w$  intersects the unit half-circle, and let  $C_{a,b} = P \cup Q_1 \cup Q_2$ :



$P$ :  $\gamma$  is  $(a, b)$ -reduced

$Q_1$ :  $TS(\gamma)$  is  $(a, b)$ -reduced

$Q_2$ :  $T^{-1}S(\gamma)$  is  $(a, b)$ -reduced

- $C_{a,b}$  is a **cross-section**, i.e. a surface inside  $SM$  that every geodesic visits infinitely many times.
- $\Lambda_{a,b}$  is a  $(u, w)$ -parametrization of  $C_{a,b}$ .

# Coding via $(a, b)$ -continued fractions

- **Reduction theory**  $\Rightarrow$  every geodesic in  $\mathcal{H}$  is  $PSL(2, \mathbb{Z})$ -equivalent to a reduced one.
- **Reduced geodesic**  $\gamma$  in  $\mathcal{H}$  from  $u$  to  $w \Leftrightarrow (u, w) \in \Lambda_{a,b}$   
 $\Rightarrow (\gamma)$  - **arithmetic code** obtained by juxtaposing of expansion for  $w = (n_0, n_1, \dots)_{a,b}$  and (dual) expansion for  $1/u = (n_{-1}, \dots)_{a',b'}$

$$(\gamma) = (\dots n_{-2}, n_{-1}; n_0, n_1, n_2, \dots)$$

- The left shift  $\sigma$  corresponds to the first return to the cross-section  $C_{a,b}$
- $PSL(2, \mathbb{Z})$ -invariance is proved via the cross-section.
- Closed geodesics have periodic coding sequences:  
 $w = (n_1, \dots, n_m)_{a,b}, \frac{1}{u} = (n_m, \dots, n_1)_{a',b'}$ .

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[GK] B. Gurevich and S. Katok, *Arithmetic coding and entropy for the positive geodesic flow on the modular surface*, Moscow Math. J., 1, no. 4 (2001), 569–582.