

# Quantum Error Correction II

## Robust Quantum Information Processing

Manny

- Threshold theorems.
- Requirements for scalability.
- Physical noise models.
- Methods for scalability.

# References

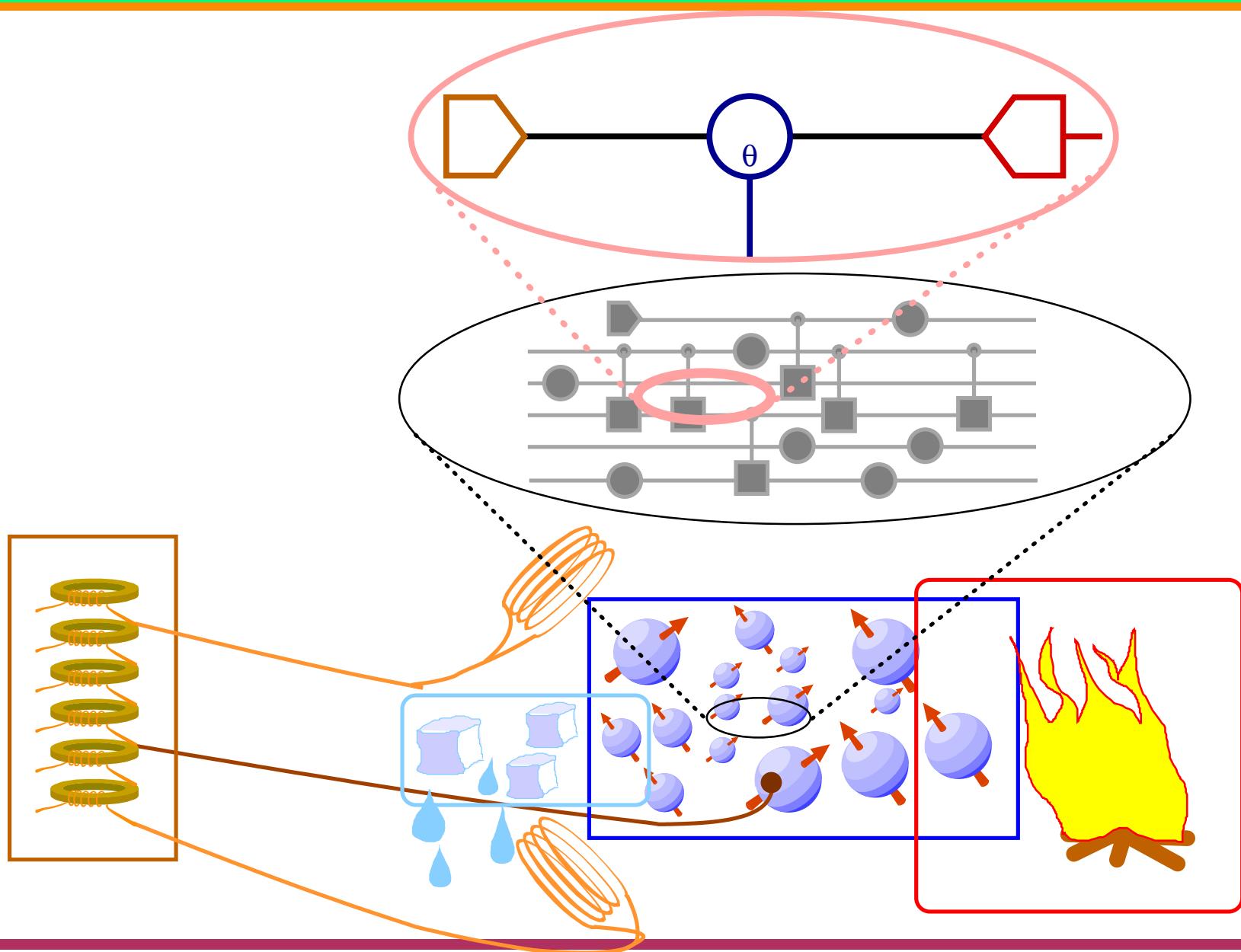
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- Prehistory:
  - Quantum Zeno effect: Misra&Sudarshan 1977 [18].
  - Deutsch 1993, Barenco&*al.* 1996 [3].
- Discovery and Theory:
  - Shor 1995 [21], Steane 1995 [23].
  - Bennett&DiVincenzo&Smolin&Wootters 1996 [4], Knill&Laflamme 1996 [12].
  - Calderbank&Shor 1996 [6], Gottesman 1996 [8], Calderbank&Rains&Shor&Sloane 1997 [5].
- Fault tolerance and threshold accuracies:
  - Shor 1996 [22], Kitaev 1997 [11].
  - Aharonov&Ben-Or 1996 [1, 2], Knill&Laflamme&Zurek 1996 [15], Gottesman&Preskill 1997 [9, 20], Dür&Briegel&*al.* 1999 [7].
  - Gottesman 1997 [9], Gottesman&Chuang 1999 [10]
- Toward subsystems:
  - Quasi-particles . . .
  - Zanardi&Rasetti 1997 [27], Lidar&Chuang&Whaley 1998 [17].
  - Viola&Knill&Lloyd 1998 [25, 24, 26].
  - Knill&Laflamme 1996 [12], Knill&Laflamme&Viola 2000 [14].

General reference: (M)ike, Ch. 10.

Nielsen&Chuang 2001 [19]

# QIP System Overview



# The Threshold Theorems

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- The Accuracy Threshold Theorem:  
Assume the *requirements for scalable computing*. If the error per gate (including “no-op”) is less than a threshold, then it is possible to efficiently quantum compute arbitrarily accurately.
  - Error Thresholds:
    - Worst case  $> 10^{-6}$ .
    - Estimate  $> 10^{-4}$ .
    - Communication  $> 10^{-2}$ .
    - Erasure  $> 10^{-2}$ .
    - Z-measurement  $\geq .5$ .

Shor 1996[22], Kitaev 1996[11], Aharonov&Ben-Or 1996[1], Knill& al. 1996[16].

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# Sufficient Requirements for Scalability

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Realize QIP task with  $M$  gates and noops,  $N$  qubits.

- Systems:
  - polylog( $M$ ) $N$  independent  $\geq 2$  state subsystems.
- Preparation:
  - Can reset subsystem to  $|0\rangle$  anytime.
- Control:
  - Universal two-system gates can be applied.
- Parallelism:
  - Can apply gates in parallel.
  - Or: Perfect long-term memory.
- Noise:
  - Sufficiently weak.
  - Quasi-independent.

# Types of Noise

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- Quantum noise in QIP:

Any unwanted effect in quantum systems.

- By temporal behavior.

- Step-wise independent, discrete.
- Markovian.
- Relaxation: Decay to equilibrium.
  - Depolarization.
  - Dissipation, thermal relaxation.
  - Decoherence: Loss of phase.
- Stationary.

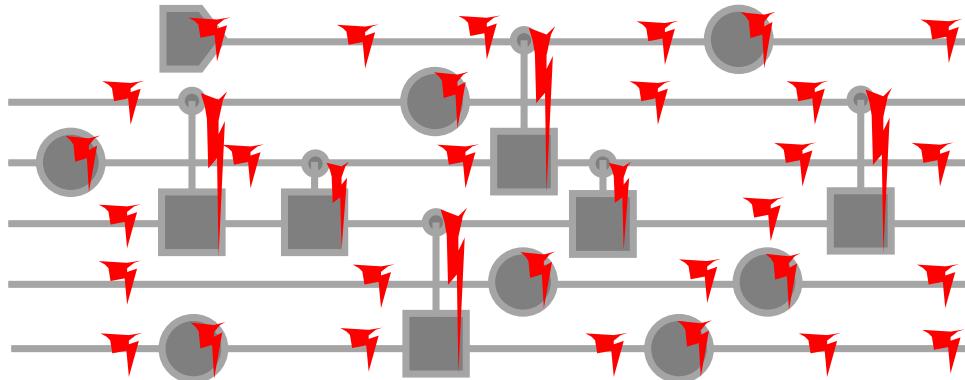
- By spatial behavior.

- Local or factorizable.
- Symmetric.
- Linear.

- By origin.

- Thermal.
- Control errors, faults.
  - Miscalibration.
  - Over/under-rotation.
- Stray fields or inhomogeneity.

# Noise Analysis



- Error locations:

For each gate  $U_i$  including  $U_i = \text{noop}$ :

$$U_{i\text{actual}} = E_i U_{i\text{intended}}$$

Locality:  $E_i$  acts on the same qubits as  $U_i$ .

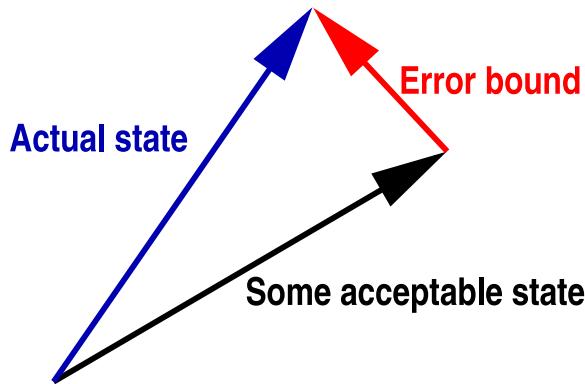
- Error expansion with “environment”:

$$|\psi\rangle \rightarrow \sum_e |e\rangle_E E_{e,n} U_n \cdots E_{e,2} U_2 E_{e,1} U_1 |\psi\rangle$$

The  $|e\rangle_E$  need not be orthogonal.

# A Measure of Noise

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- Wavefunction:

r.h.s. states need not be normalized

$$|\text{output}\rangle = |\text{acceptable}\rangle + |\text{error}\rangle$$

- Error probability is bounded by  $\langle \text{error} | \text{error} \rangle$ .

- Density operator:

Error operator need not be a state.

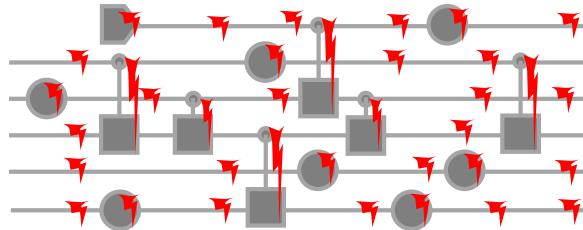
$$\rho_{\text{out}} = \rho_{\text{acceptable}} + \rho_{\text{error}}$$

- Error probability is bounded by  $(\text{tr}|\rho_{\text{error}}| + |\text{tr}\rho_{\text{error}}|)/2$ .

- Acceptable states need not be unique.

# Quasi-independent Noise

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- Noise given by error expansion:

$$|\psi\rangle \rightarrow \sum_e |e\rangle_{\mathbb{E}} E_{e,n} U_n \cdots E_{e,2} U_2 E_{e,1} U_1 |\psi\rangle$$

- The *support* of error term  $e$  is the set of  $i$  such that  $E_{e,i}$  is not  $\mathbb{I}$ .
- *Quasi-independent* with probability  $p$  if for each  $I \subseteq \{1, \dots, n\}$  probability of errors with support  $\supseteq I$  is  $\leq p^{|I|}$ .  
Threshold:  $p < 10^{-6}$  at worst.

# Independent Noise Models

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- Systematic models.

- Only one environment state  $|e\rangle_E$ .

$$\begin{aligned} |\psi\rangle &\rightarrow |e\rangle_E V_n U_n \cdots V_2 U_2 V_1 U_1 |\psi\rangle \\ &= |e\rangle_E (\mathbb{I} + E_n) U_n \cdots (\mathbb{I} + E_2) U_2 (\mathbb{I} + E_1) U_1 |\psi\rangle \end{aligned}$$

- Examples: Frequency errors, miscalibration.

Correctable by classical control.

- Random models.

- Orthogonal environments at each site.

$$\begin{aligned} |\psi\rangle &\rightarrow (|0\rangle_{E_n} \mathbb{I} + \dots + |k_n\rangle_{E_n} E_{n,k_n}) U_n \cdots \\ &\quad (|0\rangle_{E_2} \mathbb{I} + \dots + |k_2\rangle_{E_2} E_{2,k_2}) U_2 \\ &\quad (|0\rangle_{E_1} \mathbb{I} + \dots + |k_1\rangle_{E_1} E_{1,k_1}) U_1 |\psi\rangle \end{aligned}$$

- Example: Independent thermal relaxation.

# Simple Random Models

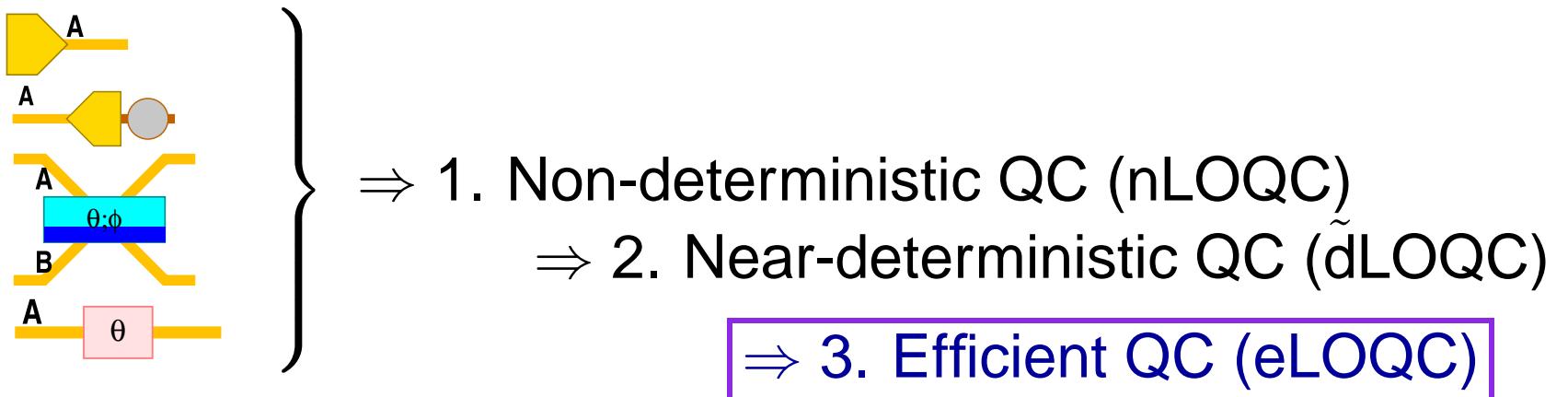
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- Random bit flips  $\mathcal{X}(p)$ .
  - One qubit:  $\sqrt{1-p}|0\rangle_E \mathbb{I} + \sqrt{p}|1\rangle_E \sigma_x$ .
  - Two qubits?  
Threshold?
- Depolarizing noise  $\mathcal{P}(p)$ .
  - One qubit:  
$$\sqrt{1 - \frac{3}{4}p} |0\rangle_E \mathbb{I} + \sqrt{\frac{1}{4}p} |1\rangle_E \sigma_x + \sqrt{\frac{1}{4}p} |2\rangle_E \sigma_y + \sqrt{\frac{1}{4}p} |3\rangle_E \sigma_z.$$
  - Two qubits?
  - Other independent models can be “twirled” into this.  
Threshold:  $p < 10^{-4}$ ?
- Erasure/detected loss  $\mathcal{L}(p)$ .
  - $\sqrt{1-p}|0\rangle_E \mathbb{I} + \sqrt{p}|1\rangle_E \mathcal{P}(1)$ .
  - E is accessible.  
Threshold:  $p < 10^{-2}$  at worst.

# eLOQC Designer Errors

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- eLOQC: Efficient linear optics quantum computation.



- Source of non-determinism:  $Z$ -measurement errors.
  - Elementary gates succeed, but...
  - ...with probability  $f$ : qubit measured.

Threshold:  $f < .5$  at worst.

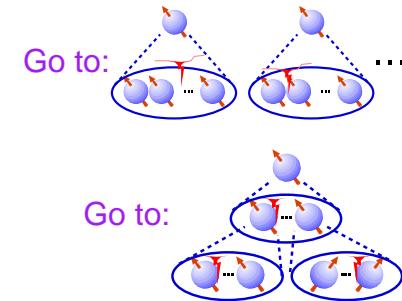
Knill&Laflamme&Milburn 2001 [13]

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# Methods for Scalability

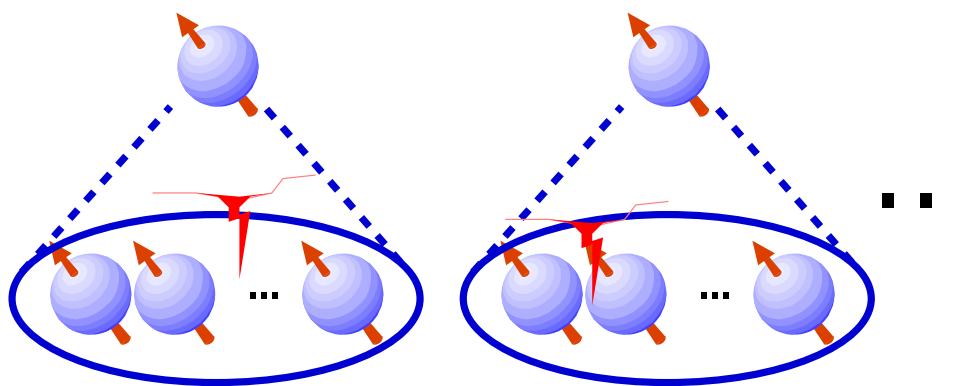
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- Error-correcting subsystems.
- Concatenation.



# Error-correcting Subsystems

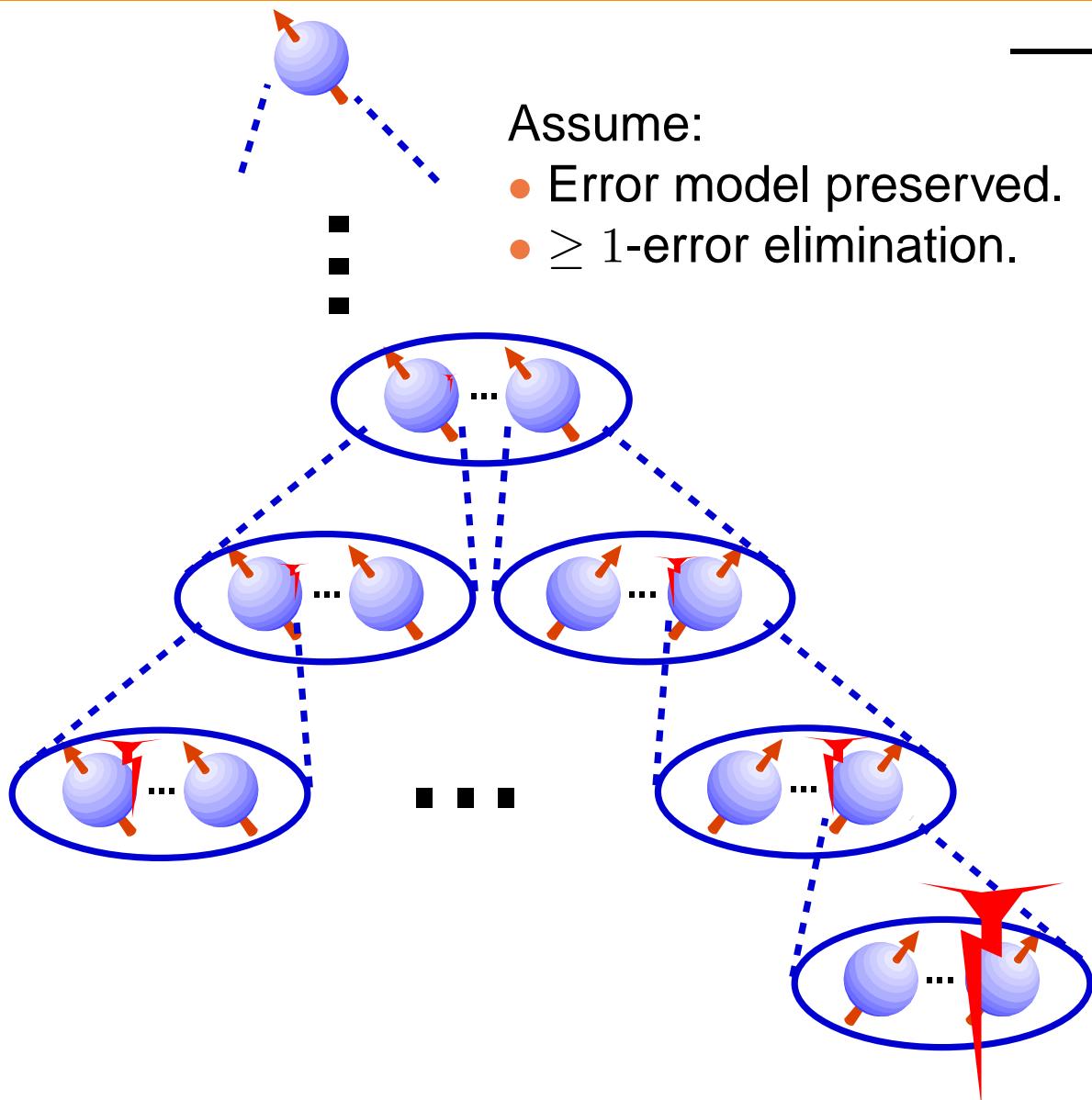
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[Back to: Methods for Scalability](#)

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# Concatenation



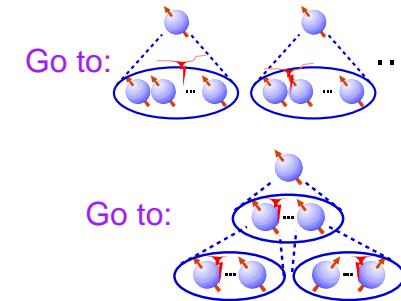
Level	Error rate
$k$	$\leq C^{2^k-1} p^{2^k}$
$\vdots$	
4	$\begin{cases} p_4 \leq C(C^3 p^{2^2})^2 \\ = C^{2^3-1} p^{2^3} \end{cases}$
3	$\begin{cases} p_3 \leq C(C p^2)^2 \\ = C^{2^2-1} p^{2^2} \end{cases}$
2	$p_2 \leq C p^2$
1	$p$

Back to: Methods for Scalability

# Methods for Scalability

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- Error-correcting subsystems.
- Concatenation.



⇒ unbounded distance communication.

# Pauli Product Operator Rotations

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- One qubit rotations and gates:

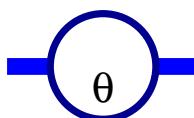
- $X_{90^\circ} = e^{-i\sigma_x\pi/4}$



- $Y_{-90^\circ} = e^{i\sigma_y\pi/4}$

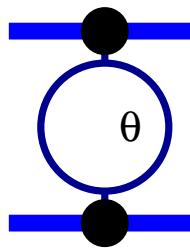


- $Z_\theta = e^{-i\sigma_z\theta/2}$

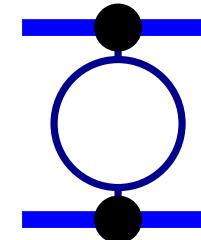


- Two qubit rotations and gates:

- $(Z^{(1)}Z^{(2)})_\theta = e^{-i\sigma_z^{(1)}\sigma_z^{(2)}\theta/2}$

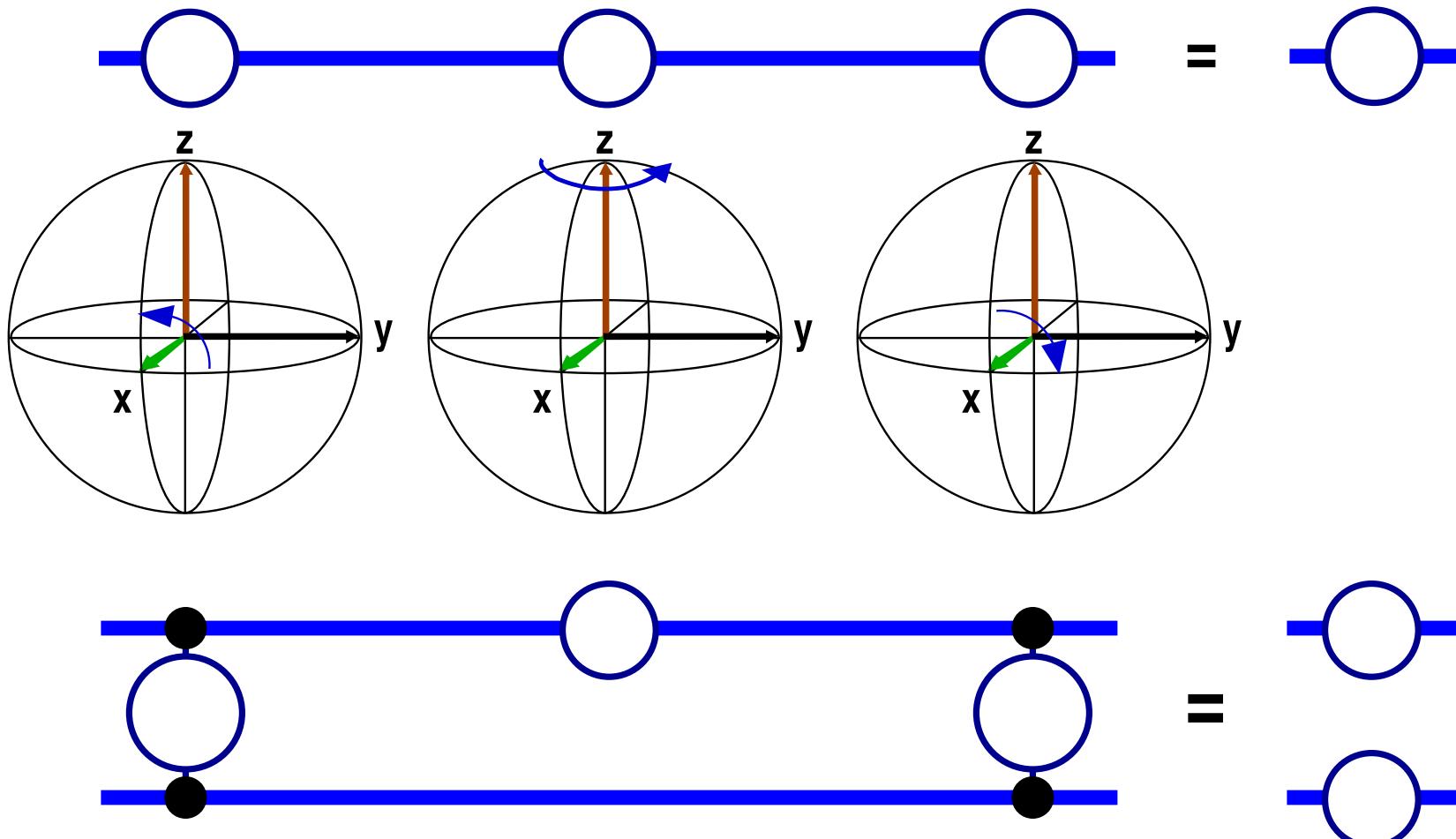


- $(Y^{(1)}Z^{(2)})_{90^\circ} = e^{-i\sigma_y^{(1)}\sigma_z^{(2)}\pi/2}$

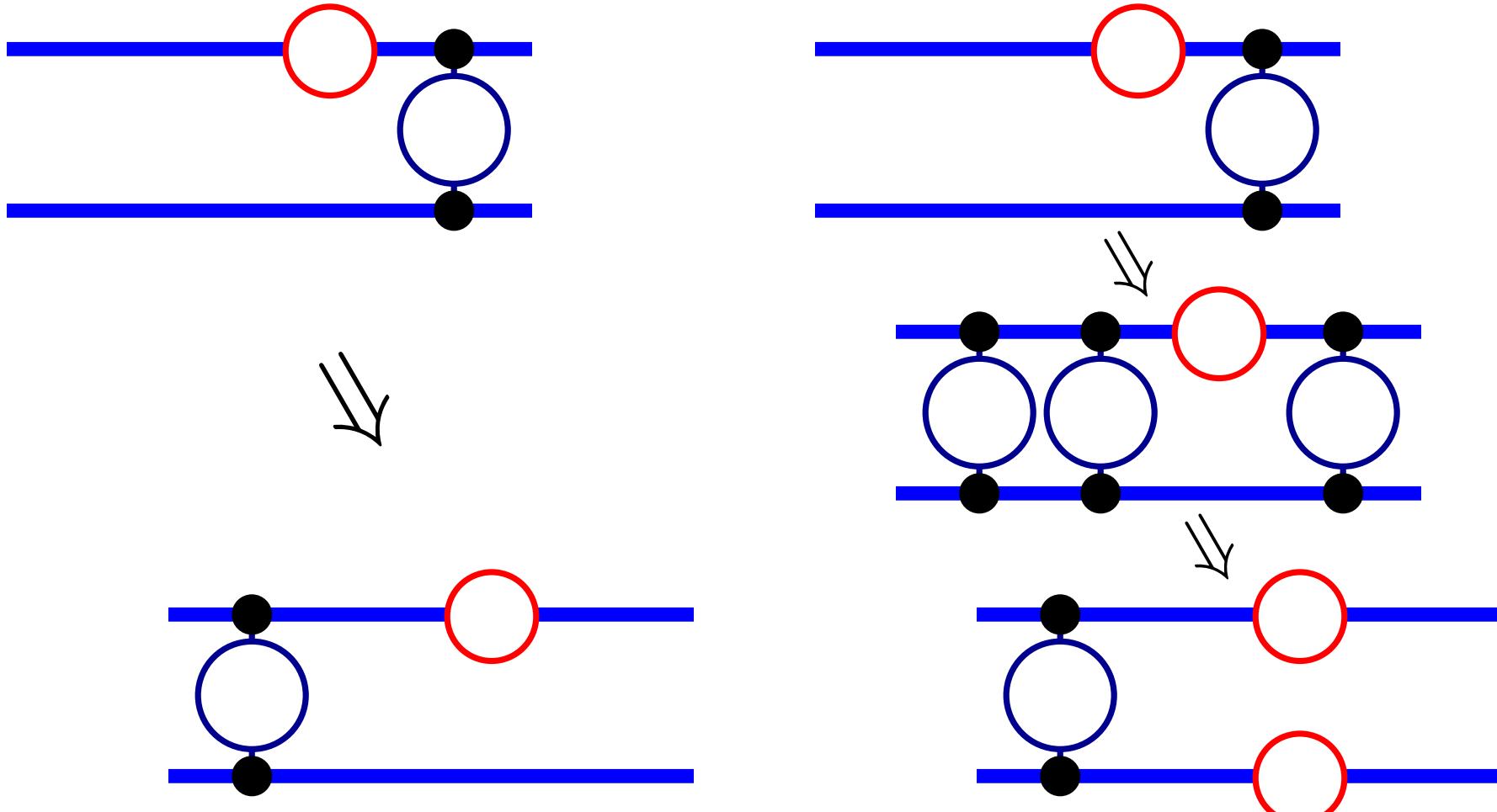


# Rotation Equivalences

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# Error Propagation



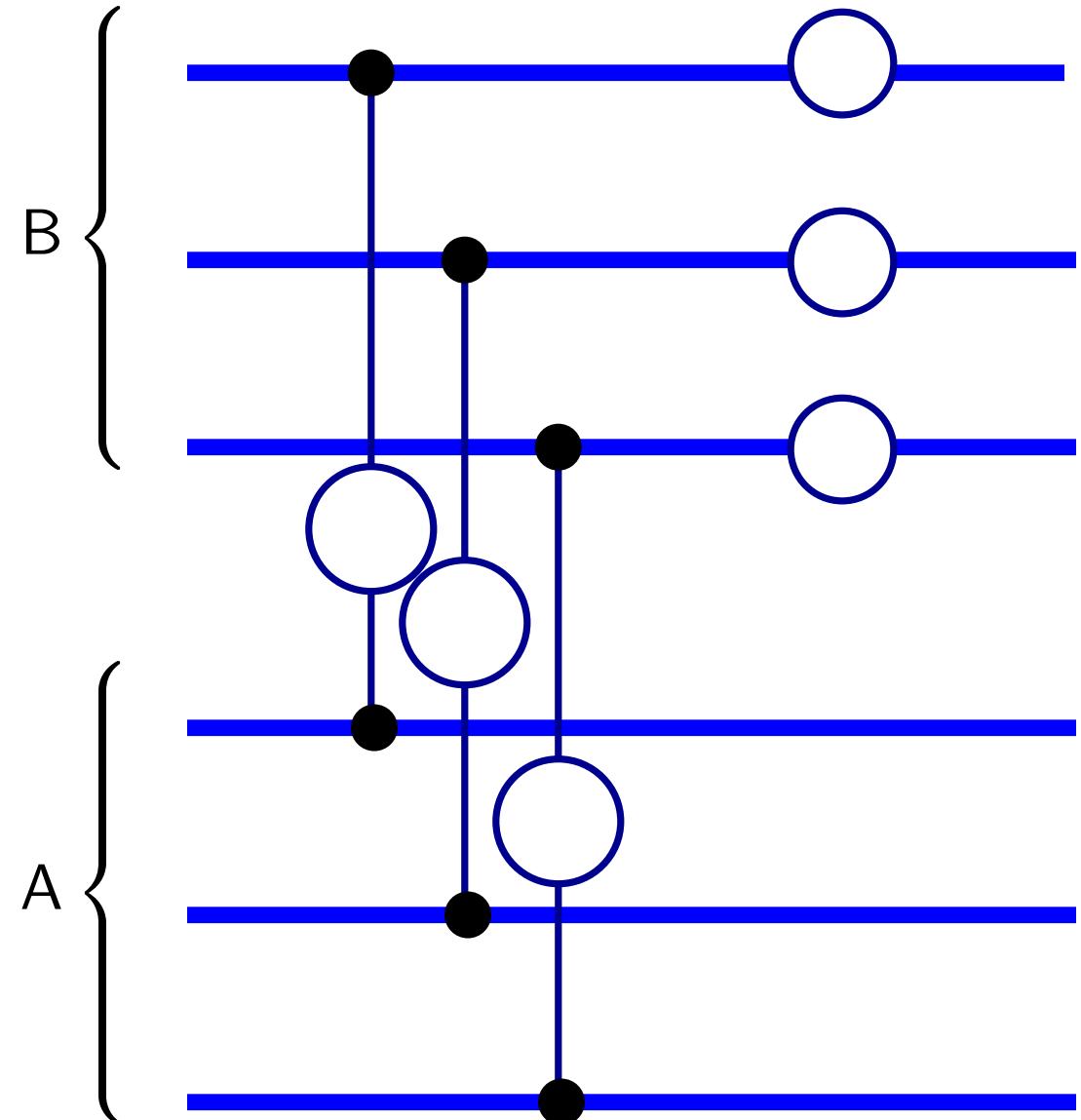
# Transversally Encoding Operations

Repetition code for bit  
flips.

Stabilizer:  $\langle Z Z I, I Z Z \rangle$

$|0\rangle_L = |000\rangle$ ,  $|1\rangle_L = |111\rangle$

- “Encoded” cnot ...



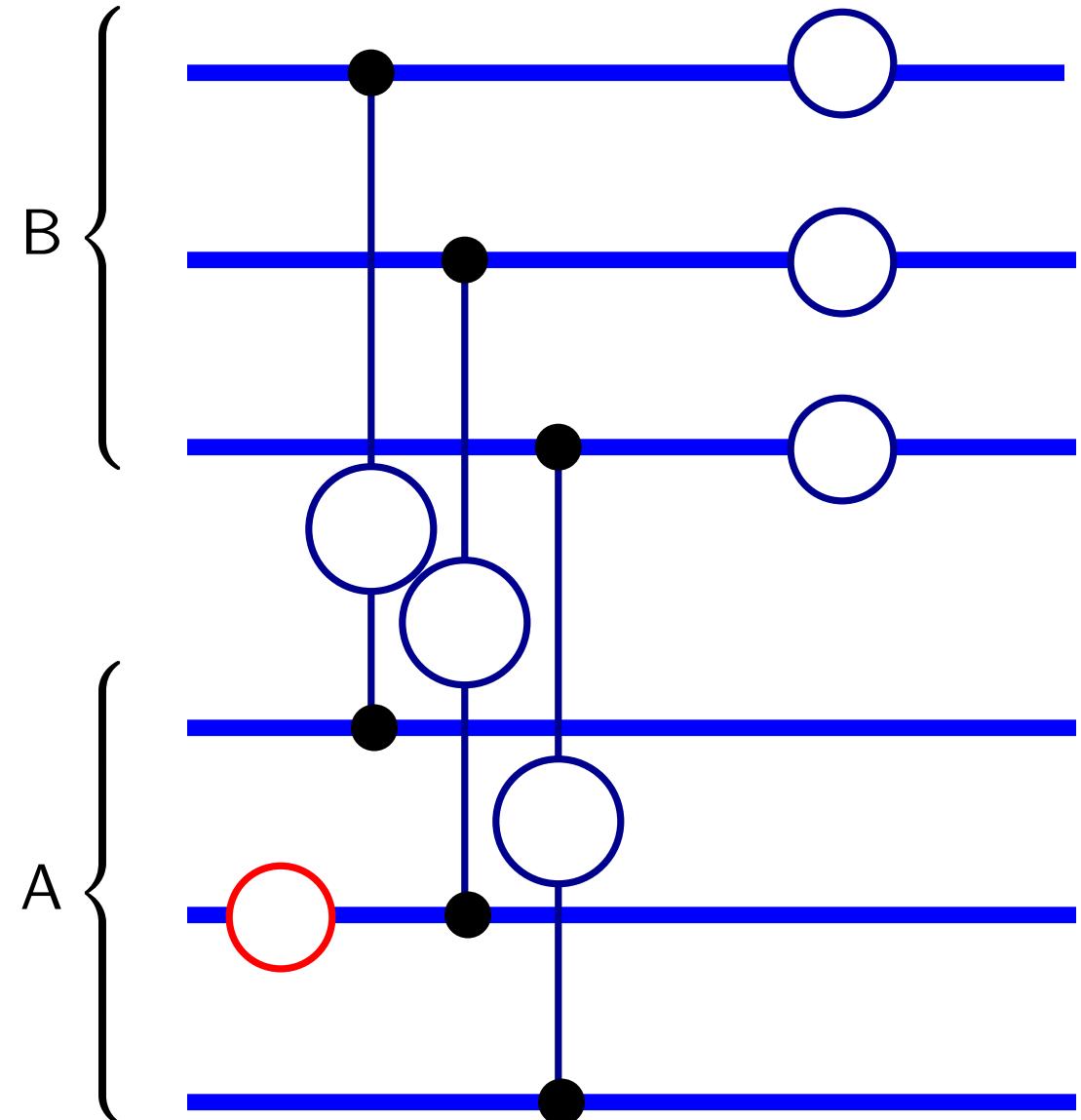
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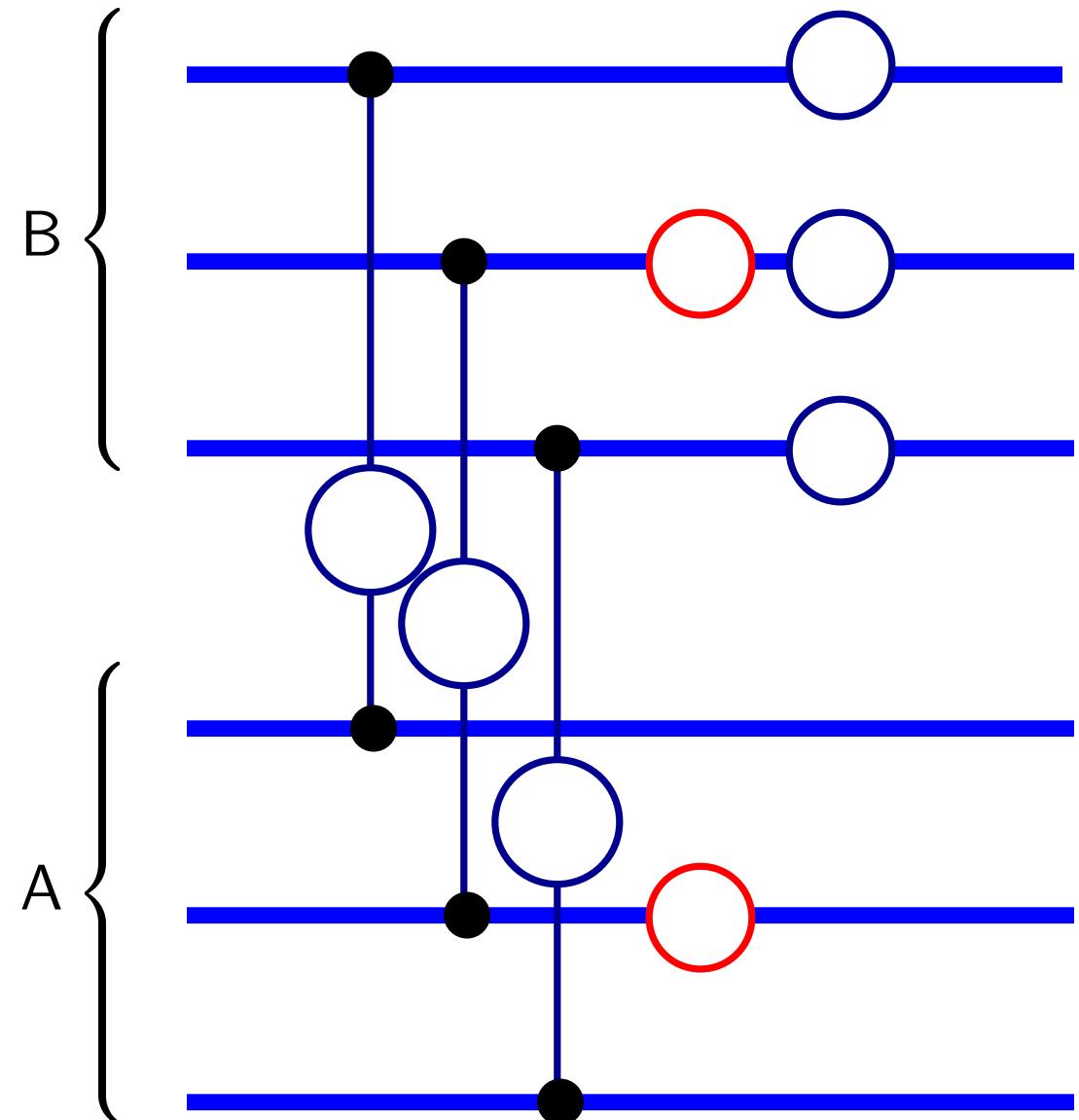
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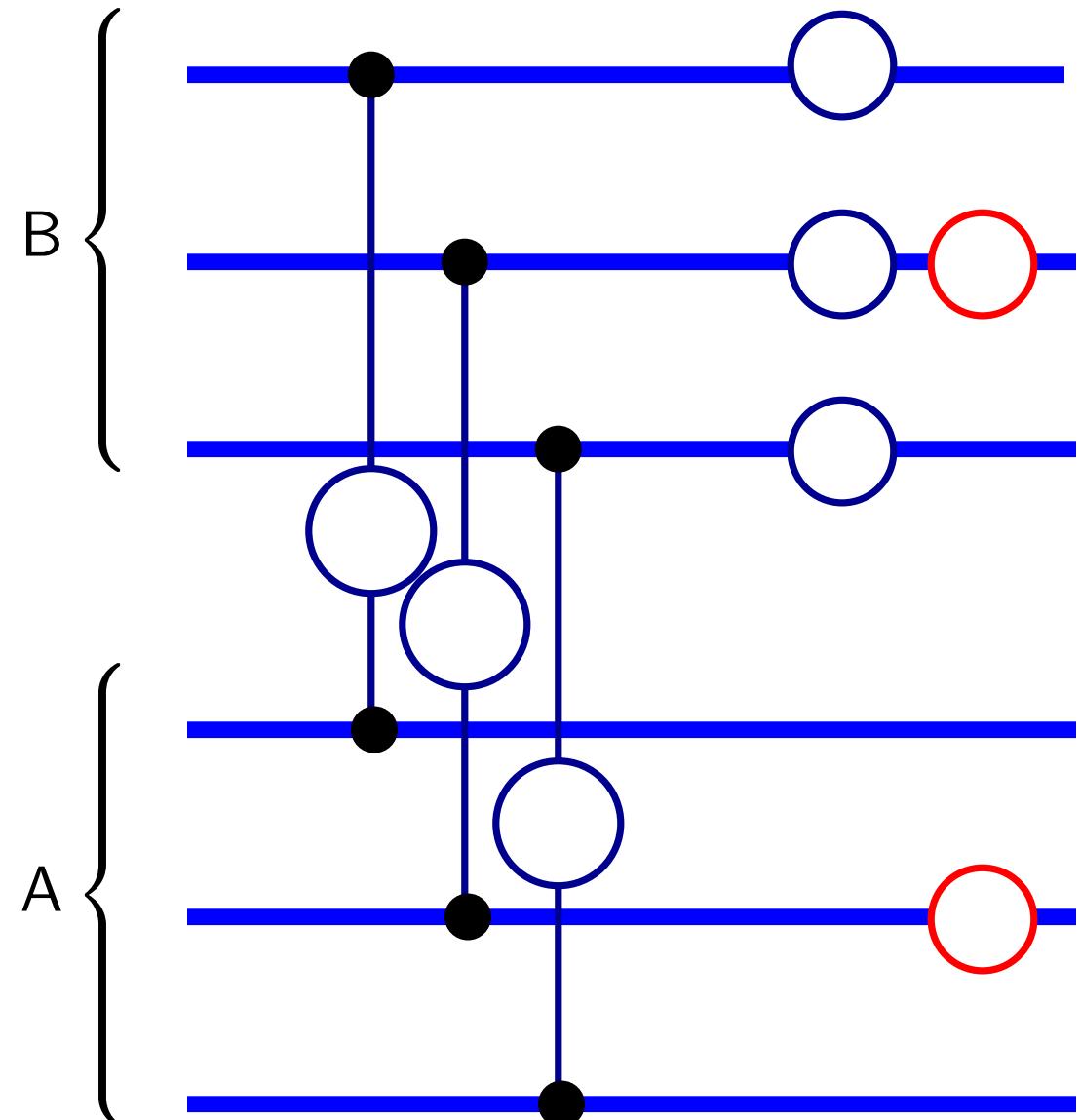
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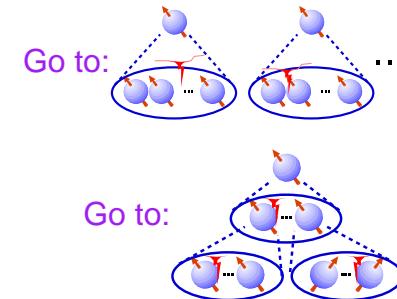
- “Encoded” cnot ...



Back to: Methods for Scalability

# Methods for Scalability

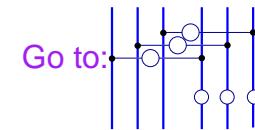
- Error-correcting subsystems.



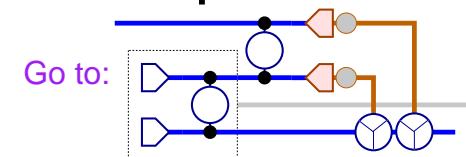
- Concatenation.

⇒ unbounded distance communication.

- Transversally encoded operations.



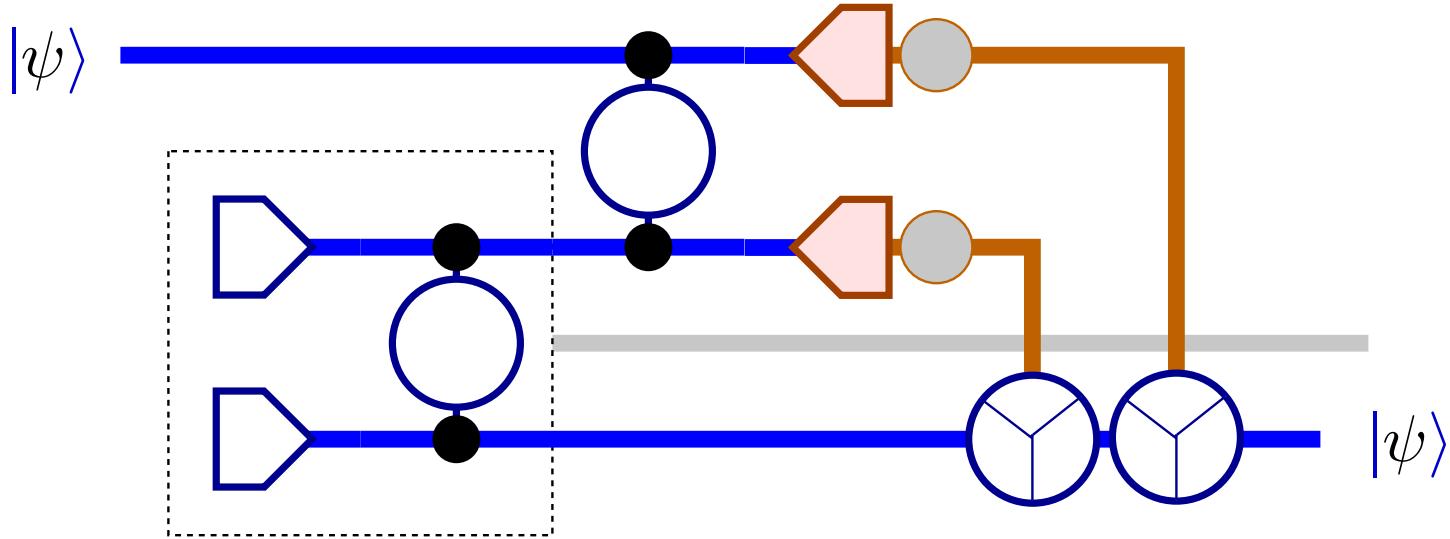
- Measurement or Operation  
= state preparation + teleportation.



- Robust error detection and recovery.

Jump to: Conclusion

# Teleportation



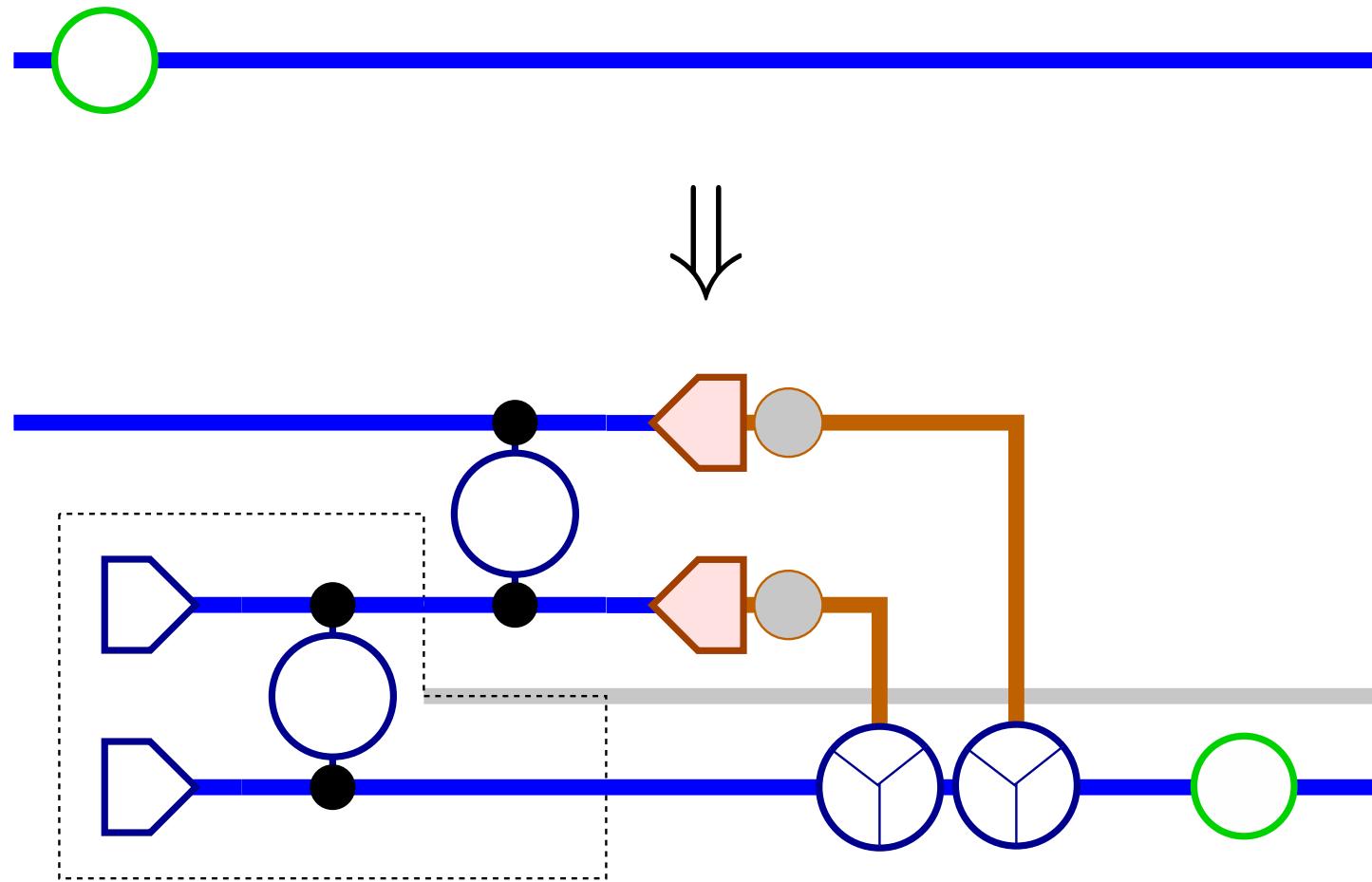
$$a = +1 \Rightarrow X_{180^\circ} , \quad a = -1 \Rightarrow \mathbb{I}$$

and

$$b = +1 \Rightarrow Z_{180^\circ} , \quad b = -1 \Rightarrow \mathbb{I}$$

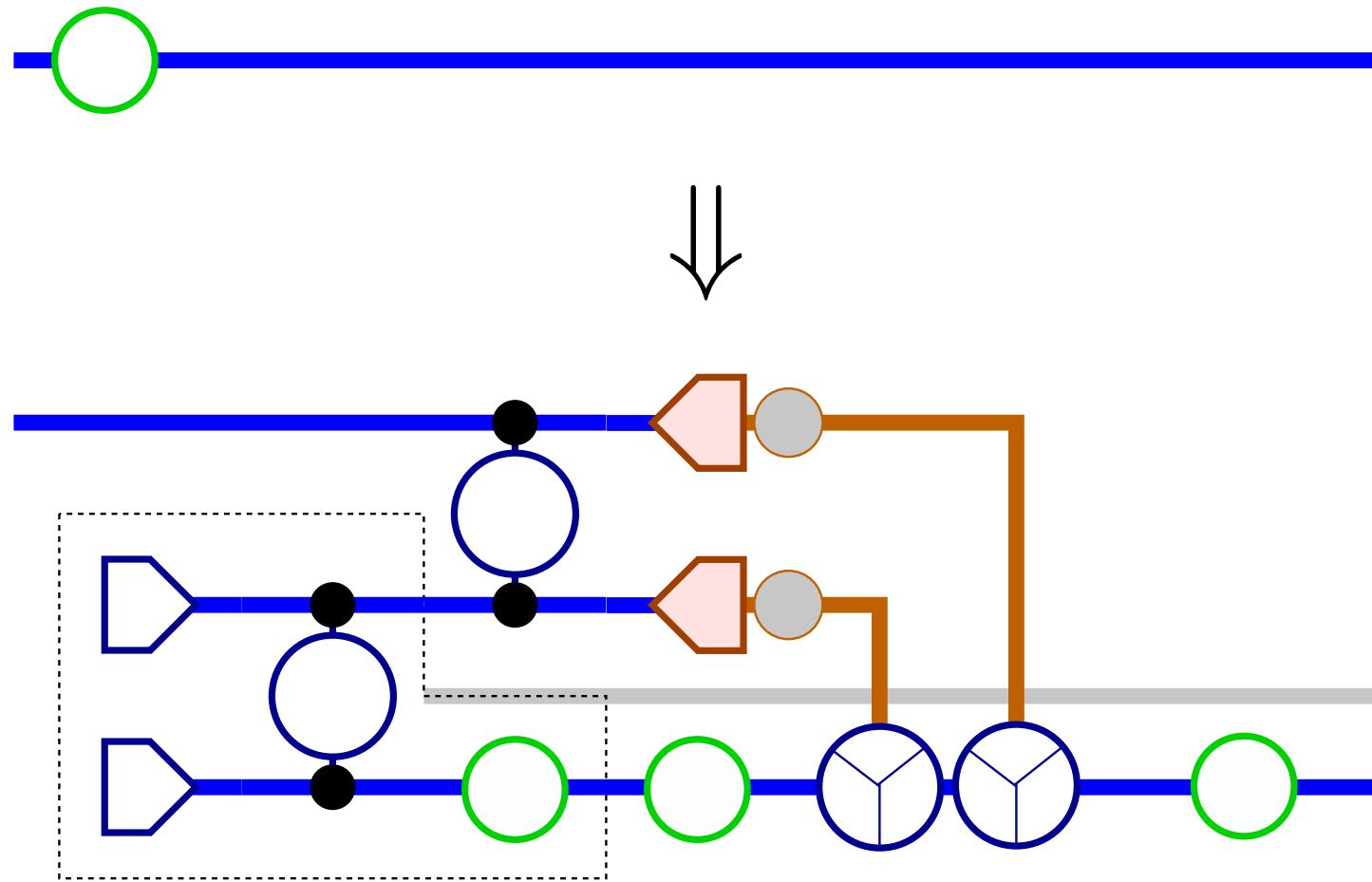
# Operation = Preparation + Teleportation I

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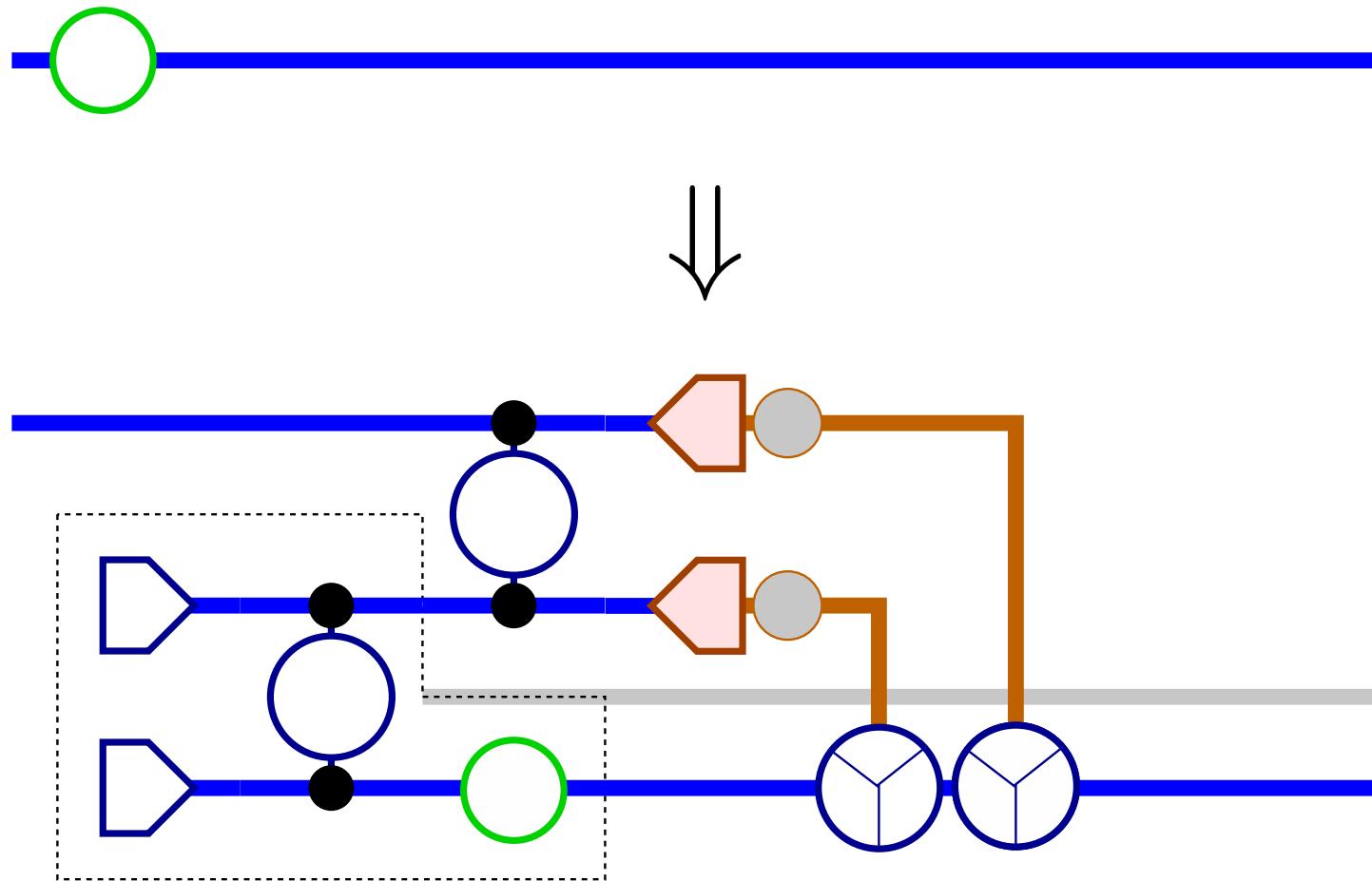
# Operation = Preparation + Teleportation I

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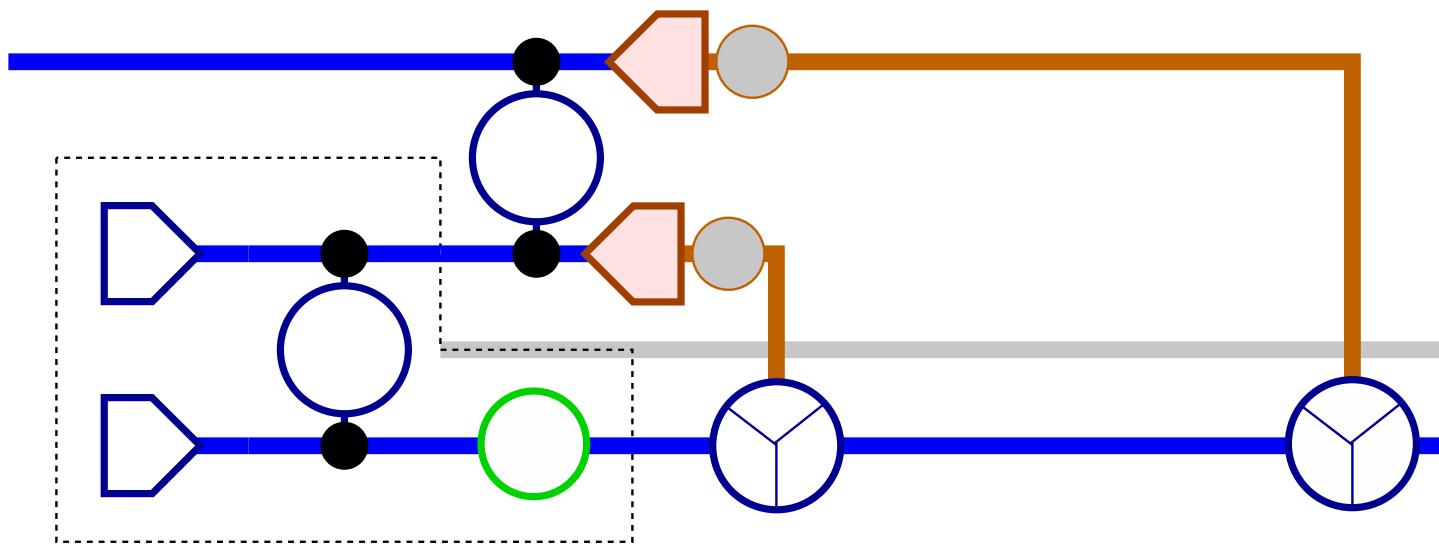
# Operation = Preparation + Teleportation I

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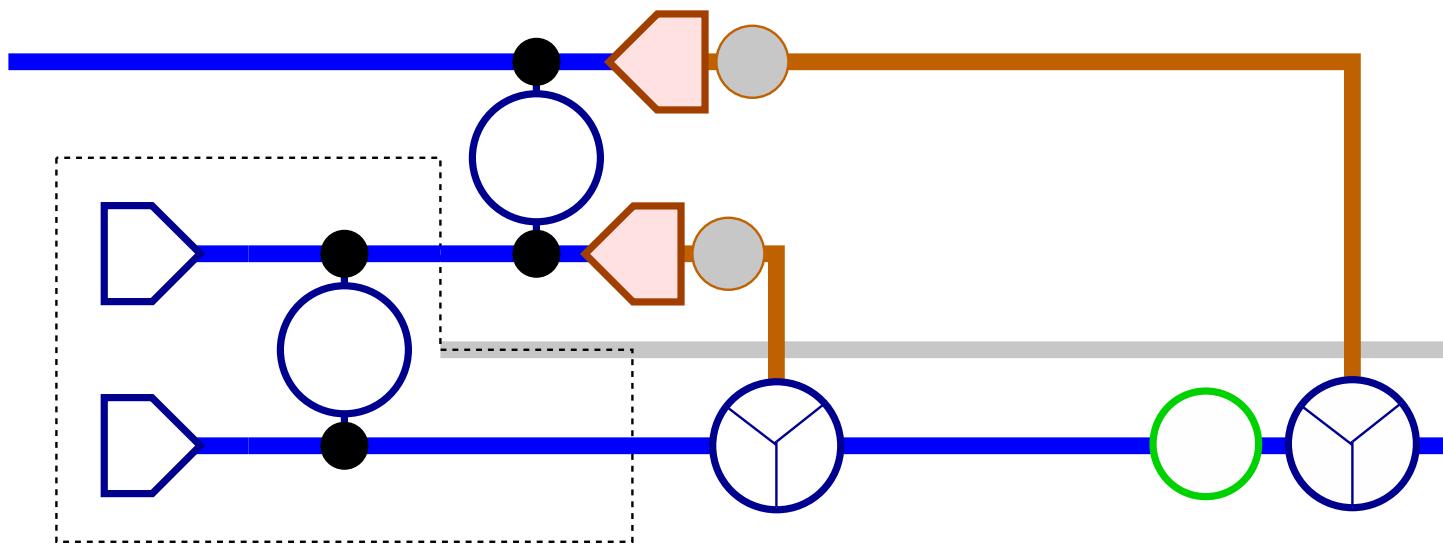
# Operation = Preparation + Teleportation II

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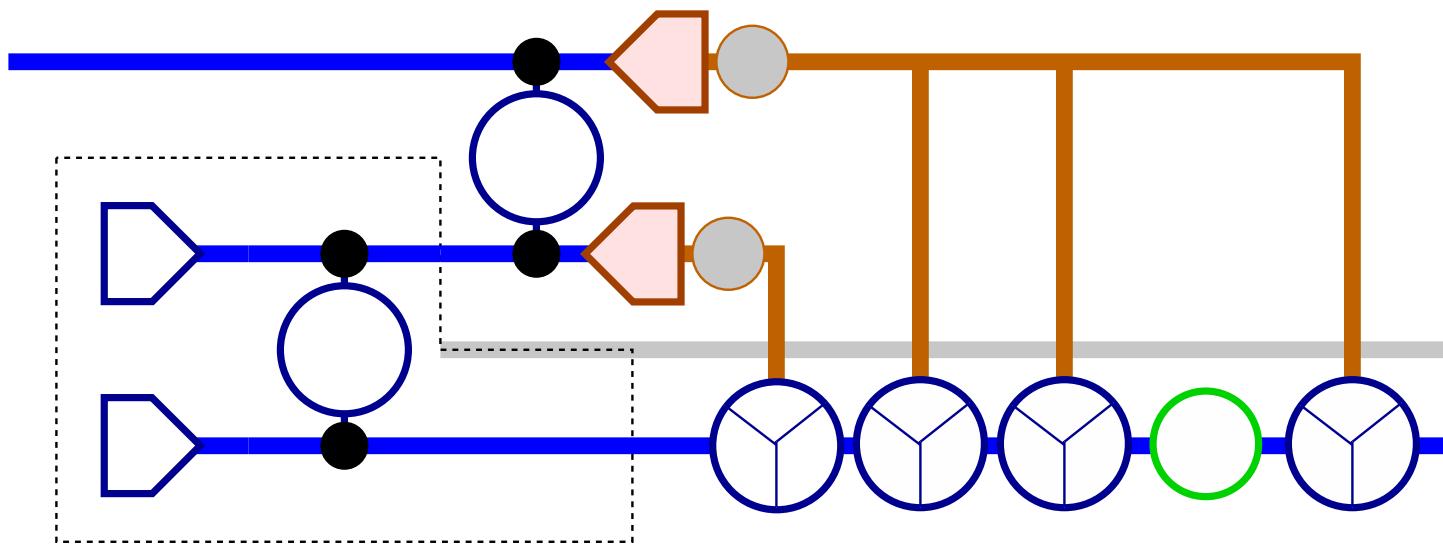
# Operation = Preparation + Teleportation II

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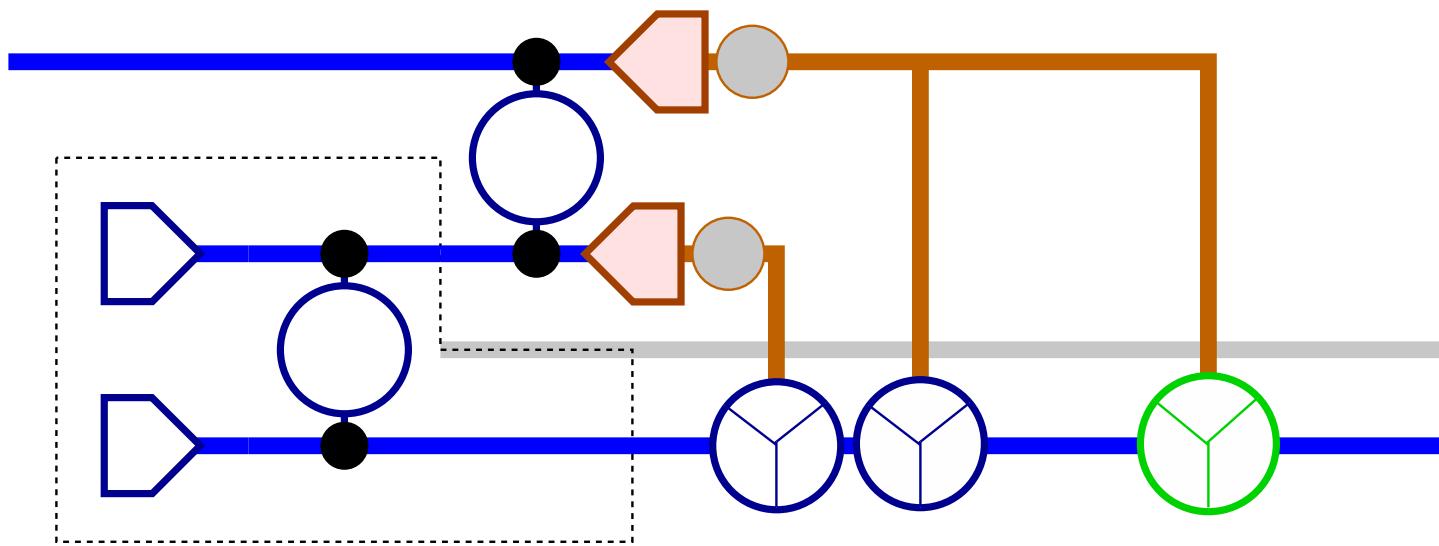
# Operation = Preparation + Teleportation II

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# Operation = Preparation + Teleportation II

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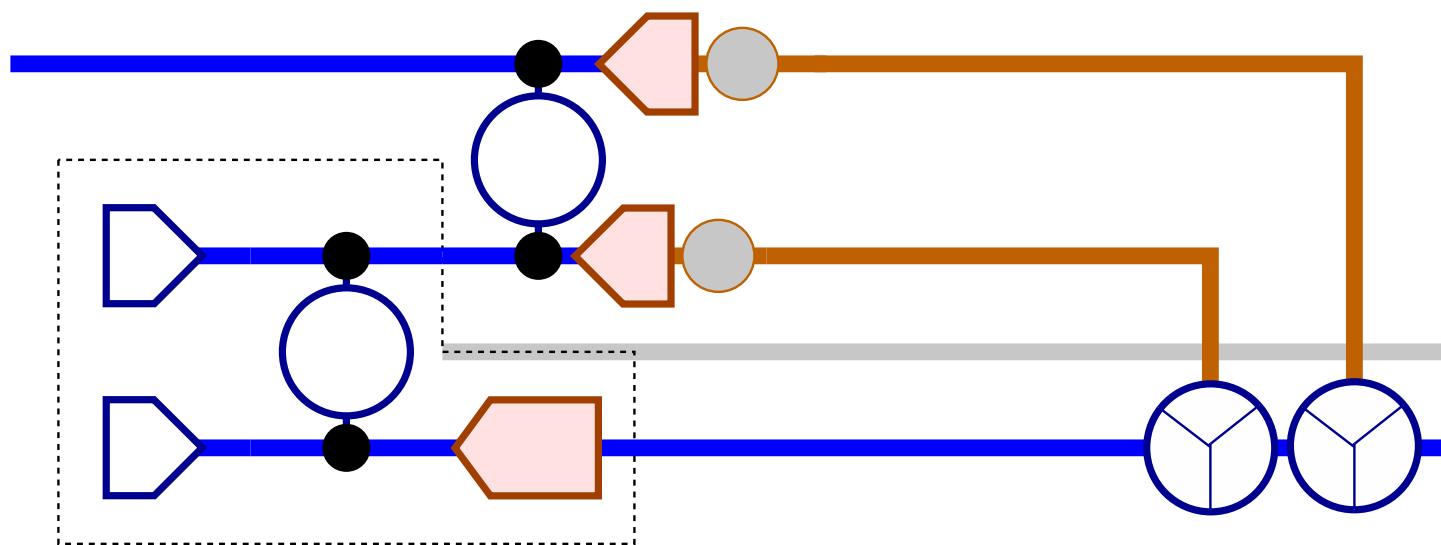


To: Advantages of Teleportation

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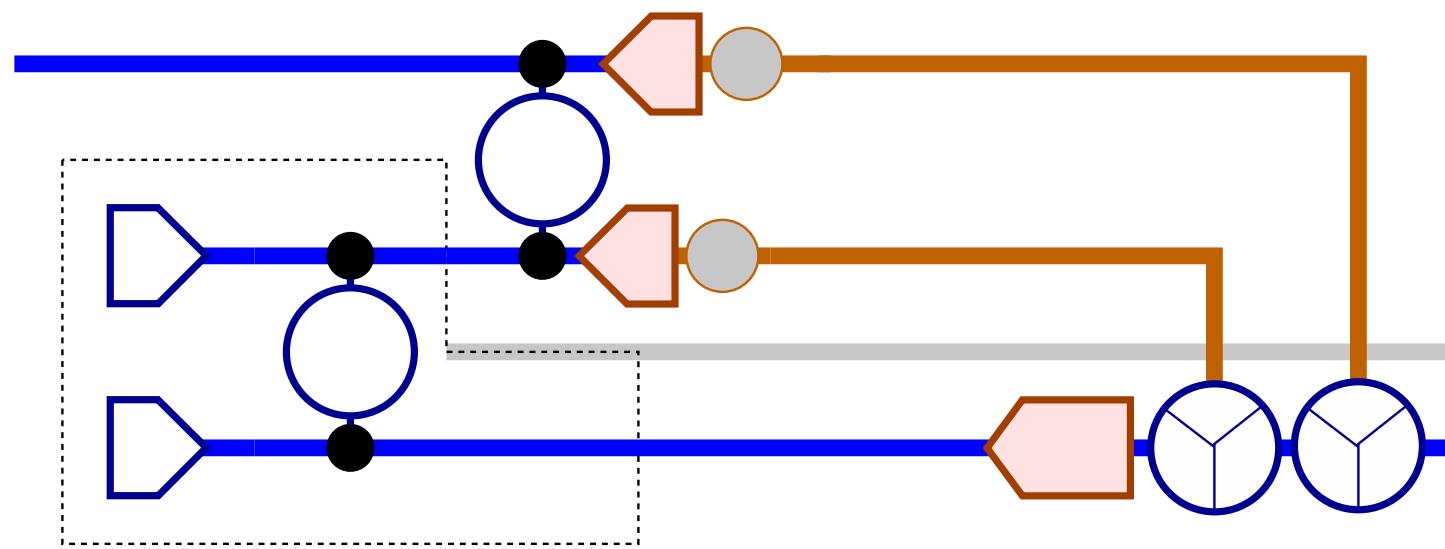
# Measurement = Preparation + Teleportation

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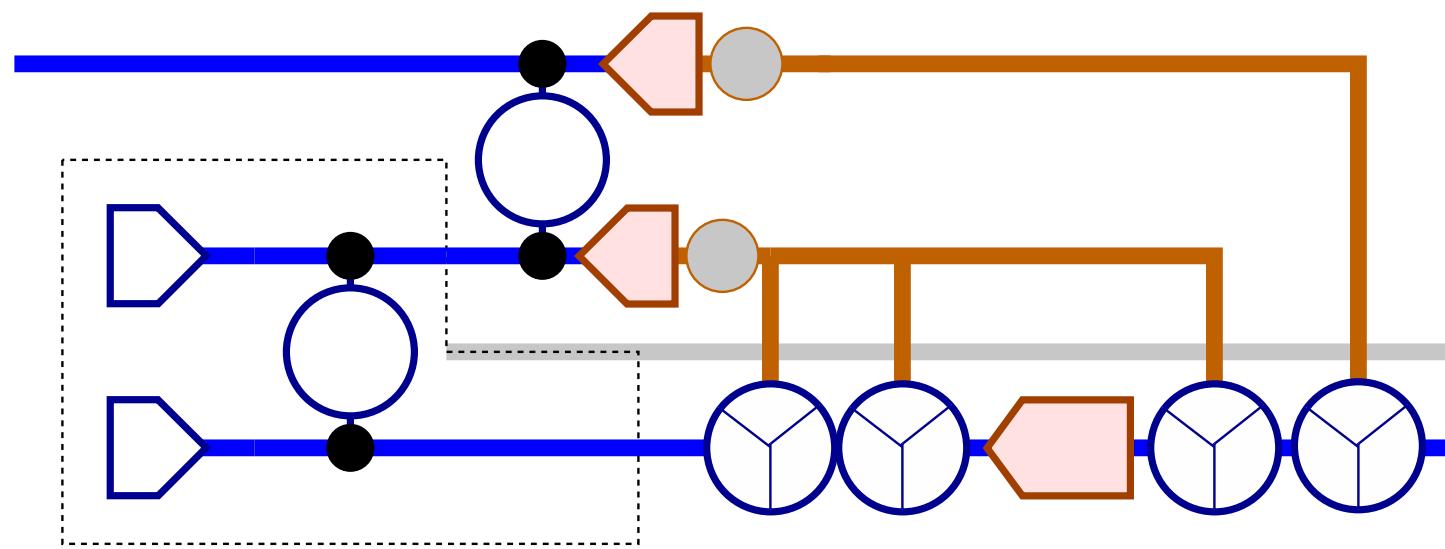
# Measurement = Preparation + Teleportation

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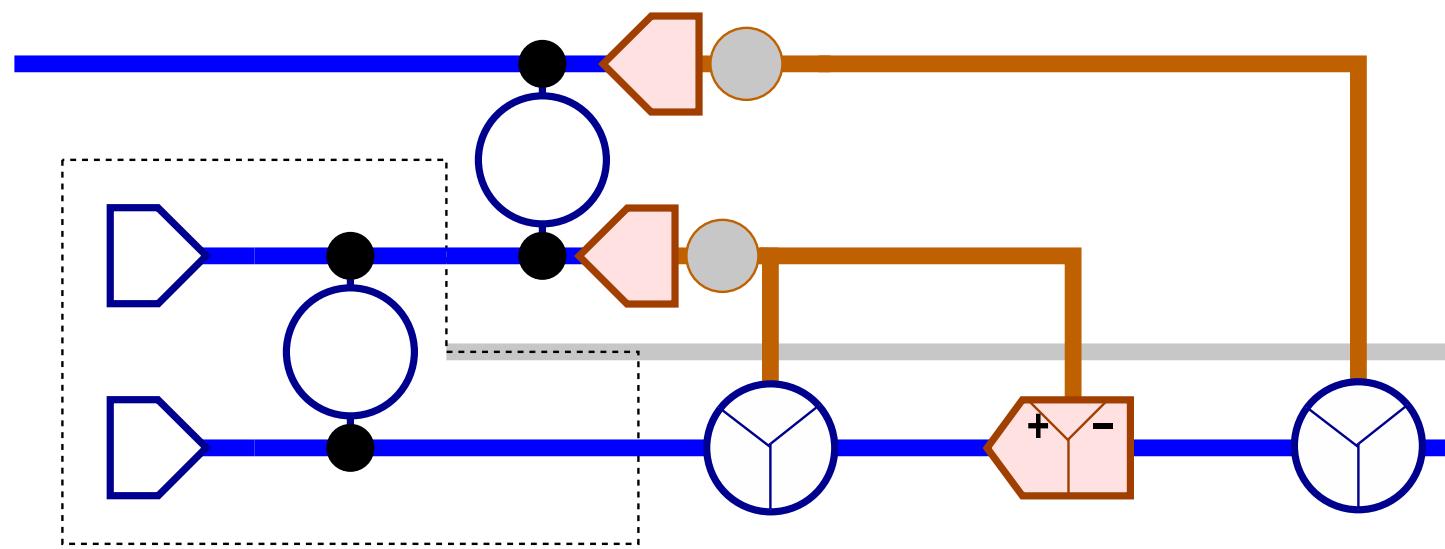
# Measurement = Preparation + Teleportation

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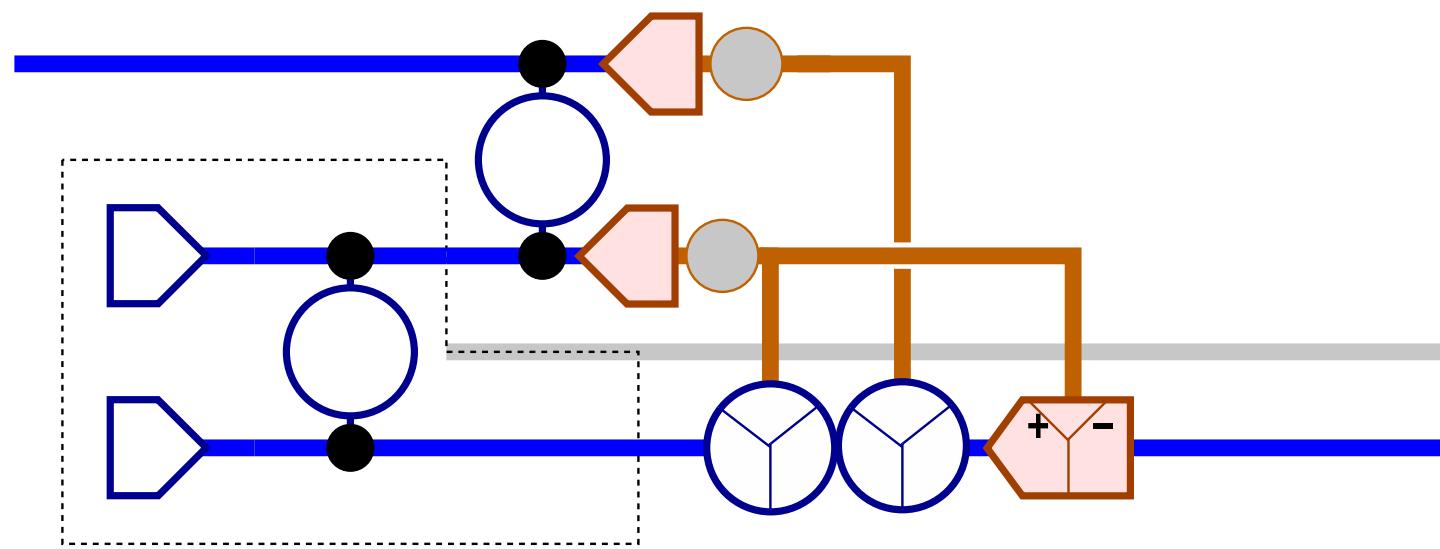
# Measurement = Preparation + Teleportation

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# Measurement = Preparation + Teleportation

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# Advantages of Teleportation

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- Transversal after successful state preparation.
  - Fault tolerant universality.
  - Robust syndrome detection for recovery from error.
- Good error detection suffices.
  - Reject attempted state preparations if errors are detected.

[Back to: Methods for Scalability](#)

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# Conclusion

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- Accuracy threshold questions:
  - Bit flip error model?
  - Erasure error model?
  - Depolarizing error model?
- Worst-case dimension of a maximum size error-detecting quantum code in an  $n$ -dimensional space subject to an  $e$ -dimensional error model?  
(Best bound known:  $\Omega(n/e^2)$ .)

# References

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